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## Geometric entanglement in topologically ordered states

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### Abstract

Here we investigate the connection between topological order and the geometric entanglement, as measured by the logarithm of the overlap between a given state and its closest product state of blocks. We do this for a variety of topologically ordered systems such as the toric code, double semion, colour code and quantum double models. As happens for the entanglement entropy, we find that for sufficiently large block sizes the geometric entanglement is, up to possible sub-leading corrections, the sum of two contributions: a bulk contribution obeying a boundary law times the number of blocks and a contribution quantifying the underlying pattern of long-range entanglement of the topologically ordered state. This topological contribution is also present in the case of single-spin blocks in most cases, and constitutes an alternative characterization of topological order for these quantum states based on a multipartite entanglement measure. In particular, we see that the topological term for the two-dimensional colour code is twice as much as the one for the toric code, in accordance with recent renormalization group arguments (Bombin *et al* 2012 *New J. Phys.* **14** 073048). Motivated by these results, we also derive a general formalism to obtain

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upper- and lower-bounds to the geometric entanglement of states with a non-Abelian group symmetry, and which we explicitly use to analyse quantum double models. Furthermore, we also provide an analysis of the robustness of the topological contribution in terms of renormalization and perturbation theory arguments, as well as a numerical estimation for small systems. Some of the results in this paper rely on the ability to disentangle single sites from the quantum state, which is always possible for the systems that we consider. Additionally we relate our results to the behaviour of the relative entropy of entanglement in topologically ordered systems, and discuss a number of numerical approaches based on tensor networks that could be employed to extract this topological contribution for large systems beyond exactly solvable models.

## 1. Introduction

Topological order (TO) [1] is an example of new physics beyond Landau's symmetry-breaking paradigm of phase transitions. Systems exhibiting this new kind of order are linked to concepts of the deepest physical interest, e.g. quasiparticle anyonic statistics and topological quantum computation (see e.g. [2]). Importantly, TO finds a realization in terms of topological quantum field theories [3], which are the low-energy limit of quantum lattice models such as the toric code and quantum double models [4], as well as string-net models [5].

A remarkable property about TO is that it influences the long-range entanglement in the wave function of the system. For instance, as proven for systems in two spatial dimensions (2D) [6, 7], the entanglement entropy  $S$  of a region, such as  $A$  in figure 1(a), of boundary size  $L \gg 1$  obeys the law

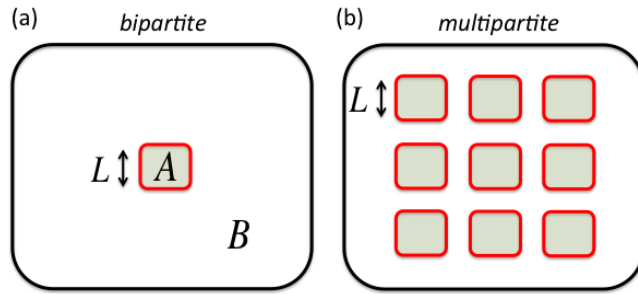
$$S = S_0 - S_\gamma + O(L^{-\nu}), \quad (1)$$

where  $S_0 \propto L$  is a non-universal term (the so-called 'boundary law' or 'area law'),  $\nu$  some exponent and  $S_\gamma$  is a universal long-distance contribution: the topological entanglement entropy.  $S_\gamma$  is non-zero for systems with TO, e.g.  $S_\gamma = 1$  for systems in the topological phase of the toric code. In general one has that

$$S_\gamma = \log \left( \sqrt{\sum_a d_a^2} \right), \quad (2)$$

where  $\{d_a\}$  are the so-called *quantum dimensions* of the associated anyon model. More recently, similar universal contributions have also been found for other bipartite entanglement measures such as the mutual information and the Rényi entropy [8, 9].

The main purpose of this paper is to investigate, in the context of systems with TO, a global measure of entanglement which captures multipartite correlations in the system: *the geometric entanglement (GE)* [10, 11], which we shall denote by  $E_G$ . This measure intuitively characterizes how well an entangled state can be approximated by a mean-field state, i.e. a product state, and its behaviour is in principle very different from that of bipartite measures such as the entanglement entropy. Can such a mean-field, global measure of entanglement be able to reveal any property about TO? As we shall see in this paper, this is indeed the case.



**Figure 1.** Two different types of partitions. (a) Bipartition with region  $A$  and region  $B$ . Region  $A$  has linear size  $L$ . This is usually the partition used in evaluations of the entanglement entropy and other bipartite entanglement measures. (b) Multipartition with many blocks, each one with linear size  $L$ . This is the partition used in the calculations of the GE. If each block contains only one spin (akin to individual degree of freedom), then the product state is a completely separable state and the GE is defined with respect to individual spins. However, each block can also contain multiple spins, and in this case the GE is defined with respect to the closest product state of the individual blocks. In this latter scenario, we can inquire how the entanglement depends on the block boundary size  $L$  and, in particular, whether there is something like a boundary law (i.e. linear scaling with  $L$ ) and whether there is a topological correction.

Most of our study is focused on exactly solvable models corresponding to fixed-points of the renormalization group (RG). We deal most prominently with the toric code model [4], but shall also discuss the double semion model [5], topological colour codes [12] as well as quantum double models [4]. Each one of these models can be recast as some particular instance of a string-net (Levin–Wen) model. Because of this, they correspond (by definition) to RG fixed-points and, therefore, they are representative of their respective topological phases. Our main result is that, for all fixed-point models considered in this paper, the GE obeys

$$E_G = E_0 - E_\gamma \quad (3)$$

with  $E_\gamma$  being identical to  $S_\gamma$  and  $E_0$  some bulk contribution. The exact form of  $E_0$  depends on short-distance details such as the size of the blocks (i.e. the number of spins) considered for the closest product state optimization. As we shall see, for *blocks* of boundary size  $L$ ,  $E_0$  will be proportional to  $L$  and also to the number of blocks  $n_b$  (more specifically,  $E_0 \propto n_b L$ ). We call this behaviour, which depends on short-distance correlations, a ‘boundary law’ for the GE, since it corresponds to a boundary term for each one of the blocks, i.e.  $E_0/n_b \sim L$ . Moreover,  $E_\gamma$  being identical to  $S_\gamma$  is an indication of the topological origin of the term  $E_\gamma$ .

Although we do not have a proof that equation (3) holds generically beyond the considered fixed-point states, we shall argue that the above law is also valid up to sub-leading corrections away from the RG fixed points, e.g. for the toric code model under the influence of non-relevant (and weak) perturbations such as magnetic fields. In this case the GE of blocks with boundary size  $L \gg 1$  seems to obey the law

$$E_G = E_0 - E_\gamma + O(L^{-\nu'}), \quad (4)$$

where again  $E_0 \propto n_b L$  is a bulk contribution,  $\nu'$  some exponent and  $E_\gamma$  the topological contribution. From here on, we shall refer to the term  $E_\gamma$  as the *topological GE*.

It is worth mentioning now a number of remarks about our results. Firstly, they provide the first evidence that TO can actually be read and assessed in a number of situations from

multipartite nature of entanglement, and, in particular, the geometric measure of entanglement employed here. Thus, this may provide an alternative way of characterizing TO in quantum states. Secondly, our results are also the first explicit and analytic examples of a boundary law behaviour for the GE in ground states of 2D quantum many-body systems. Thirdly, we will see that the topological term for topological colour codes is twice the one for the toric code. This is consistent with the recent result that the former model is equivalent to two copies of the latter [13]. Moreover, there may be potential advantages in considering the GE to characterize TO instead of other quantities. For instance, one of them is that given a (possibly approximate) description of the ground state by a tensor network such as a projected entangled pair state (PEPS) [14] or a multi-scale entanglement renormalization ansatz (MERA) [15], its computation is not too costly. This is because, in the context of tensor network methods, the calculation mainly involves overlaps with product states and optimizations over these, which can be computed quite efficiently. In this respect, some of the possible approaches using these methods will be discussed in the appendix. Moreover, the GE can also be probed experimentally with current technology using e.g. nuclear magnetic resonance via single spin measurement on all nuclear spins [16], or ultra-cold atoms in optical lattices via similar techniques to those explained in [17].

The structure of the paper is as follows. First, we briefly review the basics of the GE in section 2. After this, we deal in section 3 with the toric code model. Since this is the simplest model displaying non-trivial TO, we spend quite some time explaining many of its properties, as well as derivations of the GE of spins and blocks both for the square and honeycomb lattices. In section 4 we explain a number of important points that need to be considered in order to generalize the results obtained for the toric code to other models. Then, the double semion model is analysed in section 5, and topological colour codes are addressed in section 6. For colour codes we derive results using two independent methods: a bound approach (similar to the one used for the other models), and a direct approach valid for Calderbank–Shor–Steane (CSS) self-orthogonal codes. Depending on the particular setting we will see that one approach may provide some advantages over the other. After this, we develop in section 7 a general formalism for the bound approach, valid for states with non-Abelian symmetries, and apply this formalism to quantum double models in section 8. The situation beyond RG fixed points is considered in section 9, where the robustness under perturbations of the topological contribution is assessed by RG and perturbation theory arguments, as well as by small-size numerical calculations. We present our conclusions and possible future directions in section 10. Finally, in the appendix we describe how our results are related to those of the relative entropy of entanglement for topologically ordered states, and discuss possible numerical algorithms based on tensor networks to extract this topological component for non-exactly solvable models of large size.

## 2. Geometric entanglement

To begin with, let us recall some basic notions about the GE. Consider an  $m$ -partite normalized pure state  $|\Psi\rangle \in \mathcal{H} = \bigotimes_{i=1}^m \mathcal{H}^{[i]}$ , where  $\mathcal{H}^{[i]}$  is the Hilbert space of party  $i$ . For instance, in a system of  $n$  spins each party could be a single spin, so that  $m = n$ . But each party could also be a set of spins, either contiguous (a *block* [18]) or not.

We wish now to determine how well state  $|\Psi\rangle$  can be approximated by an unentangled (normalized) state of the parties,  $|\Phi\rangle \equiv \bigotimes_{i=1}^m |\phi^{[i]}\rangle$ . The proximity of  $|\Psi\rangle$  to  $|\Phi\rangle$  is captured by

their overlap. The entanglement of  $|\Psi\rangle$  is thus revealed by the maximal overlap [10]

$$\Lambda_{\max}(\Psi) \equiv \max_{\Phi} |\langle \Phi | \Psi \rangle|. \quad (5)$$

The larger  $\Lambda_{\max}$  is, the less entangled is  $|\Psi\rangle$ . Therefore, we quantify the entanglement of  $|\Psi\rangle$  via the quantity

$$E_G(\Psi) \equiv -\log_2 \Lambda_{\max}^2(\Psi), \quad (6)$$

where we have taken the base-2 logarithm, and which gives zero for unentangled states.  $E_G(\Psi)$  is called GE. This quantity has been studied in a variety of contexts, including critical systems and quantum phase transitions [18, 19], quantification of entanglement as a resource for quantum computation [20], local state discrimination [21], and has been recently measured in NMR experiments [16].

To make a connection between the GE and other measures of entanglement, let us consider the case of just two sets of spins. In this case  $E_G(\Psi)$  coincides with the so-called single-copy entanglement between the two sets,

$$E_1(\Psi) = -\log_2 \nu_1(\rho), \quad (7)$$

with  $\nu_1(\rho)$  the largest eigenvalue of the reduced density matrix  $\rho$  of either set [22]. As is well known, this also coincides with the  $\alpha$ -Rényi entropy

$$S_\alpha = \frac{1}{(1-\alpha)} \log(\text{tr}(\rho^\alpha)) \quad (8)$$

for  $\alpha \rightarrow \infty$ .

Unlike other measures of entanglement, the GE offers a lot of flexibility to study multipartite quantum correlations in spin systems. For instance, one can choose each party to be a single spin, but one can also choose blocks of increasing boundary length  $L$  [18] (see figure 1(b) for an example). In fact, studying how the GE changes with  $L$  provides valuable information about how close the system is to a product state under coarse-graining RG transformations.

### 3. Toric code model

As described in the introduction, here we considered a variety of models corresponding to different topological phases. The simplest of these models is the *toric code* [4]. This model is, in fact, the simplest example of a topologically non-trivial 2D system, and our results for this model provide also the grounds for the rest of the models that we shall consider (namely double semion, colour codes and quantum double models). Given its relevance for the rest of the paper, we provide now a short introduction to some of the key aspects of this model. As we shall see, this will be very useful for the forthcoming calculations.

Mathematically speaking, the toric code is the RG fixed point of the topological phase of a  $\mathbb{Z}_2$  gauge theory. The model is equivalent under local transformations (or local *moves*, or local *disentanglers*) to the Levin–Wen string model on a honeycomb lattice [5], which is by definition a RG fixed point. A more precise derivation of this property will be provided later.

To define the toric code, we consider a lattice  $\Sigma$  on a torus. In this paper we will deal with the toric code in the square and honeycomb lattices, yet the model can be defined on arbitrary lattices. Other Riemann surfaces of genus  $g$  could also be considered easily without changing

our conclusions. There are spin-1/2 (qubits) degrees of freedom attached to each link in lattice  $\Sigma$ . The model is described in terms of stars and plaquettes. A star ‘ $s$ ’ is a set of links sharing a common vertex. A plaquette ‘ $p$ ’ is an elementary face on the lattice  $\Sigma$ . For any star  $s$  and plaquette  $p$ , we consider the star operators  $A_s$  and plaquette operators  $B_p$  defined as

$$A_s \equiv \prod_{j \in s} \sigma_x^{[j]}, \quad B_p \equiv \prod_{j \in p} \sigma_z^{[j]}, \quad (9)$$

where  $\sigma_\alpha^{[j]}$  is the  $\alpha$ th Pauli matrix at link  $j$  of the lattice. Let us call respectively  $n_s$ ,  $n_p$  and  $n$  the number of stars (vertices), plaquettes (faces) and links (sites) in lattice  $\Sigma$ . Importantly, star and plaquette operators satisfy the global constraint

$$\prod_s A_s = \prod_p B_p = \mathbb{I}. \quad (10)$$

Therefore, there are  $n_s - 1$  independent star operators and  $n_p - 1$  independent plaquette operators.

With the definitions above, the Hamiltonian of the model reads

$$H_{\text{TC}} = - \sum_s A_s - \sum_p B_p. \quad (11)$$

This Hamiltonian is frustration free (i.e. all the terms in the above sum commute with each other), and can be diagonalized exactly as explained in [4]. The ground level is four-fold degenerate on a torus (4g-fold for a Riemann surface of genus  $g$ ). This degeneracy depends on the underlying topology of lattice  $\Sigma$ , and is already by itself a signature of TO. Moreover, the ground level is a stabilized space of  $\mathcal{G}_s$ , the group of all the possible products of independent star operators, of size  $|\mathcal{G}_s| = 2^{(n_s-1)}$ .

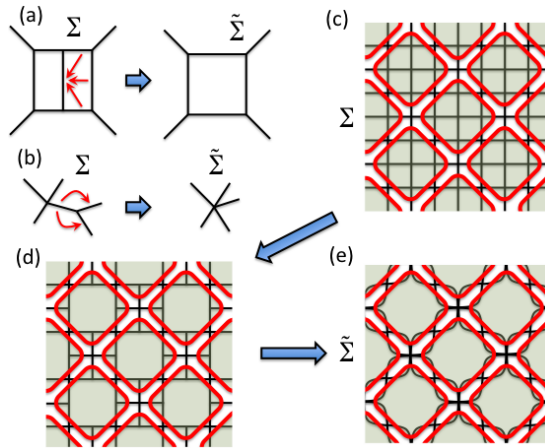
### 3.1. Ground states

In order to build a basis for the ground level subspace, let us consider a closed curve  $\gamma$  running on the links of the dual lattice  $\Sigma^*$ . We define its associated loop operator  $W_x[\gamma] = \prod_{j \in \gamma} \sigma_x^{[j]}$ , where  $j \in \gamma$  are the links in  $\Sigma$  crossed by the curve  $\gamma$  connecting the centres of the plaquettes. It is not difficult to see that the group  $\mathcal{G}_s$  can also be understood as the group generated by all the possible contractible loop operators<sup>7</sup>. Let also  $\gamma_1$  and  $\gamma_2$  be the two non-contractible loops on a torus, and define the associated string operators  $w_{1,2} \equiv W_x[\gamma_{1,2}]$ . We call  $|0\rangle$  and  $|1\rangle$  the eigenstates of  $\sigma_z$  respectively with +1 and -1 eigenvalue. With this assumptions, the ground level subspace then reads  $\mathcal{L} = \text{span}\{|i, j\rangle, i, j = 0, 1\}$ , where

$$|i, j\rangle = \frac{1}{\sqrt{|\mathcal{G}_s|}} \sum_{g \in \mathcal{G}_s} g w_1^i w_2^j |0\rangle^{\otimes n}. \quad (12)$$

It is easy to check that the four vectors  $|i, j\rangle$  are orthonormal and stabilized by  $\mathcal{G}_s$ , that is,  $g|i, j\rangle = |i, j\rangle \forall g \in \mathcal{G}_s$  and  $\forall i, j$ . These four states form a possible basis of the ground level subspace of the toric code model on a torus.

<sup>7</sup> Notice that some sets of non-contractible loops are also elements of  $\mathcal{G}_s$ , such as two parallel non-contractible ones. This is because they can be built from products of contractible loops.



**Figure 2.** (a), (b) CNOT operations that disentangle qubits from the system, thus removing their links from the lattice. In the diagram, the arrows go from controlling to target qubits; (c) example of a partition into blocks of a given boundary  $L$ , e.g.  $L = 12$  here (but all our derivations work for arbitrary  $L$ ). Qubits crossed by the boundary of a block are regarded as being inside the block; (d), (e) doing CNOT operations locally inside of each block we can remove all the stars inside of all the blocks in two steps: firstly, we apply CNOTs as in (a) to get (d) and, secondly, we apply CNOTs as in (b) to get (e). Notice that in (e) each star is made of either 4 or 8 qubits. The procedure works similarly for other blockings.

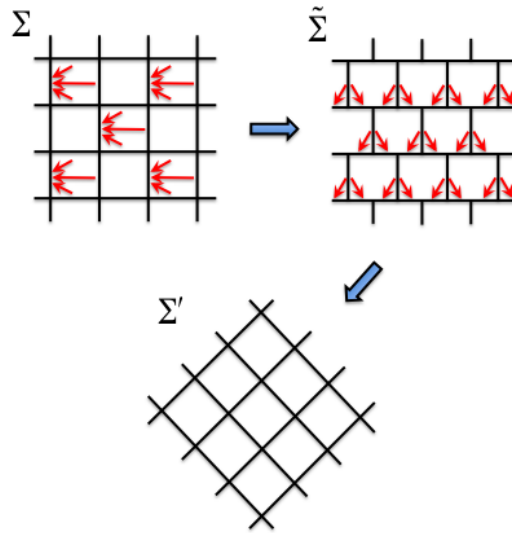
### 3.2. Excited states

Excited states of the toric code can be constructed by locally applying Pauli operators  $\sigma_x$ ,  $\sigma_y$ ,  $\sigma_z$  on the ground states  $|i, j\rangle$ . Pauli operators  $\sigma_z$  create pairs of deconfined charge–anticharge quasiparticle excitations,  $\sigma_x$  create pairs of deconfined flux–antiflux quasiparticles and  $\sigma_y$  creates a flux–antiflux and a charge–anticharge pairs. These operators can be applied to several sites of the lattice, thus creating a quasiparticle pattern that defines the excited state. The excited states of the model are then labelled by a set of quantum numbers,  $|\phi, c, i, j\rangle$ , where  $\phi$  and  $c$  are patterns denoting the position of flux-type and charge-type excitations respectively, and  $i, j$  label the ground state  $|i, j\rangle$  that was excited. In this paper we will focus on the entanglement properties of states  $|\phi, c, i, j\rangle$ .

### 3.3. Disentangling the toric code

A key property of the toric code model, which is fundamental for some of the derivations in this paper, is that its ground state  $|0, 0\rangle \equiv |i = 0, j = 0\rangle$  can be created by a quantum circuit that applies a sequence of controlled-NOT (CNOT) unitary operations over an initial separable state of all the qubits [23]. This means that it is actually possible to *disentangle* qubits from the ground state of the system simply by reversing the action of these CNOTs. Moreover, these CNOT operations are *local*, i.e. they do not span over delocalized sites in the lattice, but rather act over pairs of not-too-far-neighbouring qubits (e.g. nearest- and next-to-nearest-neighbours). This was the key observation that allowed to build an exact MERA representation of the ground states  $|i, j\rangle$  of the model [15]. The two fundamental disentangling moves are represented in figures 2(a) and (b), and leave the overall quantum state as a product state of the disentangled





**Figure 3.** Mapping of the toric code from a square lattice  $\Sigma$  to a honeycomb lattice  $\tilde{\Sigma}$  (Levin–Wen string model), and back to a square lattice  $\Sigma'$ . Each red arrow represents a CNOT operation that disentangles qubits from the system, as explained in figure 2. The arrows go from controlling to target qubits.

qubits with the rest of the system. What is more, the rest of the system is left in the ground state of a toric code model on a deformed lattice  $\tilde{\Sigma}$ , where  $\tilde{\Sigma}$  is obtained from  $\Sigma$  by removing the links that correspond to the disentangled qubits. This property turns out to be of great importance for some of our derivations.

### 3.4. Geometric entanglement of the toric code

We are now in position to study the GE of the toric code model. Given that this is a paradigmatic model of TO, we will perform a detailed analysis. Firstly, we will provide a number of lemmas and theorems that will bring us towards an expression for the GE of the toric code in the square lattice, both for spins and blocks, in a particular basis. Secondly, we will make a similar study to the case of the honeycomb lattice. There we will see that different choices of ground states may give rise to different expressions for the GE of spins with the block size being *unity*, yet the long-wavelength properties are the same and equal to those for the square lattice after *blocking of spins* are considered. This is expected, as both lattices can be mapped between each other by means of CNOT disentangling operations (as shown in figure 3), and therefore they should correspond to the same RG fixed point. Here, we would like to emphasize that blocking is essential as the topological contribution is a long-wavelength property, and, thus, that the behaviour of the GE for sufficiently large block sizes is the same independently of which of the two lattices we choose. (Without blocking, this conclusion cannot be reached.)

**3.4.1. A couple of lemmas.** Let us start by considering two lemmas that will be quite useful in our derivations:

**Lemma 1.** *All the eigenstates  $|\phi, c, i, j\rangle$  of the toric code Hamiltonian have the same entanglement properties.*

**Proof.** Let us choose a reference state from the ground level basis, e.g.  $|0, 0\rangle$ . First of all, notice that the four basis states  $|i, j\rangle$  in the ground level are related to  $|0, 0\rangle$  by local unitary operators, since the string operators  $w_1$  and  $w_2$  are tensor products of  $2 \times 2$  identity operators and Pauli- $x$  matrices. Since these operators act locally on each spin, they do not change the entanglement of the quantum state. Moreover, the excited states  $|\phi, c, i, j\rangle$  are created locally by applying tensor products of  $2 \times 2$  identity and Pauli- $x$  and  $z$  operators to  $|i, j\rangle$ . The states resulting from these operations have then the entanglement properties of state  $|i, j\rangle$ , which in turn are the same as those of the reference state  $|0, 0\rangle$ . Thus, all states  $|\phi, c, i, j\rangle$  have the same entanglement properties.  $\square$

**Lemma 2.** *Consider an arbitrary set of blocks of qubits in lattice  $\Sigma$ , where each block is regarded as an individual party. Then, the entanglement properties of the ground state  $|0, 0\rangle$  are the same as those of state  $|\widetilde{0}, \widetilde{0}\rangle$ , the ground state of the toric code model in the deformed lattice  $\widetilde{\Sigma}$  obtained after disentangling as many qubits as possible using CNOTs locally inside of each block.*

**Proof.** Given the ground state  $|0, 0\rangle$  in  $\Sigma$ , we start by doing CNOT disentangling operations locally inside of each block in order to disentangle as many qubits as possible. As a result, state  $|0, 0\rangle$  transforms into state  $|0, 0\rangle_{\text{disentangled}} \equiv |e_1\rangle \otimes \cdots \otimes |e_p\rangle \otimes |\widetilde{0}, \widetilde{0}\rangle$ , where  $|e_k\rangle$  is the quantum state for the  $k$ th disentangled qubit,  $k = 1, \dots, p$  ( $p$  is the number of disentangled qubits), and  $|\widetilde{0}, \widetilde{0}\rangle$  is the  $i = 0, j = 0$  ground state of a toric code Hamiltonian in the deformed lattice  $\widetilde{\Sigma}$ . Since all CNOT operations are done locally inside of each block, states  $|0, 0\rangle$  and  $|0, 0\rangle_{\text{disentangled}}$  have the same entanglement content if we regard the blocks as individual parties. What is more, the entanglement in the system is entirely due to the qubits in  $|\widetilde{0}, \widetilde{0}\rangle$ , which proves the lemma.  $\square$

**3.4.2. Square lattice.** We are now in position to consider the entanglement properties of the four ground level states  $|i, j\rangle$  as well as the excited states  $|\phi, c, i, j\rangle$ . Our main result here is that, for these states, the GE of blocks consists of a boundary term plus a topological contribution. Keeping the above two lemmas in mind, we now present the following theorem about the GE of spins in the square-lattice toric code:

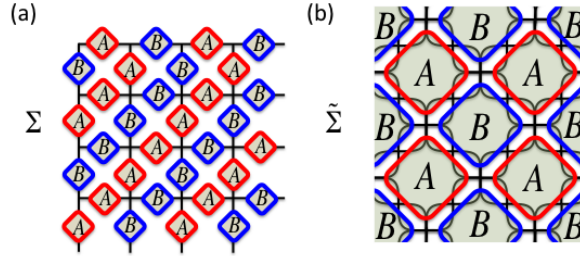
**Theorem 1.** *For the toric code Hamiltonian in a square lattice  $\Sigma$ , the four ground states  $|i, j\rangle$  for  $i, j = 0, 1$  and also the excited states  $|\phi, c, i, j\rangle$  have all the same GE of spins and is given by*

$$E_G = n_s - 1, \quad (13)$$

where  $n_s$  is the number of stars in  $\Sigma$ .

**Proof.** Using lemma 1 we can restrict our attention to the GE for the state  $|0, 0\rangle$ . In this setting, we call *computational basis* the basis of the many-body Hilbert space constructed from the tensor products of the  $\{|0\rangle, |1\rangle\}$  local basis for every spin, which is an example of *product basis* for the spins.

Our proof follows from upper and lower bounding the quantity  $E_G(0, 0)$  in equation (6) for the ground state  $|0, 0\rangle$ . First, from the expression in equation (12) for the ground states we immediately have that the absolute value of the overlap with any state of the computational basis



**Figure 4.** Bipartite lattices of spins and blocks: (a) sets of spins  $A$  and  $B$  for theorem 1 and a  $4 \times 4$  torus; and (b) sets of blocks  $A$  and  $B$  for theorem 2.

is  $|\mathcal{G}_s|^{-1/2}$ , and therefore  $\Lambda_{\max} \geq |\mathcal{G}_s|^{-1/2}$ . From here, we get  $E_G(0, 0) \leq \log_2 |\mathcal{G}_s|$ , which gives an upper bound.

Next, to derive a lower bound we use the fact that if we group the  $n$  spins into two sets  $A$  and  $B$ , then  $\Lambda_{\max} \leq \Lambda_{\max}^{[A:B]}$ , where  $[A : B]$  means that a partition of the system with respect to the two sets  $A$  and  $B$  is considered. Thus, we have that  $E_G(0, 0) \geq E_G(0, 0)^{[A:B]}$ .

The trick to find a useful lower bound is to find an appropriate choice of sets  $A$  and  $B$ . In our case, we consider e.g. the bipartition shown in figure 4(a) for even  $\times$  even lattices (other cases can be considered similarly). Then, we use a lemma by Hamma *et al* in [6], that the reduced density matrix  $\rho_A$  of  $|0, 0\rangle$  for subsystem  $A$  satisfies  $\rho_A^2 = (|\mathcal{G}_s(A)||\mathcal{G}_s(B)|/|\mathcal{G}_s|)\rho_A$ , where  $\mathcal{G}_s(A/B)$  is the subgroup of  $\mathcal{G}_s$  acting trivially on subsystem  $A/B$ . Importantly, for  $A$  and  $B$  chosen as explained above, it happens that  $\mathcal{G}_s(A)$  and  $\mathcal{G}_s(B)$  are *trivial groups consisting of only the identity element*. With this in mind, we see that the reduced density matrix  $\rho_A$  has eigenvalues either zero or  $|\mathcal{G}_s|^{-1}$ , which is  $|\mathcal{G}_s|$ -fold degenerate. This means that  $(\Lambda_{\max}^{[A:B]})^2 = |\mathcal{G}_s|^{-1}$  and hence  $E_G(0, 0) \geq \log_2 |\mathcal{G}_s|$ .

Combining the two bounds, we get  $E_G(0, 0) = \log_2 |\mathcal{G}_s| = n_s - 1$ , and from here equation (13) for the GE for spins follows immediately.  $\square$

Importantly, the techniques used in theorem 1 can also be used to deal with blocks of spins whenever the blocks form a bipartite lattice. In general, one may first disentangle as many qubits as possible inside the blocks. Then the remaining qubits are left in an entangled state where the GE is equal to the number of remaining independent star operators, which amounts to a boundary law term for each block (possibly with a sub-leading correction depending on how the blocks are chosen) plus a topological term. As an example of this, let us present the following theorem:

**Theorem 2.** *Given a partition of the square lattice  $\Sigma$  into  $n_b$  blocks of boundary size  $L$  as indicated in figure 2(c) (where  $L$  is measured in number of qubits), then the four ground states  $|i, j\rangle$  for  $i, j = 0, 1$  and also the excited states  $|\phi, c, i, j\rangle$  of the toric code Hamiltonian have all the same GE of blocks and is given by*

$$E_G = \frac{n_b L}{4} - 1. \quad (14)$$

**Proof.** First, notice that lemmas 1 and 2 imply that we can entirely focus on the ground state  $|\widetilde{0}, 0\rangle$  of a toric code model in the deformed lattice  $\widetilde{\Sigma}$  from figure 2(e). As shown in figures 2(d)

and (e), it is always possible to remove *all* the stars inside of each block in  $\Sigma$  just by doing CNOTs locally inside of each block. As a result, the stars in the deformed lattice  $\tilde{\Sigma}$  correspond to those in  $\Sigma$  that lay among the blocks. It is easy to see that, for  $n_b$  blocks of boundary  $L$ , there are  $\tilde{n}_s = n_b L/4$  of such stars.

The rest of the proof follows as in theorem 1, by upper and lower bounding  $E_G(0, 0)$  for the partition with respect to blocks. For the upper bound, we use the fact that  $\Lambda_{\max}^{[\text{blocks}]} \geq \Lambda_{\max}^{[\text{spins}]} \geq |\tilde{\mathcal{G}}_s|^{-1/2}$ , where the first inequality again used the property that if larger blocks are considered then the maximum overlap is also larger, and  $\tilde{\mathcal{G}}_s$  is the corresponding group of contractible loop operators on  $\tilde{\Sigma}$ . From here  $E_G(0, 0) \leq \log_2 |\tilde{\mathcal{G}}_s|$  follows. For the lower bound, we divide the system into two sets  $A$  and  $B$  of blocks as indicated in figure 4(b). Following a similar reasoning as in theorem 1, it is possible to see again that no element  $g \in \tilde{\mathcal{G}}_s$  will act trivially on  $A$  or  $B$  except for the identity element. From this point, the rest of the proof is simply equivalent to the proof for theorem 1.  $\square$

**3.4.3. Interpretation of the previous theorems.** Let us now discuss equations (13) and (14) in detail. First, let us remind that a variety of works have shown the existence of a ‘boundary law’ for the entanglement entropy of many 2D systems (see e.g. [24]), including models with TO [6, 7]. It is then remarkable that, according to the first term of these equations, the entanglement per block obeys also a boundary law, i.e. it is proportional to the size  $L$  of the boundary of the block. To the best of our knowledge, these results are the first example of a boundary law behaviour for a multipartite (rather than bipartite) measure of entanglement in 2D.

However, the second term in equations (13) and (14) is far more intriguing and important. Its existence is caused by the global constraint from equation (10) on star operators which, in turn, allow for the topological degeneracy of the ground state. Thus, this term is *of topological nature, and quantifies the pattern of long-range entanglement* present in topologically ordered states.

In order to clarify further the meaning of the topological term, let us consider the bipartite case of a block of spins of boundary size  $L$  and its environment (i.e. the rest of spins). In the bipartite case, the geometric entanglement  $E_G$  coincides with the single-copy entanglement  $E_1 = -\log v_1(\rho)$ , with  $v_1(\rho)$  the largest eigenvalue of the reduced density matrix  $\rho$  of the block. Since this density matrix is proportional to a projector (see [6] and also [9]), we have that  $E_1 = S$ , with  $S$  the entanglement entropy of the bipartition. Thus, for this case we have the chain of equalities

$$E_G = E_1 = S. \quad (15)$$

We also know that for a system with TO, the entanglement entropy satisfies  $S = S_0 - S_\gamma + O(L^{-\nu})$ , where  $S_0$  is some boundary law term and  $S_\gamma$  is the topological entropy. As explained in [6], for the toric code model the global constraints in equation (10) imply that  $S_\gamma = 1$ . Thus, with the convention from equation (4), we have that the *topological GE* is given, in this case, by

$$E_\gamma = S_\gamma. \quad (16)$$

The topological origin of  $E_\gamma$  is thus clear. Notoriously, this topological contribution is maintained when promoting the entanglement measure from the bipartite to the multipartite scenario.

**3.4.4. Honeycomb lattice.** For the honeycomb lattice, most of the derivations are equivalent to the square lattice, yet there is a small difference: while the square lattice is self-dual, the honeycomb lattice is not. This means, among other things, that the toric code in the honeycomb lattice is not self-dual with respect to a change of star- and plaquette-operators (which is actually the case for the square lattice). In practice, this implies that different choices of ground state basis may have different values of the GE of spins. However, this is only a short-distance property, since the GE of blocks has the same topological contribution as for the square lattice. In order to illustrate this more explicitly, we provide calculations for two specific ground states.

*The  $|0, 0\rangle$  ground state.* We first consider the  $|0, 0\rangle$  ground state obtained as

$$|0, 0\rangle = \frac{1}{\sqrt{|\mathcal{G}_s|}} \sum_{g \in \mathcal{G}_s} g|0\rangle^{\otimes n}, \quad (17)$$

which is nothing but the  $i = 0, j = 0$  ground state in the notation of equation (12). This time,  $\mathcal{G}_s$  is the group generated by the product of all possible star operators in the *honeycomb* lattice. As before,  $|\mathcal{G}_s| = 2^{n_s-1}$ , where  $n_s$  is the number of stars (vertices) in the honeycomb lattice. For this state, we provide the following observation (not a theorem):

**Observation 1.** *For the toric code Hamiltonian in a honeycomb lattice  $\Sigma$ , the GE of spins for the ground state  $|0, 0\rangle$  is upper- and lower-bounded by*

$$n_p - 1 \leq E_G \leq n_p + 1, \quad (18)$$

where  $n_p$  is the number of plaquettes (faces) in  $\Sigma$ . Moreover, numerically it looks like the upper bound is saturated, so that

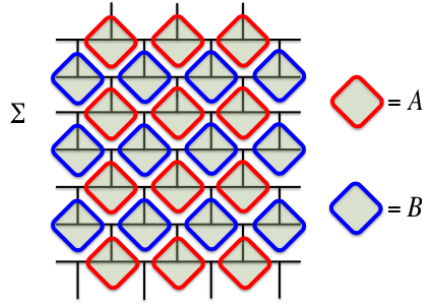
$$E_G = n_p + 1. \quad (19)$$

The above is an observation, and not a theorem, since equation (19) has only been checked numerically. Yet, let us prove now the upper and lower bounds in equation (18). First, the maximal overlap can be easily upper bounded again by  $\Lambda_{\max} \geq |\mathcal{G}_s|^{-1/2} = 2^{(1-n_s)/2}$ . Nevertheless, unlike in the case of the square lattice, this time we can actually have a better bound, i.e.

$$\begin{aligned} \Lambda_{\max} &\geq \langle + |^{\otimes n} |0, 0\rangle \\ &= \frac{1}{\sqrt{2^n}} \frac{1}{\sqrt{|\mathcal{G}_s|}} \sum_{x \in \{0,1\}^n} \sum_{g \in \mathcal{G}_s} \langle x | g |0\rangle^{\otimes n} \\ &= \sqrt{\frac{|\mathcal{G}_s|}{2^n}} = \sqrt{\frac{2^{2n_p-1}}{2^{3n_p}}} = \frac{1}{\sqrt{2^{n_p+1}}}, \end{aligned} \quad (20)$$

which is also smaller than  $|\mathcal{G}_s|^{-1/2} = 2^{(1-n_s)/2}$  using the fact that  $n = 3n_p$  and  $n_s = 2n_p$  for the honeycomb lattice. This gives the desired upper bound on the GE,  $E_G \leq n_p + 1$ .

To obtain the lower bound on GE (or equivalently upper bound on the maximal overlap), our strategy is again similar to the one for the square lattice. This time we consider the bipartition illustrated in figure 5. There is a star operator in each of the smallest units of the sets  $A$  (red) and  $B$  (blue). It is easy to convince oneself that the nontrivial subgroup  $\mathcal{G}_s(A)$  of  $\mathcal{G}_s$  is generated by the  $n_s/4$  star operators residing in all units in  $A$ , and similarly for  $\mathcal{G}_s(B)$ . The size of these



**Figure 5.** Bipartition of edges on the honeycomb lattice (or equivalently the brickwall lattice)  $\Sigma$ . We have two sets:  $A$  (red-grey) and  $B$  (blue-white). Edges crossed by the block boundaries are regarded as inside the block.

**Table 1.** GE computed numerically for finite honeycomb lattices. We use a  $N \times N$  brickwall structure to represent the honeycomb lattice with periodic boundary conditions, where  $n_s = N^2$ ,  $n = (3/2)N^2$  and  $n_p = n - n_s$  (with even  $N$ ). Our calculations are accurate with a  $10^{-6}$  of relative error.

$N \times N$	$n_s$	$n$	$n_p$	$E_G$
$2 \times 2$	4	6	2	3
$2 \times 4$	8	12	4	5
$2 \times 6$	12	18	6	7
$2 \times 8$	16	24	8	9
$4 \times 4$	16	24	8	9
$4 \times 6$	24	36	12	13

groups is therefore  $|\mathcal{G}_s(A)| = |\mathcal{G}_s(B)| = 2^{n_s/4}$ . The reduced density matrix of e.g. set  $A$  satisfies  $\rho_A^2 = (|\mathcal{G}_s(A)\rangle\langle\mathcal{G}_s(B)|/|\mathcal{G}_s|)\rho_A$ , and hence the maximal overlap  $\Lambda_{\max}^2$  is upper bounded by

$$\Lambda_{\max}^2 \leq \frac{|\mathcal{G}_s(A)\rangle\langle\mathcal{G}_s(B)|}{|\mathcal{G}_s|} = \frac{1}{2^{n_p-1}}, \quad (21)$$

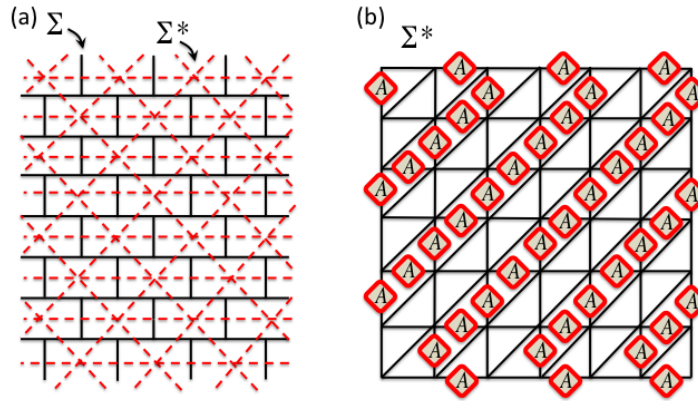
which leads to  $n_p - 1 \leq E_G$ , as claimed. Combining this bound with the one computed previously, we arrive at the result that we mentioned before, namely,  $n_p - 1 \leq E_G \leq n_p + 1$ .

In order to see whether any of these bounds is tight, we have carried out a numerical computation for small system sizes and confirmed that it is the upper bound which is saturated, i.e.  $E_G = n_p + 1$ . This can be seen in table 1. We would like to emphasize that even though we do not obtain  $-1$ , this is only a short-range property (where each block size is unity), as  $E_G$  for blocks will give rise to the correct contribution; see below.

*The  $|+, +\rangle$  ground state.* We now consider the  $|+, +\rangle$  ground state. By definition, this ground state is obtained by acting on the reference state  $|+\rangle^{\otimes n}$  with the elements of the group  $\mathcal{G}_p$  of plaquette operators, namely,

$$|+, +\rangle \equiv \frac{1}{\sqrt{|\mathcal{G}_p|}} \sum_{g \in \mathcal{G}_p} g|+\rangle^{\otimes n}. \quad (22)$$

Notice that, unlike in the square lattice (where there is a duality between stars and plaquettes), in the honeycomb lattice the  $|+, +\rangle$  ground state does not need to have, necessarily, the same



**Figure 6.** (a) Honeycomb lattice (or brickwall lattice)  $\Sigma$  in black lines, and its dual triangular lattice  $\Sigma^*$  in dotted red lines. (b) Bipartition of edges in  $A$  and  $B$  sets on the triangular lattice (deformed to fit in to a square lattice). For simplicity of the figure, only the spins in the  $A$  set are shown (in red). The rightmost edges are identified with the leftmost edges, and the bottom edges are identified with the top edges, since periodic boundary conditions are used.

properties as the  $|0, 0\rangle$  ground state from the previous section (at least not the same short-distance properties). For this state we can actually find its GE exactly, and provide the following theorem:

**Theorem 3.** *For the toric code Hamiltonian in a honeycomb lattice  $\Sigma$ , the GE of spins for the ground state  $|+, +\rangle$  is given by*

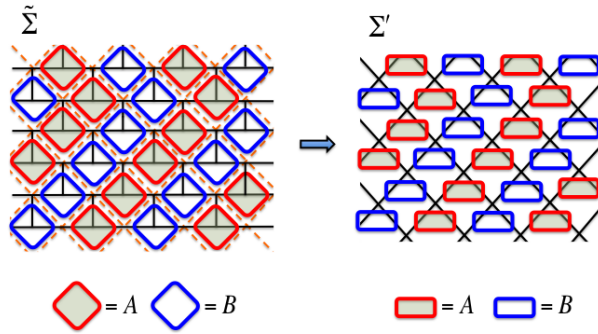
$$E_G = n_p - 1, \quad (23)$$

where  $n_p$  is the number of plaquettes (faces) in  $\Sigma$ .

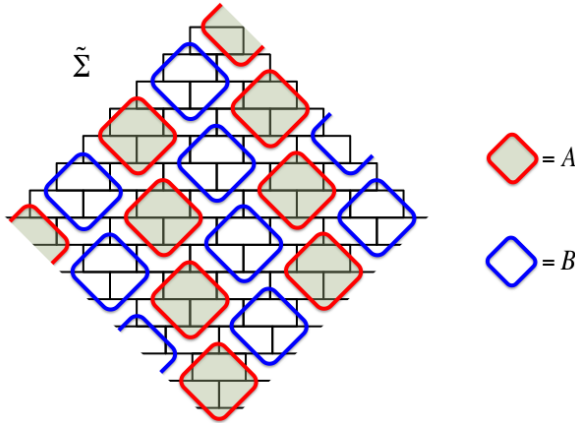
**Proof.** We can obtain the GE analytically as follows: first, we use the duality that the state  $|+, +\rangle$  in the honeycomb lattice  $\Sigma$  is equivalent to the  $|0, 0\rangle$  ground state of the toric code on the dual lattice  $\Sigma^*$ , which is the triangular lattice. After this, we simply follow the same approach as before by finding appropriate upper- and lower-bounds. The trivial upper bound is given by  $n_p - 1 \leq E_G$ , with  $n_p$  the number of plaquettes in the original honeycomb lattice. To find the lower bound, we consider the bipartition given in figure 6, which is similar to the one that we used square lattice (figure 2). Using the same methods as before, the expression for the GE follows.  $\square$

*Blocking the honeycomb lattice.* The fact that the  $|0, 0\rangle$  state on different lattices gives different constant contributions to the GE of spins turns out to be a short-distance property that disappears when blocks of spins are considered. Blocking corresponds to coarse-graining and gives long-wavelength properties.

This is easy to notice, if one realizes that in just one RG step we can map the honeycomb to the square lattice. Let us explain this in more detail. It is well known that the toric code in the honeycomb and square lattices can be mapped to one another by simple RG (disentangling) operations acting locally on the lattice. These transformations (from the square



**Figure 7.** Blocking and bipartition of a honeycomb lattice on a torus. The blocking is such that the CNOT operations from figure 3, which map the honeycomb lattice (black) to the square lattice (orange), fall within the blocks, so they do not affect the calculation of the GE. The GE of three-site blocks for the honeycomb lattice (left) then maps to the GE of two-site blocks for the square lattice (right). Larger block sizes could also be considered.



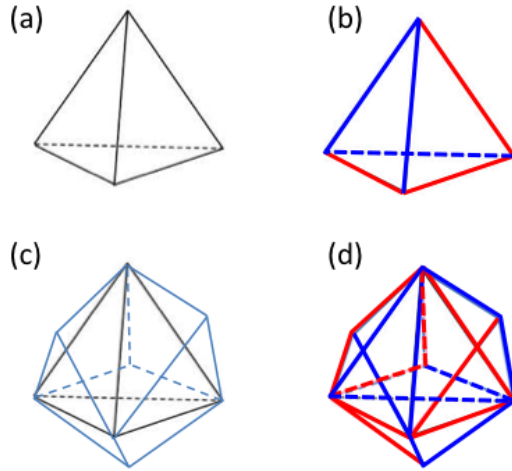
**Figure 8.** Blocking and bipartition of the honeycomb lattice on a torus, for larger blocks. Each block is defined by edges contained in red or blue square, where the colouring indicates the bipartition. Larger blocks can be considered similarly.

to the honeycomb lattice, and back to the square) are represented in the diagram of figure 3. In fact, these moves can also be understood in terms of entanglement renormalization steps [15], where some of the initial qubits become disentangled from the rest after every step.

Importantly, after mapping the honeycomb to the square lattice, the GE follows exactly the laws explained in the previous section, and hence it is evident that  $E_\gamma = 1$  also holds in this case after a suitable blocking of the spins (so that the RG CNOTs fall within the blocks).

For instance, one could choose the blocking shown in figure 7 and obtain  $E_\gamma = 1$ , but similar results hold for other choices of blocks as well. One could also consider larger blocking such as the one shown in figure 8. It is interesting to know that one-step of blocking is sufficient to remove the short-range behaviour and result in a correct topological contribution. Needless to say, in the process of blocking the corresponding ‘boundary laws’ for the bulk contribution of the GE of blocks also follow.





**Figure 9.** (a) A tetrahedron; (b) bipartition of edges in a tetrahedron (red/blue); (c) the same topology but with one tetrahedron grown onto each face; (d) bipartition (red/blue) of the edges in (c).

**3.4.5. Brief discussion of the honeycomb lattice results.** Let us comment briefly on the above results. Firstly, we have seen that the GE of spins is, in this case, different for the  $|0, 0\rangle$  and  $|+, +\rangle$  states. One could also infer from here that the topological contribution to the GE is also different. However, we wish to stress that the topological GE is a long-range contribution, and can only be extracted reliably after appropriate blocking of the spins. This is exactly what has been observed when considering the blocking of spins in the honeycomb lattice. Secondly, it is clear that the same conclusions that applied to excited states and other ground states on the square lattice apply also to the honeycomb lattice, i.e. once chosen a ground state basis, all the states in the basis as well as the corresponding (quasiparticle) excitations have the same amount of entanglement as the ‘reference’ ground state (which in this case is either  $|0, 0\rangle$  or  $|+, +\rangle$ ), as long as they are related by local operations.

**3.4.6. A different topology: the sphere.** In addition to the topology of a torus, we also consider that of a sphere. The simplest geometry is a tetrahedron, shown in figure 9(a). The toric code ground state on this topology is unique.

Using the same technique of upper and lower bounding the maximal overlap, one can easily show that its GE of spins (i.e. one-site blocks) is given by  $E_G = n_s - 1$ , where  $n_s = 4$  is the number of vertices and is related to the number of spins  $n$  via  $n = dn_s/2$ , where  $d = 3$  is the degree of a vertex. The bipartition that gives rise to this result is shown in figure 9(b).

In order to obtain a scaling we add one tetrahedron to each face, which can be iterated in order to add more sites to the system. The first step of such a procedure is shown in figure 9(c). In general it is easy to obtain  $E_G = n_s - 1 = 2n/d - 1$ , where  $n_s$ ,  $n$  and  $d$  are, respectively, the number of vertices, the number of spins (edges), and the vertex degree of the associated geometry. This result shows that the  $-1$  contribution is not only present in the torus topology, but in the sphere topology as well.

#### 4. Beyond the toric code

After our detailed study on the toric code model, we now turn to study the GE in other topological models that also correspond to fixed points of RG. From now on, some parts of our study will not be as detailed as for the toric code, and will focus mainly on some of the fundamental properties. Yet, in order to make everything as clear as possible, we need to see how some of the main properties of the toric code apply also to other models. In particular, there are three main points to consider:

1. *Excitations and other ground states.* As for the toric code, excited states and other ground states that can be obtained by local operations acting on some reference ground state will share the same entanglement properties (cf lemma 1). Therefore, in many cases we shall only discuss the entanglement of some ‘reference’ ground state.
2. *Disentangling general stabilizer states.* Very importantly, the disentangling property is not exclusive of the toric code only, but it holds for any stabilizer state as well. More concretely, *any state that is stabilized by a set of operators can be built from a quantum circuit acting on some separable initial state, where the gates of the quantum circuit involve the basis of the stabilizer group*<sup>8</sup>. If this basis of stabilizers can be thought of as acting locally on a lattice, then the quantum gates of the corresponding quantum circuit will also be local in the lattice. Moreover, there will be a set of unitary gates in the quantum circuit that will correspond to each one of the elements of the group basis, and which will take into account the existence of that particular basis element. That is: removing the action of the gates corresponding to a particular basis element effectively removes that stabilizer from the basis of the stabilizer group (e.g. removing a plaquette or star operator from the toric code).

Therefore, simply by *reversing* the action of such a quantum circuit on the stabilized state, we will always manage to *disentangle* the state by means of sufficiently local unitary operations. At every step, the remaining quantum state will be the state stabilized by the remaining set of stabilizer operators. In particular, this implies that there is always a MERA representation for any stabilizer state (not necessarily topological), and the toric code model is just a particular case.

This property turns out to be quite important, since it allows us to extend straightforwardly the results for the GE of blocks beyond the toric code: when blocking is considered, simply perform disentangling operations inside of each block in order to remove the maximum of stabilizers within each block, so that the remaining stabilizers will always be between spins in different blocks. By construction, the remaining state will usually obey an area law for the GE of blocks, plus a possible topological term. Because of this, from now on the GE of blocks will not be discussed in detail except for a couple of examples.

3. *The non-Abelian case.* It is indeed possible to generalize our results to systems with non-Abelian symmetries, such as quantum double models. However, in order to deal with these models, it will be worthwhile to develop part of our formalism in the most general fashion. This will be done in section 7, and quantum double models will be considered subsequently in section 8.

<sup>8</sup> The discussion around equation (7.160) in chapter 7 of the lecture notes of J Preskill, [www.theory.caltech.edu/people/preskill/ph229/](http://www.theory.caltech.edu/people/preskill/ph229/)

## 5. Double semion model

The double semion model is given by the spin model on the honeycomb lattice

$$H_{\text{DS}} = - \sum_s A_s - \sum_p B'_p, \quad (24)$$

where  $A_s$  and  $B'_p$  are mutually commuting and given by

$$A_s \equiv \prod_{j \in s} \sigma_x^{[j]}, \quad B'_p \equiv - \prod_{k \in \text{legs of } p} i^{(1-\sigma_x^{[k]})/2} \prod_{j \in p} \sigma_z^{[j]}. \quad (25)$$

This is, the star operators are the same as for the toric code, whereas the plaquette operators are related via

$$B'_p = -B_p \prod_{k \in \text{legs of } p} i^{(1-\sigma_x^{[k]})/2} \quad (26)$$

with  $B_p$  the toric code plaquette operator. As for the toric code, star and plaquette operators satisfy the non-local constraint

$$\prod_s A_s = \prod_p B'_p = \mathbb{I}. \quad (27)$$

This model is known to be also topologically ordered, yet its topological phase differs from the one of the toric code model. More precisely, while the toric code corresponds to the topological phase of a  $\mathbb{Z}_2$  gauge theory, the double semion model corresponds to a  $U(1) \times U(1)$  Chern–Simons theory [5].

One of the ground states of this double semion model is almost identical to the toric code ground state  $|+, +\rangle$ , except that each term in the decomposition equation (22) is multiplied by a factor  $(-1)^{X_c}$ , where  $X_c$  counts the number of loops formed by spins in the  $|-\rangle$  state. More specifically, this ground state is given by

$$|+, +\rangle_{\text{DS}} \equiv \frac{1}{\sqrt{|\mathcal{G}_p|}} \sum_{g \in \mathcal{G}_p} (-1)^{X_c} g |+\rangle^{\otimes n}, \quad (28)$$

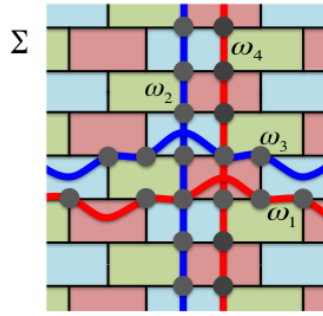
where again  $\mathcal{G}_p$  is the group of all possible products of plaquette operators  $B'_p$ . Of course, the state  $|+, +\rangle$  is by construction a stabilized state of the group  $\mathcal{G}_p$ . We provide now the following theorem:

**Theorem 4.** *For the double semion model in a honeycomb lattice  $\Sigma$ , the GE of spins for the ground state  $|+, +\rangle_{\text{DS}}$  is given by*

$$E_G = n_p - 1, \quad (29)$$

where  $n_p$  is the number of plaquettes (faces) in  $\Sigma$ .

**Proof.** Remember the ground state  $|+, +\rangle$  for the toric code in the honeycomb lattice. This state can always be written as  $|+, +\rangle = \sum c_{ijk\dots} |ijk\dots\rangle$  with  $c_{ijk\dots} \geq 0$  in the local basis  $\{|+\rangle, |-\rangle\}$  for every spin. From here, it follows that  $|\sum_{ijk\dots} p_i q_j \dots c_{ijk\dots}| \leq \sum_{ijk\dots} |p_i q_j \dots c_{ijk\dots}|$ , where the equality holds if  $p_i, q_j, \dots$  are all non-negative up to a common phase, and the maximum is achieved by choosing  $p_i, q_j, \dots$  to be non-negative. Next, let us define  $d_{ijk\dots} = c_{ijk\dots} e^{i\phi}$ , where  $\phi$  is an arbitrary phase. Then it follows that  $|\sum_{ijk\dots} p_i q_j \dots d_{ijk\dots}| \leq \sum_{ijk\dots} |p_i q_j \dots c_{ijk\dots}|$ .



**Figure 10.** Colour code on a honeycomb (brickwall) lattice  $\Sigma$ , which is an example of a 2-colex. There is a spin on every vertex, and every plaquette is coloured with either red, green or blue. Non-local string operators  $w_1, w_2, w_3$  and  $w_4$  are also shown, where periodic boundary conditions on a torus are assumed. Grey dots indicate the sites where string operators are acting.

A maximization over the left hand side gives the maximal overlap, say, for the double semion model by choosing appropriate phase factors and coefficients. But in any case, this will be smaller than maximizing the right hand side whose maximum is, in this particular case, the maximal overlap for the toric code. From here the result just follows.  $\square$

From the above theorem it follows that the topological contribution is  $E_\gamma = 1$ , which is again equal to the topological entropy for this model. Moreover, even if we do not discuss it here in detail, we expect the GE of blocks to obey similar laws as for the toric code model, since the ground state can also be disentangled by local moves [25]. Similarly, the entanglement properties of other ground states and excited states obtained by acting locally on  $|+, +\rangle$ , remain identical to those of  $|+, +\rangle$  (cf lemma 1).

## 6. Colour code models

Topological colour codes [12] are defined as follows: qubits reside on *vertices* (not links!) of a lattice that is a 2-colex. A 2-colex  $\Sigma$  is a 2D lattice embedded in a torus of arbitrary genus  $g$ , with the following properties: (a) every vertex of the lattice has degree 3, i.e. precisely three edges meet at each vertex; (b) the faces of the lattice are 3-colourable, i.e. the faces can be coloured with 3-colours (e.g. red, green and blue) in such a way that no two adjacent faces have the same colour, see e.g. figure 10. Two types of commuting operators are defined on each plaquette (or face)  $p$ ,

$$B_p^X \equiv \prod_{j \in p} \sigma_x^{[j]}, \quad B_p^Z \equiv \prod_{j \in p} \sigma_z^{[j]}. \quad (30)$$

The Hamiltonian of the colour code is thus

$$H_{CC} = - \sum_p (B_p^X + B_p^Z). \quad (31)$$

For this model in 2D lattices, the ground space has degeneracy  $2^{4g}$ , where  $g$  is the genus of the underlying Riemann surface. For instance, on the torus this degeneracy is 16.

Let us define the group  $\mathcal{G}_p^X$  as the group generated by the products of all possible plaquette operators  $B_p^X$  for all three colours. Notice that the product of all plaquette operators  $B_p^X$

corresponding to the same colour  $c$  corresponds to the same element of the group  $\mathcal{G}_p^X$  (namely, the action of a  $\sigma_x$  operator on all the qubits). Therefore, we have the non-local constraint

$$\prod_{p \in \text{red}} B_p^X = \prod_{p \in \text{green}} B_p^X = \prod_{p \in \text{blue}} B_p^X = \prod_{v \in \Sigma} \sigma_x^{[v]}. \quad (32)$$

The above relation implies that the size of the group is  $|\mathcal{G}_p^X| = 2^{n_p-2}$ , with  $n_p$  the number of plaquettes in lattice  $\Sigma$ .

The ground level subspace of the colour code model is constructed in a very similar way as we did for the toric code. Specifically, for the colour code on a torus the ground space reads  $\mathcal{L} = \text{span}\{|i, j, k, l\rangle, i, j, k, l = 0, 1\}$ , where

$$|i, j, k, l\rangle = \frac{1}{\sqrt{|\mathcal{G}_p^X|}} \sum_{g \in \mathcal{G}_p^X} g w_1^i w_2^j w_3^k w_4^l |0\rangle^{\otimes n}. \quad (33)$$

Operators  $\{w_1, w_2, w_3, w_4\}$  above correspond to non-contractible loop operators, see figure 10(b). We do not enter here on the specific action of these operators on every spin in the loop, and just mention that it is local. The 16 vectors  $|i, j, k, l\rangle$  are orthonormal and stabilized by  $\mathcal{G}_p^X$ , and therefore form a possible basis of the ground level space of the model on a torus.

As for the toric code, the action of loop operators corresponds to local unitary operations on a reference state  $|0, 0, 0, 0\rangle$ , and therefore do not change the entanglement properties (cf lemma 1). Therefore, from now on we shall consider only the entanglement properties of the ground state  $|0, 0, 0, 0\rangle$ , defined as

$$|0, 0, 0, 0\rangle = \frac{1}{\sqrt{|\mathcal{G}_p^X|}} \sum_{g \in \mathcal{G}_p^X} g |0\rangle^{\otimes n}. \quad (34)$$

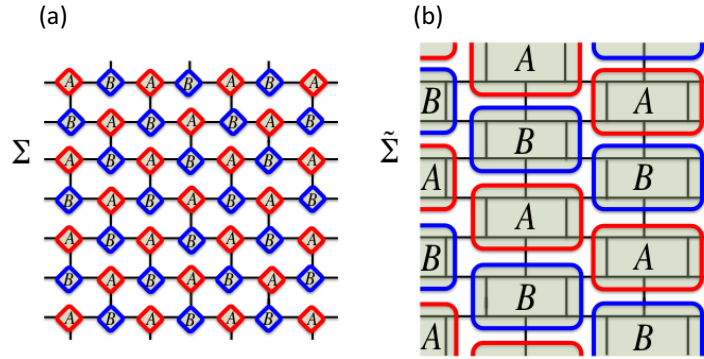
In what follows we show two different approaches to compute the GE for this model. The first approach is based, as all the calculations before, on finding upper and lower bounds for the GE. This approach works well for e.g. finding the GE of spin and blocks on regular lattices such as the honeycomb lattice. The second approach uses a different technique, building on recent work [26]. The main idea is to exploit that colour code ground states are instances of so-called Calderbank–Shor–Steane states of self-orthogonal type. This refers to the fact that  $|0, 0, 0, 0\rangle$ , when expanded in the standard basis, has the form

$$|0, 0, 0, 0\rangle = \frac{1}{\sqrt{|S|}} \sum_{x \in S} |x\rangle, \quad (35)$$

where  $S \subseteq \mathbb{Z}_2^n$  is a collection of bit strings with a particular structure i.e. it is a self-orthogonal classical linear code. This property is in fact displayed by colour code states associated with arbitrary lattices (2-colexes) and for surfaces of arbitrary genus. As a result, the second approach will lead to a general calculation of the GE for arbitrary colour codes.

### 6.1. Bound method

Let us consider the colour code on a honeycomb lattice  $\Sigma$ . For this lattice, and with this method, we can provide the following theorem:



**Figure 11.** (a) Natural bipartition of the honeycomb lattice  $\Sigma$  into  $A$  and  $B$  sets (respectively red and blue). (b) Possible blocking and bipartition of lattice  $\tilde{\Sigma}$ , obtained after applying disentglers corresponding to the stabilizers within each block to the honeycomb lattice  $\Sigma$  of the colour code. We also show the relevant bipartition into  $A$  and  $B$  sets (again red and blue).

**Theorem 5.** For the colour code model in a honeycomb lattice  $\Sigma$ , the GE of spins for the ground state  $|0, 0, 0, 0\rangle$  is given by

$$E_G = n_p - 2, \quad (36)$$

where  $n_p$  is the number of plaquettes (faces) in  $\Sigma$ .

**Proof.** Firstly, since the group  $\mathcal{G}_p^X$  has size  $|\mathcal{G}_p^X| = 2^{n_p-2}$ , we obtain a lower bound on the maximal overlap  $\Lambda_{\max} \geq |\mathcal{G}_p^X|^{-1/2}$ , and hence an upper bound on the GE given by  $E_G \leq n_p - 2$ . Secondly, in order to derive a good lower bound on the GE, we consider the natural bipartition for the honeycomb lattice, see figure 11(a). One can convince oneself easily that any group element in  $\mathcal{G}_p^X$  must have non-trivial support contained in both the  $A$  and  $B$  sets, and therefore the subgroups  $\mathcal{G}_p^X(A)$  and  $\mathcal{G}_p^X(B)$  are both trivial and contain only the identity element. For the colour code, the reduced density matrix of  $A$  also satisfies the property  $\rho_A^2 = (|\mathcal{G}_p^X(A)\rangle\langle\mathcal{G}_p^X(B)|/|\mathcal{G}_p^X|)\rho_A$  [27]. Therefore,  $\Lambda_{\max}^2 \leq 1/|\mathcal{G}_p^X|$ , and hence  $E_G \geq n_p - 2$ . We therefore arrive at  $E_G = n_p - 2$ , which proves the theorem.  $\square$

Some comments are in order. Firstly, we see that the term of topological origin is  $E_\gamma = 2$ , which is once again identical to the topological entanglement entropy [27]. Secondly, this topological GE is twice that of the toric code, which is a direct consequence that the colour code model we consider here is equivalent to two copies of the toric code under some given RG procedure [13]. Thirdly, a similar result holds also for the ground state associate with the group  $\mathcal{G}_p^Z$  generated by the product of  $B_p^Z$  operators. Fourthly, excitations and other ground states related by local operations will have the same entanglement. And fifthly, the GE of blocks will consist again of an area-law term plus a topological contribution which will be again  $E_\gamma = 2$ . This can be seen, e.g. by choosing the blocking from figure 11(b), and appropriately disentangling qubits within each block.

## 6.2. Colour codes and self-orthogonal classical codes

Here we provide a generalization of theorem 5 which will allow to compute the GE of colour code states in a much broader setting i.e. for arbitrary lattices (2-colexes) and arbitrary genus. The technique used to arrive at this generalization is entirely different from the proof of theorem 5. It involves connections between colour codes and classical coding theory which were found in the recent work [26].

We consider a colour code defined on an arbitrary 2-colex embedded in a surface of arbitrary genus. In analogy with (34) we consider the state

$$|\mathbf{0}\rangle = \frac{1}{\sqrt{|\mathcal{G}_p^X|}} \sum_{g \in \mathcal{G}_p^X} g |0\rangle^{\otimes n}. \quad (37)$$

We will show here that the result in theorem 5 extends to all states  $|\mathbf{0}\rangle$ :

**Theorem 6.** *Consider the colour code ground state  $|\mathbf{0}\rangle$  for an arbitrary 2-colex  $\Sigma$  embedded in a surface of arbitrary genus. The GE of spins for this state is*

$$E_G = n_p - 2. \quad (38)$$

Before proving the result, we develop some preliminary material. A linear subspace  $C \subseteq \mathbb{Z}_2^n$  is called a (classical) binary linear code of length  $n$ . The elements of  $C$  are called its codewords. For every binary linear code  $C$  of length  $n$ , define an  $n$ -qubit state

$$|C\rangle := \frac{1}{\sqrt{|C|}} \sum_{u \in C} |u\rangle. \quad (39)$$

We shall refer to any state of this kind as a CSS state. It easily follows from (37) that  $|\mathbf{0}\rangle$  is a CSS state. More precisely, we have  $|\mathbf{0}\rangle \equiv |C_p^X\rangle$  where the linear code  $C_p^X$  is defined as follows: writing every element  $g \in \mathcal{G}_p^X$  as

$$g = X^{u_1} \otimes \dots \otimes X^{u_n} \quad (40)$$

with  $u_i \in \mathbb{Z}_2$ , the code  $C_p^X$  is simply given by the collection of all bit strings  $(u_1, \dots, u_n)$  arising from elements of  $\mathcal{G}_p^X$  in this way. In particular, one has

$$|C_p^X| = |\mathcal{G}_p^X| = 2^{n_p-2}, \quad (41)$$

which will be important below.

It turns out that colour code states are CSS states of a particular kind. A linear code  $C$  is called self-orthogonal if every two codewords are orthogonal:  $s^T t = 0$  for every  $s, t \in C$ . A CSS state for which the underlying classical code is self-orthogonal will be denoted *CSS state of self-orthogonal type*. In [26] the following was shown:

**Lemma 3.** *Consider a colour code associated with an arbitrary 2-colex  $\Sigma$ . Then the ground state  $|\mathbf{0}\rangle = |C_p^X\rangle$  is a CSS state of self-orthogonal type.*

Next we provide a calculation of the GE of arbitrary CSS states of self-orthogonal type—by lemma 3 this will yield the GE of colour code ground states. Note however that not all CSS states of self-orthogonal type are colour code ground states i.e. our results extend beyond the colour code setting. It is also interesting to point out that toric code ground states, even though they are CSS states, are not of self-orthogonal type.

Consider an  $n$ -qubit system and let  $|\Phi\rangle = \bigotimes |\phi^{[i]}\rangle$  be an  $n$ -qubit product state, where  $|\phi^{[i]}\rangle = a_i|0\rangle + b_i|1\rangle$ . Define an associated  $n$ -qubit tensor product operator as follows:

$$A := \bigotimes_i A_i, \quad \text{where} \quad A_i := \begin{bmatrix} a_i & b_i \\ b_i & -a_i \end{bmatrix}. \quad (42)$$

Now consider an arbitrary self-orthogonal linear code  $C$  of length  $n$  and the associated  $n$ -qubit CSS state  $|C\rangle$ . The following lemma, proved in [26], relates the overlap  $\langle C|\Phi\rangle$  to the expectation value  $\langle C|A|C\rangle$ :

**Lemma 4.** *Let  $C$  be a binary self-orthogonal linear code of length  $n$  and let  $|\Phi\rangle$  be an  $n$ -qubit product state. Then*

$$\sqrt{|C|} \langle C|\Phi\rangle = \langle C|A|C\rangle. \quad (43)$$

Lemma 4 will now be used to prove:

**Theorem 7.** *Let  $|C\rangle$  be an arbitrary  $n$ -qubit CSS state of self-orthogonal type. Then the GE of spins is*

$$E_G(|C\rangle) = \log |C|. \quad (44)$$

**Proof.** As before, let  $\Lambda_{\max}$  denote the maximal value of  $|\langle C|\Phi\rangle|$  when  $|\Phi\rangle$  ranges over all product states. Firstly, taking an arbitrary  $u \in C$  yields  $|\langle C|u\rangle| = |C|^{-\frac{1}{2}}$  so that we obtain the lower bound  $\Lambda_{\max} \geq |C|^{-\frac{1}{2}}$ . Secondly, since  $|C\rangle$  has real and non-negative amplitudes in the standard basis, the maximal value of  $|\langle C|\Phi\rangle|$  will be reached by a product state which has real, non-negative amplitudes as well. Consider an arbitrary product state  $|\Phi\rangle$  with real and non-negative amplitudes and the associated tensor product operator  $A$  as in (42). Using lemma 4 we find

$$\begin{aligned} |\langle C|\Phi\rangle| &= \frac{1}{\sqrt{|C|}} |\langle C|A|C\rangle| \leq \frac{1}{\sqrt{|C|}} \|A\| \\ &= \frac{1}{\sqrt{|C|}} \prod_i \|A_i\|, \end{aligned} \quad (45)$$

where  $\|\cdot\|$  denotes the operator norm. Since each  $|\phi^{[i]}\rangle$  is a unit vector with real and non-negative amplitudes, each  $A_i$  is a real orthogonal matrix. Hence  $\|A_i\| = 1$ . This proves the upper bound  $\Lambda_{\max} \leq |C|^{-\frac{1}{2}}$  which matches the previously found lower bound. The result now follows immediately.  $\square$

Combining lemma 3, theorem 7 and equation (41) immediately proves theorem 6.

## 7. General formalism for the bound method

In order to extend our results to the non-Abelian setting (e.g. quantum double models), it is worthwhile to generalize the bounding technique that we have used so far. Looking back, we have calculated the GE of quantum states with a certain (Abelian) group symmetry, and our bounding strategy involved two steps. Firstly, we chose suitable product states to lower bound the maximal overlap (and hence to upper bound the GE). Secondly, we picked certain



bipartitions and determined the maximal overlap with bipartite product states in order to upper bound the maximal overlap with multipartite product states (and hence to lower bound the GE). As we have seen, this procedure gives *tight* bounds in almost all the cases considered.

We shall now distill the essence of this bounding strategy and derive general bounds on the GE of quantum states with a non-Abelian group symmetry. Other than placing all of our previous examples under a common roof these bounds may even be of independent interest beyond the present study of TO.

### 7.1. Overview

Let  $G$  be a finite group. Suppose its  $n$ -fold direct product  $\mathcal{G} = G^{\times n}$  acts on some  $m$ -partite Hilbert space  $\mathcal{H} = \mathcal{H}_1 \otimes \dots \otimes \mathcal{H}_m$  via product operators  $O^t = O_1^t \otimes \dots \otimes O_m^t$ . Consider states of the form

$$\sum_{t \in \mathcal{G}} O^t |\Phi\rangle, \quad (46)$$

where  $|\Phi\rangle = |\phi^{[1]}\rangle \otimes \dots \otimes |\phi^{[m]}\rangle$  is some reference product state such that any two local states  $O_k^t |\phi^{[k]}\rangle$  and  $O_k^u |\phi^{[k]}\rangle$  are either identical or orthogonal.

We can define the *global stabilizer*  $\mathcal{G}_\Phi \subset \mathcal{G}$  as the subgroup which leaves the reference state  $|\Phi\rangle$  invariant. Note that  $\mathcal{G}_\Phi$  need *not* be normal in general, so the collection  $\mathcal{G}' := \mathcal{G}/\mathcal{G}_\Phi$  of cosets is not necessarily a group. It is clear that the state

$$|\Psi\rangle = \frac{1}{\sqrt{|\mathcal{G}'|}} \sum_{t \in \mathcal{G}'} O^t |\Phi\rangle \quad (47)$$

is correctly normalized. Picking any component  $O^t |\Phi\rangle$  we get a trivial lower bound on the maximal overlap  $\Lambda_{\max}(\Psi)$  and hence by (6) an upper bound  $E_G(\Psi) \leq n \log_2 |G| - \log_2 |\mathcal{G}_\Phi|$ .

Now let  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$  a bipartition of the Hilbert space into subsystems  $A$  and  $B$  with the reference state  $|\Phi\rangle = |\phi_A\rangle \otimes |\phi_B\rangle$  partitioned accordingly. For each operator  $O^t$  we define a truncated operator  $O_X^t$  as the restriction of  $O^t$  to  $\mathcal{H}_X$  and the identity in  $\mathcal{H}_{\bar{X}}$ . We furthermore define the *local stabilizer*  $\mathcal{G}_\Phi(X) \subset \mathcal{G}$  for subsystem  $X$  as the subgroup whose truncated operators  $O_X^t$  leave the reference state  $|\Phi\rangle$  invariant. It is clear that the global stabilizer is contained in any local stabilizer, so  $\mathcal{G}'_\Phi(X) := \mathcal{G}_\Phi(X)/\mathcal{G}_\Phi$  makes sense. Note that again  $\mathcal{G}'_\Phi(X)$  is not necessarily a group. Nevertheless, for bipartitions with the property

$$\mathcal{G}_\Phi(A) \supseteq \mathcal{G}_\Phi(B) \quad (48)$$

we will find (see lemma 6 below) that the largest eigenvalue of the reduced density operator  $\rho_A$  equals  $|\mathcal{G}_\Phi(A)||\mathcal{G}_\Phi(B)|/(|\mathcal{G}'||\mathcal{G}_\Phi|)$  which by (7) immediately implies a lower bound on the GE. We can summarize both bounds on the GE in the following

**Theorem 8. (General bound).** *Let  $|\Psi\rangle$  be a state of the form (46). If there is a bipartition  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$  obeying (48), the GE of  $|\Psi\rangle$  is bounded by*

$$-\log_2 \frac{|\mathcal{G}_\Phi(A)||\mathcal{G}_\Phi(B)|}{|\mathcal{G}'_\Phi|} \leq E_G(\Psi) - n \log_2 |G| \leq -\log_2 |\mathcal{G}_\Phi|. \quad (49)$$

As is evident the derivation of this general bound is *independent* of the reference state  $|\Phi\rangle$  or any physical model in which the state  $|\Psi\rangle$  might arise, including any geometry or topology that might be involved. So the problem of obtaining actual bounds for specific states in specific

models with specific geometries (topologies) reduces to the much simpler problem of analysing the interplay between the group  $\mathcal{G}$ , the global stabilizer  $\mathcal{G}_\Phi$  and the local stabilizers  $\mathcal{G}_\Phi(X)$  in the particular case at hand. Additionally, specific models will often allow us to characterize these stabilizers in purely geometric (combinatorial) terms, which provides a constructive way to find bipartitions obeying (48). Remarkably, this strategy even yields *exact* values for the GE in many interesting cases as we will show (and have seen already). Namely, with a mild assumption in addition to (48) we immediately get

**Theorem 9.** *Let  $|\Psi\rangle$  a state of the form (46). If there is a bipartition  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$  such that*

$$\mathcal{G}_\Phi(A) = \mathcal{G}_\Phi(B) = \mathcal{G}_\Phi, \quad (50)$$

*the GE of  $|\Psi\rangle$  is given by*

$$E_G(\Psi) = n \log_2 |G| - \log_2 |\mathcal{G}_\Phi|. \quad (51)$$

We would like to mention that if  $G$  is Abelian then both  $\mathcal{G}$  and  $\mathcal{G}'_\Phi(X)$  are indeed groups and one can determine the spectrum of  $\rho_A$  directly in terms of these. This has been discussed in [6] for the special case  $G = \mathbb{Z}_2$ . There our  $\mathcal{G}$  is denoted by  $G$  and our  $\mathcal{G}'_\Phi(A)$  by  $G_A$  (and similarly for the other subsystem  $B$ ).

## 7.2. Some details

Here we will supply some additional details glossed over before. First the precise definitions of global and local stabilizers.

**Definition 1** (Global stabilizer).

$$\mathcal{G}_\Phi := \{t \in \mathcal{G} \mid O^t |\Phi\rangle = |\Phi\rangle\}. \quad (52)$$

**Definition 2** (Local stabilizer).

$$\mathcal{G}_\Phi(X) := \{t \in \mathcal{G} \mid O_X^t |\Phi\rangle = |\Phi\rangle\}. \quad (53)$$

This shows how we obtain arbitrary reduced density operators from the state (47).

**Lemma 5** (Reduced density operator). *For any bipartition  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$  into subsystems  $A$  and  $B$  one has*

$$\rho_A = \frac{1}{|\mathcal{G}||\mathcal{G}_\Phi|} \sum_{u \in \mathcal{G}} \sum_{v \in \mathcal{G}_\Phi(B)} O_A^u |\phi_A\rangle \langle \phi_A| O_A^{vu^{-1}}. \quad (54)$$

**Proof.** Avoiding the sum over cosets, we may rewrite (47) as

$$|\Psi\rangle = \frac{1}{\sqrt{|\mathcal{G}||\mathcal{G}_\Phi|}} \sum_{t \in \mathcal{G}} O_A^t |\phi_A\rangle \otimes O_B^t |\phi_B\rangle,$$

hence the global density operator reads

$$\rho = \frac{1}{|\mathcal{G}||\mathcal{G}_\Phi|} \sum_{u,v \in \mathcal{G}} O_A^u |\phi_A\rangle \langle \phi_A| O_A^{v^{-1}} \otimes O_B^u |\phi_B\rangle \langle \phi_B| O_B^{v^{-1}}.$$

By assumption  $\{O_B^t|\phi_B\rangle \mid t \in \mathcal{G}/\mathcal{G}_\Phi(B)\}$  is a set of mutually orthogonal states in subsystem  $B$  so we can use it to take the partial trace

$$\mathrm{tr}_B(\rho) = \frac{1}{|\mathcal{G}||\mathcal{G}_\Phi|} \sum_{u,v \in \mathcal{G}} O_A^u |\phi_A\rangle \langle \phi_A| O_A^{vu^{-1}} \cdot \langle \phi_B| O_B^v | \phi_B\rangle.$$

Since the local expectation value is non-zero precisely for  $v \in \mathcal{G}_\Phi(B)$  we obtain

$$\mathrm{tr}_B(\rho) = \frac{1}{|\mathcal{G}||\mathcal{G}_\Phi|} \sum_{u \in \mathcal{G}} \sum_{v \in \mathcal{G}_\Phi(B)} O_A^u |\phi_A\rangle \langle \phi_A| O_A^{vu^{-1}}.$$

□

We turn to those bipartitions obeying (48).

**Lemma 6** (Spectrum of reduced density operators). *Let  $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$  a bipartition such that  $\mathcal{G}_\Phi(A) \supseteq \mathcal{G}_\Phi(B)$ . Then the spectrum of  $\rho_A$  is flat with non-zero eigenvalues*

$$\frac{|\mathcal{G}_\Phi(A)||\mathcal{G}_\Phi(B)|}{|\mathcal{G}||\mathcal{G}_\Phi|}. \tag{55}$$

**Proof.** From lemma 5 we conclude that

$$\rho_A = \frac{|\mathcal{G}_\Phi(B)|}{|\mathcal{G}||\mathcal{G}_\Phi|} \sum_{t \in \mathcal{G}} O_A^t |\phi_A\rangle \langle \phi_A| O_A^{t^{-1}},$$

and now the observation

$$\rho_A^2 = \frac{|\mathcal{G}_\Phi(A)||\mathcal{G}_\Phi(B)|}{|\mathcal{G}||\mathcal{G}_\Phi|} \rho_A$$

is enough to prove the claim. □

## 8. Quantum double models

Consider a planar graph  $\Gamma = (V, E, F)$  with vertices  $V$ , edges  $E$  and faces  $F$ , and sizes  $|V| = n_s$ ,  $|E| = n$  and  $|F| = n_p$ . The Hilbert space of the quantum double model is defined as  $\mathcal{H} := \mathcal{H}_1 \otimes \dots \otimes \mathcal{H}_{|E|}$  where each local  $\mathcal{H}_i \simeq \mathbb{C}G$ . Vertex projectors (acting on stars) read

$$A_s = \frac{1}{|G|} \sum_{g \in G} A_s^g,$$

where each individual vertex operator  $A_s^g$  acts on a vertex (star)  $s \in V$  by multiplying all edges meeting at  $s$  by  $g \in G$  from the *left*, provided all edges point towards  $s$ . Otherwise an edge is multiplied by  $g^{-1}$  from the *right*. It is clear that vertex operators on different vertices commute. Similarly, plaquette projectors are defined as

$$B_p = \delta\left(\prod_{i \text{ along } \mathcal{C}_p} g_i, e\right),$$

which select configurations where the ordered product of group elements along an oriented circuit  $\mathcal{C}_p$  around  $p$  is the unit element of  $G$ . The Hamiltonian is the sum of these mutually

commuting projectors over vertices and plaquettes, namely

$$H_{D(G)} = - \sum_{s \in V} A_s - \sum_{p \in F} B_p.$$

We identify the direct product group from our general discussion in the previous section as  $\mathcal{G} := G^{\times|V|}$  and obtain its natural action on the Hilbert space by collecting individual vertex operators into the *joint* vertex operator

$$O^t := \prod_{s \in V} A_s^{g_s},$$

where  $t = (g_1, \dots, g_{|V|}) \in \mathcal{G}$ . Clearly, any joint vertex operator has product form: given a directed edge  $(s, s')$  with value  $|x\rangle$  an element  $t = (\dots, g_s, \dots, g_{s'}, \dots) \in \mathcal{G}$  is easily seen to act locally as

$$O_{(s,s')}^t |x\rangle = |g_{s'}\rangle x g_s^{-1}, \quad (56)$$

thus  $O^t = O_1^t \otimes \dots \otimes O_{|E|}^t$ .

We can obtain a ground state of the quantum double model by choosing the reference product state  $|e\rangle := \bigotimes_E |e\rangle$  and projecting it on the common +1 eigenspace of all vertex operators

$$|\Omega\rangle = \prod_{s \in V} A_s |e\rangle = \frac{1}{|\mathcal{G}|} \sum_{t \in \mathcal{G}} O^t |e\rangle. \quad (57)$$

In order to proceed we need to figure out the actual form of the global and local stabilizers in the quantum double model for our particular reference state. As for elements of the global stabilizer  $\mathcal{G}_e$ , it follows immediately from (56) that  $g_{s'} = g_s$  iff vertices  $s$  and  $s'$  are connected by an edge. In particular, we have  $\mathcal{G}_e \simeq G$  if the graph  $\Gamma$  is connected. Note that  $\mathcal{G}_e$  is *not* normal in  $\mathcal{G}$  generally. As far as the local stabilizers  $\mathcal{G}_e(X)$  are concerned, we can equivalently define them combinatorially. Partition the vertices  $V$  into clusters  $V_i(X)$  based on whether vertices are connected by an edge in  $E_X$ . (That is, two vertices are in the same cluster iff they can be connected by a path which lies completely in  $E_X$ .) Then  $t = (g_1, \dots, g_{|V|}) \in \mathcal{G}_e(X)$  iff its components  $g_s$  are constant on clusters. If  $E_X$  connects *all* vertices into a single cluster then  $\mathcal{G}_e(X) \simeq G$ . It turns out that local and global stabilizers are related quite favourably:

**Lemma 7** (Joint local stabilizer).  $\mathcal{G}_e(A) \cap \mathcal{G}_e(B) = \mathcal{G}_e$ .

**Proof.** We want to cover the vertices  $V$  with a set  $\{V_i(X)\}$  of *overlapping* clusters obtained from any of the subsystems  $A$  or  $B$ . Then if  $t = (g_1, \dots, g_{|V|}) \in \mathcal{G}_e(A) \cap \mathcal{G}_e(B)$  its components  $g_s$  must be constant on each cluster, and hence constant on the whole of  $V$  because of the overlaps. Clearly, we can obtain a global cover by constructing a local cover for the vertices along each (edge) path  $\gamma$ .

So let  $\gamma$  be a path starting at vertex  $s$  and let  $V_1(A)$  be the enveloping cluster of  $s$  as given by subsystem  $A$ . If  $\gamma$  never leaves this cluster we are done. Otherwise we may assume that  $s$  lies at the boundary of  $V_1(A)$ , hence we will reach a distinct cluster  $V_2(A)$  by advancing a *single* edge  $(s, s')$  along  $\gamma$ . Clearly, this edge must be in  $E_B$  and thus there exists a cluster  $V_1(B)$  containing both  $s$  and  $s'$ . Then by induction  $\{V_1(A), V_1(B), V_2(A), \dots, V_l(A)\}$  is a local overlapping cover along  $\gamma$ .  $\square$

So in order to find a suitable bipartition with  $\mathcal{G}_e(A) = \mathcal{G}_e(B) = \mathcal{G}_e$  we simply need to select a subset  $E_A$  of edges such that both  $E_A$  and  $E_B = E \setminus E_A$  have the single cluster property. Indeed, for the square, triangular and Kagome lattice such bipartitions exist, see figures 4(a) and 6 for examples. Using theorem 9 we then obtain

**Theorem 10.** *For any quantum double model on a square, triangular or Kagome lattice  $\Sigma$ , the ground state  $|\Omega\rangle$  has the GE*

$$E_G = \alpha|V| - \log_2 |G|, \quad (58)$$

where  $\alpha = \log_2 |G|$ .

Note that  $|V| = n_s$  (number of stars) is proportional to the volume of the surface.

From Aguado and Vidal [15] and the proof of theorem 2 it becomes clear that we can calculate the GE of blocks of linear size  $k \geq 2$  directly from the renormalized graphs  $\Gamma_k$  shown in figure 2(e). On a square lattice of linear size  $2k\lambda$  there are  $N_k = 2\lambda^2$  such blocks, and the renormalized graph  $\Gamma_k$  has  $|V_k| = N_k(k+1)$  vertices. We merely state the result for the GE of blocks.

**Theorem 11.** *For any quantum double model on a square lattice  $\Sigma$ , the ground state  $|\Omega\rangle$  has the GE of blocks of size  $k$*

$$E_G = \alpha|V_k| - \log_2 |G|, \quad (59)$$

where  $\alpha = \log_2 |G|$ .

Clearly,  $|V_k|$  is proportional to *both* the renormalized volume of the surface *and* the boundary area of a single block (with a natural value of  $4k$ ). This is exactly the generalization of (14) to the non-Abelian case. As a final remark, we notice that the topological contribution to the GE for quantum double models is given by  $E_\gamma = \log_2 |G|$ . This again coincides with the corresponding value of the topological entanglement entropy.

## 9. $E_\gamma$ away from fixed renormalization group points

So far we just discussed the topological contribution to the GE for specific ground states of models that turn out to be RG fixed points, and therefore representative of their respective topological phases. But, what if we are away from the RG fixed point? Is  $E_\gamma$  robust under a perturbation? In this section we address briefly this question, and arrive to the conclusion that, indeed,  $E_\gamma$  is a robust property of the topological phase.

There are several ways of checking the robustness of  $E_\gamma$  under perturbations. One possibility is to perform a numerical analysis. In this sense, it should be possible to do large-scale calculations using tensor network methods. While we shall not carry out these calculations explicitly in this paper, we will explain different strategies based on tensor network algorithms in the appendix. Nevertheless, we provide a small-size exact calculation at the end of this section. Another option is a perturbation theory analysis. Yet, a more intuitive alternative is the following argumentation based on RG fixed points: all the models considered here are RG fixed points and hence representatives of their respective topological phases. The long-distance properties of any state in one of these phases do not change under local RG transformations and, thus, are equivalent to those of the fixed point. This, in particular, is true for the long-range

pattern of entanglement, and hence for  $E_\gamma$ . Thus, any non-relevant and short-range perturbation driving the Hamiltonian away from the fixed point will produce ground states with the same  $E_\gamma$ . Nevertheless, we expect a change in the short-range pattern of entanglement, and hence in the bulk term corresponding to a boundary law. For short-range perturbations the change involves modifications of the pre-factor of the boundary law as well as the possible appearance of sub-leading  $O(L^{-\nu'})$  corrections. Hence, equation (4) applies when away from the fixed point.

To double-check the above claim, we now perform a simple perturbation theory analysis of the robustness of the GE of the toric code model under external short-range perturbations such as magnetic fields. This analysis provides upper and lower bounds for the GE, and complements the above argumentation on the robustness of  $E_\gamma$  based on RG.

Let us then add a perturbation to the toric code Hamiltonian on the square lattice, and see how the ground state changes. For simplicity, we consider the case of an infinite plane, where the ground state  $|0, 0\rangle$  is non-degenerate, and hence we can use non-degenerate perturbation theory. The perturbed Hamiltonian will be

$$H^\lambda = H_{\text{TC}} + \lambda V, \quad (60)$$

where  $H_{\text{TC}}$  is the toric code Hamiltonian,  $V$  is the perturbation and  $\lambda \ll 1$ . Non-degenerate perturbation theory says that the new ground state can be approximated as

$$|0, 0\rangle^\lambda \approx |0, 0\rangle + \lambda \sum_{\phi, c} \frac{\langle \phi, c, 0, 0 | V | 0, 0 \rangle}{E_{0,0} - E_{\phi, c}} |\phi, c, 0, 0\rangle, \quad (61)$$

where  $E_{0,0}$  is the ground state energy and  $E_{\phi, c}$  is the energy of the excited state  $|\phi, c, 0, 0\rangle$ .

We now consider the case in which the perturbation is an homogeneous magnetic field, e.g. in the  $x$  direction,

$$V = \sum_{j=1}^n \sigma_x^{[j]} \quad (62)$$

(the case of  $z$  and  $y$  directions can be considered similarly). It is easy to check that in this case, the normalized perturbed ground state becomes

$$|0, 0\rangle^\lambda \approx C \left( |0, 0\rangle - \frac{\lambda}{\Delta} \sum_{j=1}^n \sigma_x^{[j]} |0, 0\rangle \right) \quad (63)$$

with  $\Delta$  the energy gap to create a pair of flux and anti-flux quasiparticles, and  $C = (1 + n\lambda^2/\Delta^2)^{-1/2}$  a normalization constant.

Our aim now is to estimate the maximum overlap of the previous state with a product state of blocks of boundary size  $L$ . This can be done as follows: first, and as in the unperturbed case, we apply CNOT operations locally inside of the blocks so that qubits are disentangled in the unperturbed ground state. By doing this, we can focus on the entanglement of the state

$$|0, 0\rangle_{\text{disentangled}}^\lambda \approx C \left( |\widetilde{0}, \widetilde{0}\rangle - \frac{\lambda}{\Delta} \sum_{j=1}^n \sigma_x^{[j]} |\widetilde{0}, \widetilde{0}\rangle \right) \otimes |e_1\rangle \otimes \cdots \otimes |e_p\rangle, \quad (64)$$

where  $|e_k\rangle$  is the quantum state for the  $k$ th disentangled qubit. The above equation is indeed equivalent to

$$|0, 0\rangle_{\text{disentangled}}^\lambda \approx C \left( |\widetilde{0}, 0\rangle \otimes |e_1\rangle \otimes \cdots \otimes |e_p\rangle - \frac{\lambda}{\Delta} \sum_{j=1}^{n_b} S_x^{[j]} |\widetilde{0}, 0\rangle \otimes |e_1\rangle \otimes \cdots \otimes |e_p\rangle - \frac{\lambda}{\Delta} |\widetilde{0}, 0\rangle |\omega_{1,\dots,p}\rangle \right), \quad (65)$$

where  $S_x^{[j]}$  is the total spin in the  $x$  direction for the  $L$  spins in the boundary of block  $j$ , and

$$|\omega_{1,\dots,p}\rangle = \sum_{j=1}^p \sigma_x^{[j]} |e_1\rangle \otimes \cdots \otimes |e_p\rangle. \quad (66)$$

Now we find upper and lower bounds to the maximum overlap of this state with a product state of the blocks. A lower bound can be easily obtained by the product state  $|0\rangle^{\otimes(n-p)} \otimes |e_1\rangle \otimes \cdots \otimes |e_p\rangle$ . Noticing that  $|e_k\rangle$  is either  $|0\rangle$  or  $|+\rangle$  [15], and that the  $|+\rangle$  contributions come only from qubits close to the boundary of the block, we have that

$$C |\widetilde{g}|^{-1/2} \left( 1 - \frac{\omega n_b L \lambda}{\Delta} \right) \leq \Lambda_{\text{max}}^{\text{[blocks]}}, \quad (67)$$

for some positive  $\omega = O(1)$  constant. The following upper bound can also be found easily:

$$\Lambda_{\text{max}}^{\text{[blocks]}} \leq C |\widetilde{g}|^{-1/2} \left( 1 + \frac{n_b L \lambda}{\Delta} + \frac{\lambda}{\Delta} \right). \quad (68)$$

Using the above bounds, one can check that for the GE we obtain, in the limit  $\lambda \ll \Delta$  and  $L \gg 1$ ,

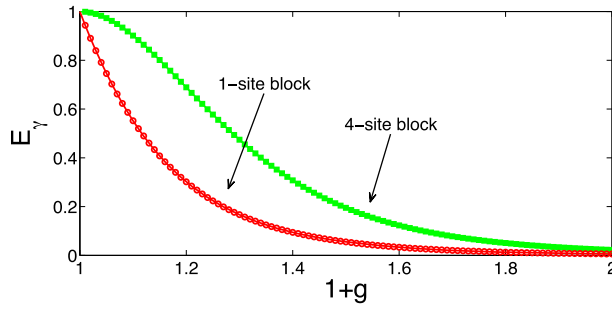
$$\left( \frac{1}{4} + \frac{2\omega\lambda}{\Delta} \right) n_b L - 1 \geq E_G^\lambda \geq \left( \frac{1}{4} - \frac{2\lambda}{\Delta} \right) n_b L - 1. \quad (69)$$

The above equation is compatible with a leading change in the GE in the pre-factor of the boundary law. Also, the fact that both bounds leave the topological component  $E_\gamma = 1$  untouched seems to indicate that this is actually robust under the perturbation. Moreover, implementing finite- $L$  corrections to these bounds provides  $O(L^{-\nu'})$  corrections, which is consistent with our previous claims.

In order to further check the above arguments, we have computed numerically the GE for the toric code with  $n$  spins on the square lattice, exactly, for sizes up to  $n = 16$ , and where we added a perturbation that amounts to introducing a string tension in the Hamiltonian. More specifically, we perturbed the system by considering the PEPS tensors of the toric code ground state  $|0, 0\rangle$  (see Verstraete *et al* [14]) and modifying the non-zero components that correspond to having ‘spin up’ in all sites. That is, the components  $A_{\alpha\beta\gamma\delta}^i$  of the perturbed-PEPS tensors are given by

$$\begin{aligned} A_{1111}^\uparrow &= 1 + g, & A_{2211}^\downarrow &= 1, \\ A_{2222}^\uparrow &= 1 + g, & A_{1122}^\downarrow &= 1, \end{aligned} \quad (70)$$

where the upper tensor index is the physical index and the rest are the bond indices (in the PEPS there is also another tensor like this but rotated  $90^\circ$ , so that we have an  $ABAB \dots$  periodicity with a two-site unit cell). For  $g = 0$  one recovers the unperturbed toric code, whereas for  $g = \infty$  one has a polarized state in the  $z$ -direction. This is the same kind of perturbation that was



**Figure 12.** Estimated topological contribution to GE for the  $|0, 0\rangle$  ground state of the toric code on a square lattice, for systems up to 16 sites, for one-site (red) and four-site (green) blocks.

considered previously in [35], which adds a string tension with a somehow similar effect to adding a magnetic field in the  $z$ -direction to the Hamiltonian.

For such a perturbed topological state we did an exact calculation of the closest product state, for one-site blocks (single spins) and also for four-site blocks, up to  $n = 16$  (i.e. up to 16 one-site blocks, and up to 4 four-site blocks). From this we estimated  $E_\gamma$  by considering the GE as a function of the number of blocks, fitting it to a straight line and extrapolating the number of blocks to zero. Calculations for larger systems required a significant amount of computational resources, and therefore could not be implemented with this approach (for these one would need e.g. the tensor network techniques explained in the appendix). A summary of our results can be found in figure 12. Here we see that in the case of single-site blocks, the topological GE is not stable under the perturbation: it drops very quickly to zero without any sign of phase transition instead of staying close to  $E_\gamma = 1$  for a while. However, the four-site block calculation shows that  $E_\gamma$  also drops to zero, but quite slower than in the case of single-site blocks. This is a clear indication that in the case of four-site blocks the topological GE is more robust under perturbation. We expect, thus, that as the size of the blocks becomes larger, the topological contribution to the GE becomes more robust. At this point we are constrained by the sizes in the exact calculations, but we take these results as a first-principle indication that the topological contribution to the GE tends to be robust under perturbation for sufficiently large block sizes, which is in fact compatible with the RG and perturbation theory arguments above.

## 10. Conclusions

In this paper, we have studied the GE in various topologically ordered states that correspond to fixed-points of RG. These are the toric code, double semion, colour code and quantum double models, in a variety of lattices. Generically, we found that the GE is typically composed of a boundary law term times the number of blocks plus a topological term. This is remarkable, since it is a signature of TO in multipartite entanglement properties of the state, rather than bipartite. The topological term for the colour code is twice that for the toric code, which is consistent with the recent result that the former model is equivalent to two copies of the latter [13]. Away from RG fixed-points, we argued that the same type of behaviour holds up to a possible subleading term, and thus the topological contribution is a robust property of topological phases of matter. The numerical evaluation of  $E_\gamma$  for non-analytic cases is possible. This could be done e.g. in the context of Tensor Network methods (such as PEPS algorithms), as is usually done for



the topological entropy [28]. We leave the specific implementation of these methods for future work.

In all the cases considered here, the topological GE turns out to be equivalent to the topological entropy. It would be good to prove, in general, whether this is always true or not, in order to understand if the GE can provide more information about TO than the one that is already in the topological entropy. However, even if this were not the case, the GE would still be a useful tool to extract the quantum dimensions of the anyon model associated to a given topologically ordered state. For scenarios where reduced density matrices are hard to evaluate numerically, one may thus believe that looking at the GE may be more efficient. Moreover, the GE may also be a useful tool in order to extract *minimally entangled states* within the ground subspace, and hence the complete topological characterization of the system<sup>9</sup>.

Finally, we believe that other models could possibly be analysed with the techniques that we used in this paper. Specifically, similar results should also apply to the A-phase of Kitaev's honeycomb model [29] for which the toric code is (in some limiting cases) an effective model, as well as spin-liquid states with an emergent  $\mathbb{Z}_2$  gauge symmetry [30]. The possibility of studying the ground states of string-net models of Levin and Wen [5] in general, as well as topological quantum field theories [3], is also left for future investigation.

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## Appendix A. Relative entropy of entanglement

We shall see in the following that for most of the ground states discussed in this paper, they can also be characterized by the so-called relative entropy of entanglement and in fact their value is identical to that of GE. The relative entropy  $S(\rho||\sigma)$  between two states  $\rho$  and  $\sigma$  is defined via

$$S(\rho||\sigma) \equiv \text{Tr}(\rho \log_2 \rho - \rho \log_2 \sigma), \quad (\text{A.1})$$

which is evidently not symmetric under exchange of  $\rho$  and  $\sigma$ , and is non-negative, i.e.  $S(\rho||\sigma) \geq 0$ . The relative entropy of entanglement (RE) for a mixed state  $\rho$  is defined to be the minimal relative entropy of  $\rho$  over the set of separable mixed states [31, 32]:

$$E_R(\rho) \equiv \min_{\sigma \in \mathcal{D}} S(\rho||\sigma) = \min_{\sigma \in \mathcal{D}} \text{Tr}(\rho \log_2 \rho - \rho \log_2 \sigma), \quad (\text{A.2})$$

<sup>9</sup> Work in progress.

where  $\mathcal{D}$  denotes the set of all separable states. With any separable state  $\sigma$ , we can obtain an immediate upper bound on  $E_R(\rho) \leq \text{Tr}(\rho \log_2 \rho - \rho \log_2 \sigma)$ . For pure state  $\rho = |\Psi\rangle\langle\Psi|$ , such as the ground state of the toric code and the colour code, we have further that

$$E_R(|\Psi\rangle\langle\Psi|) \leq -\langle\Psi| \log_2 \sigma |\Psi\rangle. \quad (\text{A.3})$$

In general, the task of finding the RE for arbitrary states  $\rho$  involves a minimization over all separable states, and this renders the computation of the RE very difficult. It has been shown that for pure state  $|\Psi\rangle$ , the its GE lower bounds its REE, i.e.  $E_R(\Psi) \geq E_G(\Psi)$  [33]. Using the technique in [33], we show that the GE is identical to the REE for the ground states considered here.

Let us illustrate this by the  $|0, 0\rangle$  ground state of the toric code on the square lattice. Since  $|0, 0\rangle = \sum_{g \in \mathcal{G}_s} g|0 \dots 0\rangle / \sqrt{|\mathcal{G}_s|}$ , where  $\mathcal{G}_s$  is the group generated by star operators, each component  $|g\rangle \equiv g|0 \dots 0\rangle$  is a product state. We can thus construct a mixed separable state

$$\sigma = \frac{1}{|\mathcal{G}_s|} \sum_{g \in \mathcal{G}_s} |g\rangle\langle g|, \quad (\text{A.4})$$

and hence obtain an upper bound on the REE,

$$\begin{aligned} E_R(|0, 0\rangle\langle 0, 0|) &\leq \log_2 |\mathcal{G}_s| \langle 0, 0| \sum_{g \in \mathcal{G}_s} (|g\rangle\langle g|) |0, 0\rangle \\ &= \log_2 |\mathcal{G}_s|. \end{aligned} \quad (\text{A.5})$$

This upper bound turns out to be identical to the value of the GE, a lower bound on REE. Therefore, we have that  $E_R(|0, 0\rangle\langle 0, 0|) = E_G(|0, 0\rangle\langle 0, 0|) = n_s - 1$ . Such an equality of GE and REE can also be understood from the group symmetry [34]. We can straightforwardly apply similar arguments to ground states of the double semion, colour code and the quantum double models.

## Appendix B. Some numerical tensor network algorithms to compute $E_g$

In what follows we sketch, at a conceptual level, several methods based on tensor networks to numerically compute the topological GE. For concreteness we focus on the case of having a representation of the relevant topological quantum state given by a 2D PEPS (see e.g. verstrate *et al* [14]). We assume that the system is finite, translationally invariant and is defined on the surface of a torus. We also assume that the size of the system is sufficiently large, so that relatively large blocks can be considered.

Within this framework, the topological GE can be extracted by e.g. doing a finite-size scaling of the GE of blocks, with respect to different block sizes. This GE of blocks needs to be computed by optimizing the fidelity of the 2D PEPS on the torus with a product state of the blocks. The key point is how to implement this optimization efficiently, considering that the number of sites within each block can be quite large, and that the total size of the system is also large.

Here two different strategies are proposed. These are based on renormalizing/not renormalizing the tensors within the blocks.

1. *Product-PEPS/MPS*. We consider a product state of the blocks such that the quantum state for each block is itself a finite  $L \times L$  PEPS, or a (snake) matrix product state (MPS) for  $L^2$

sites. The maximization of the fidelity is then carried over the tensors of these finite-PEPS or finite-MPS, which can be implemented efficiently.

2. *Tensor-renormalization.* Another approach is to consider again blocks of size  $L \times L$ , but where we compute a renormalized PEPS tensor for the block. This could be achieved by using different tensor-renormalization strategies, such as e.g. SRG [36] and HOSRG [37] but adapted to a finite 2D system with open boundary conditions. As a result of this tensor-renormalization, the topological quantum state is represented by a ‘renormalized’ PEPS where each tensor corresponds to a block. Once this is computed successfully (which may be non-trivial), the necessary optimization to compute the GE can be done using a product state directly for the renormalized single sites.

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