

I N S T I T U T F Ü R P L A S M A P H Y S I K
G A R C H I N G B E I M Ü N C H E N

Solution of an ABEL type integral
equation in the presence of noise.

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IPP/6/29

November 1964

Die nachstehende Arbeit wurde im Rahmen des Vertrages zwischen dem Institut für Plasmaphysik GmbH und der Europäischen Atomgemeinschaft über die Zusammenarbeit auf dem Gebiete der Plasmaphysik durchgeführt.

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IPP/6/29

R. Gorenflo
Y. KovetzSolution of an ABEL type
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November, 1964 (in English).

ABSTRACT:

The solution $i(r)$, $0 \leq r \leq 1$, of the integral equation

$$J(x) = 2 \int_{r=|x|}^1 \frac{i(r) r dr}{\sqrt{r^2 - x^2}},$$

arising in spectroscopy, can be obtained from $J(x)$ by a half-order differentiation process after substituting $r^2 = 1-s$, $x^2 = 1-t$. Therefore noise in the measured values of J is amplified when computing i by usual numerical methods. In the computation of $i(r)$ two undesirable effects may arise:

(a) lack of smoothness, (b) intervals where $i(r) < 0$, although for physical reasons we should have $i(r) \geq 0$.

(a) and (b) can be avoided as follows:- fitting $J(x)$ as the sum of an optimum number of suitable orthogonal functions, and using the extra information $i(r) \geq 0$ as a restriction, leads by discretization to a quadratic programming problem for the coefficients of the fitting expansion. A modification of this method is applicable when, instead of a continuous $J(x)$ only a few (e.g. 7 or 8) mean values of J (contaminated with noise) are measured over adjacent intervals in $-1 \leq x \leq 1$.

A. The case of continuously measured $J(x)$

1. Introduction

The integral transform

$$(1) \quad i(r) = -\frac{1}{\pi} \int_{x=r}^1 \frac{dJ(x)}{\sqrt{x^2 - r^2}}$$

where $-1 \leq x \leq 1$, $0 \leq r \leq 1$, $J(1) = 0$, $J(-x) = J(x)$, solves the ABEL type integral equation

$$(2) \quad J(x) = 2 \int_{r=|x|}^1 \frac{i(r) r dr}{\sqrt{r^2 - x^2}},$$

which arises in spectroscopy when, by "side-on" observation of a cylindrical source, a laterally integrated intensity $J(x)$ is measured from which the radial intensity $i(r)$ is to be computed (see for example [1] or [5], where further references are also given). Fig. 1 illustrates the geometry.

Substituting

$$(3) \quad r^2 = 1-s, \quad x^2 = 1-t, \quad i(r) = g(s), \quad J(x) = G(t)$$

in (1) yields

$$(4) \quad g(s) = \frac{1}{\pi} \int_{t=0}^s (s-t)^{-1/2} dG(t)$$

which, in terms of operational calculus ([4], 290 etc.), is half-order differentiation of G . Therefore noise of the measured function $J(x)$ appears amplified in the computed function $i(r)$. Two unwanted effects may arise when applying usual numerical procedures for computing $i(r)$:

- a) lack of smoothness in $i(r)$, that is irregular high-frequency oscillations, superimposed on $i(r)$, which do not have a physical meaning.
- b) intervals in which $i(r) < 0$, although for physical reasons we should have $i(r) \geq 0$ everywhere.

In figures 2 and 3 examples of effects (a) and (b) are presented.

A heuristic theory of noise amplification is developed in [5], 22-27. The essential results are as follows:

Inverting (2) by any standard discrete method is a process of linear transformation of a vector (J_0, J_1, \dots, J_N) representing readings of J in equidistant points x , into a vector (i_0, i_1, \dots, i_N) . By a suitable idealization the errors of J are

$$(5) \quad \delta_n = \varepsilon_n + \eta_n,$$

where the strongly correlated ε_n stand for the systematic error of the continuously working measuring device, whereas the η_n are reading and rounding errors of the scanning process giving (J_0, J_1, \dots, J_N) from the measured continuous function $J(x)$. The smooth ε_n may yield effect (b); the η_n are responsible for (a).

Let us consider the η_n only (this is justified if the ε_n are sufficiently small). Because of linearity we can approximately evaluate their influence by applying (1) to

$$(6) \quad J(x) = \begin{cases} \eta_{j-1}, & (j-1)/N \leq x < j/N, \quad 1 \leq j \leq N, \\ 0, & x = 1, \end{cases}$$

where the η_j are uncorrelated random variables with expectation 0 and standard deviation σ . The worst case is to be expected where $r=0$. From

$$(7) \quad i(0) = \frac{N}{\pi} \left(\eta_0 - \sum_{j=2}^N \frac{\eta_{j-1}}{(j-1)j} \right)$$

we conclude that $i(0)$ has expectation 0 and standard deviation

$$(7') \quad CN\sigma \text{ with } C = \frac{1}{\pi} \left\{ 1 + \sum_{j=2}^N \frac{1}{(j-1)^2 j^2} \right\}^{1/2}.$$

We have $C \approx 0.36$ for $N \gg 1$. MONTE CARLO calculations confirmed that this is realistic in magnitude.

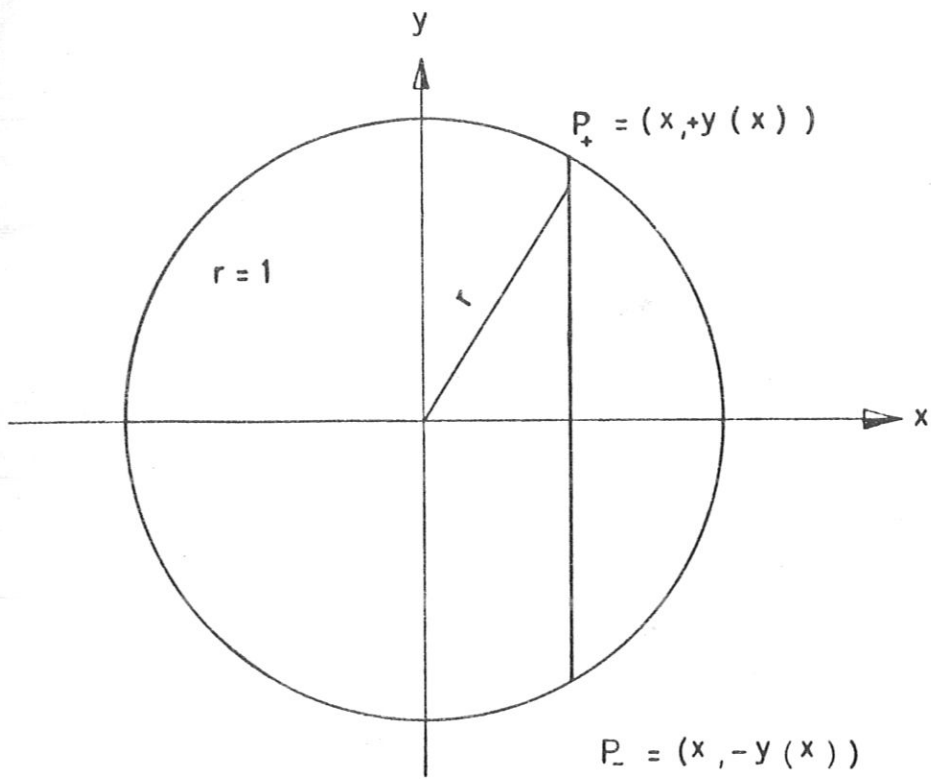


Fig 1

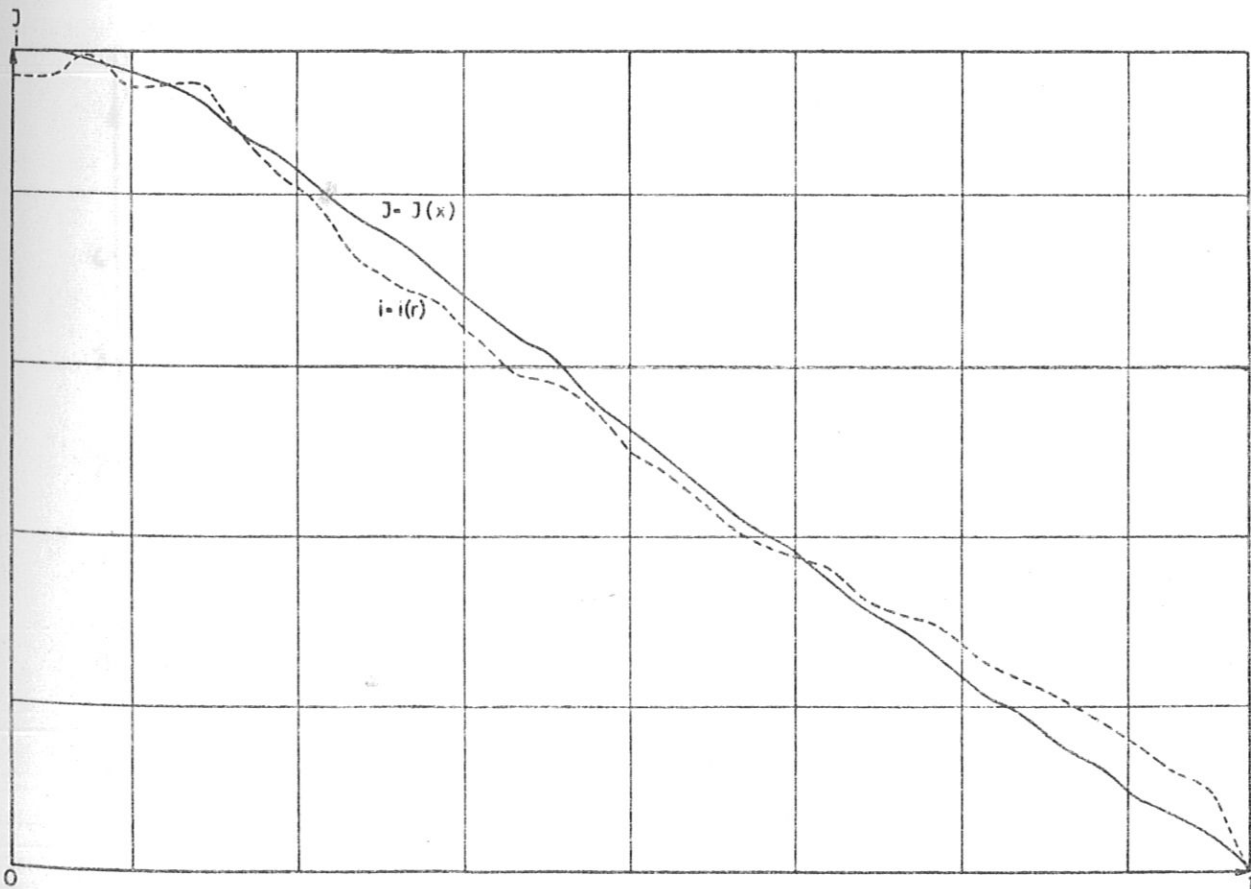


Fig 2

An illustration of effect (a): $i(r)$ not smooth.

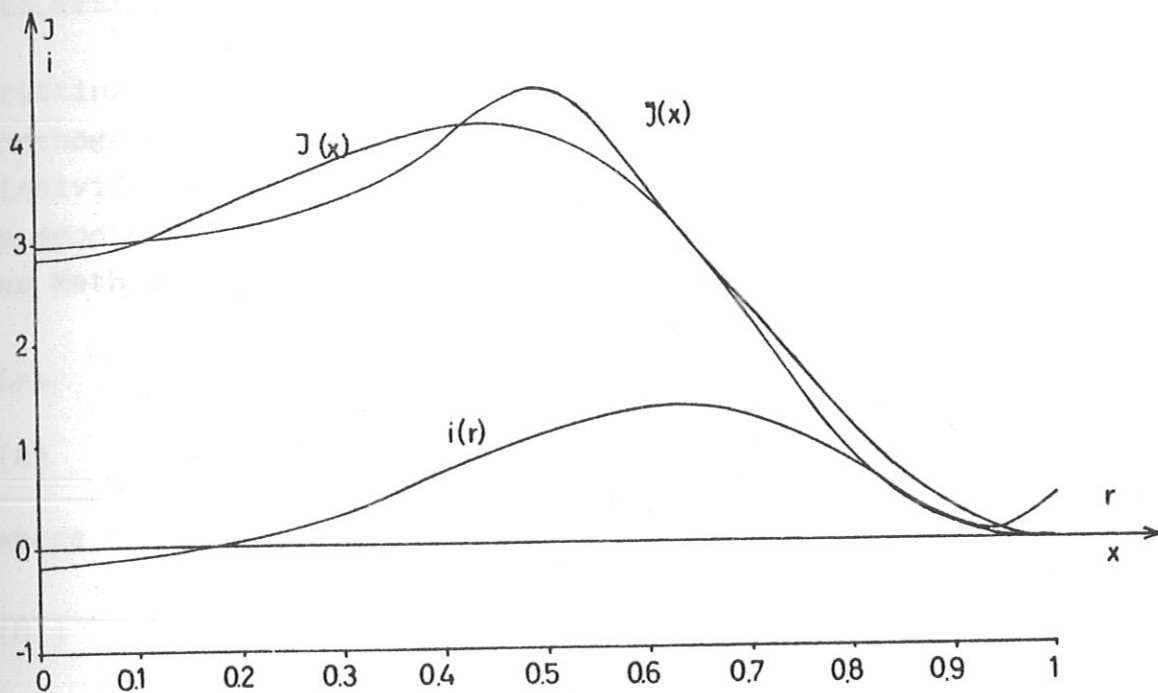


Fig.: 3

An illustration of effect (b): $i(r)$ negative near $r = 0$, $\bar{J}(x)$ = measured function, $J(x)$ = unrestricted least-squares fit to $\bar{J}(x)$, $k = 4$.

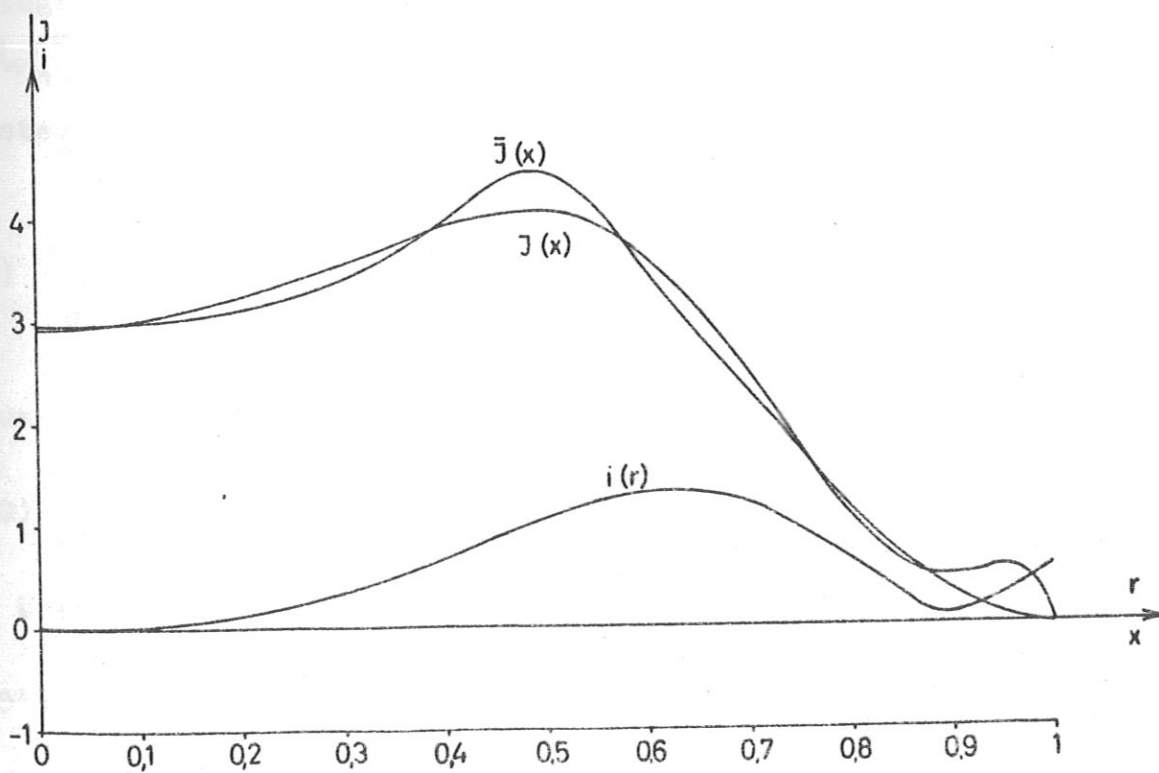


Fig.: 4

$\bar{J}(x)$ as in fig. 3, $J(x)$ = restricted least-squares fit, $i(r) \geq 0$ everywhere, $k = 4$.

2. Smoothing by using GEGENBAUER polynomials

Fitting $J(x)$ as the sum of an optimum number of suitable orthogonal functions, for which the transforms (1) are individually known, is a natural procedure for obtaining a smooth function $i(\tau)$. One such procedure is developed as Method D in [5]; numerical experiences are also given.

Let

$$(8) \quad W_n^*(x) = \sqrt{1-x^2} U_n(x),$$

where

$$(8') \quad U_n(x) = \sum_{p=0}^n \beta(p, n) x^{2p},$$

with coefficients

$$(8'') \quad \beta(p, n) = 4^p \binom{-3/2}{n+p} \binom{n+p}{n-p},$$

is (in TRICOMI's notation, [8], 177 etc.) the GEGENBAUER polynomial $C_{2n}^{3/2}(x)$. For even functions $J(x)$, square-integrable in $-1 \leq x \leq 1$, the system

$\{W_n^*, n=0, 1, 2, \dots\}$ is a complete orthogonal system with

$$(9) \quad \int_0^1 W_m^*(x) W_n^*(x) dx = \begin{cases} 0, & m \neq n \\ \frac{(2n+1)(2n+2)}{4n+3}, & m = n. \end{cases}$$

Thus

$$(10) \quad J(x) = \sum_{n=0}^k c_n W_n^*(x)$$

is transformed by (1) into

$$(10') \quad i(\tau) = \frac{1}{2} \sum_{n=0}^k c_n P_n^*(\tau)$$

where

$$(11) \quad P_n^*(r) = \sum_{p=0}^n \gamma(p, n) r^{2p}.$$

The γ may be computed by solving the triangular systems of linear equations

$$(12) \quad \sum_{\lambda=p}^n b(p, \lambda) \gamma(\lambda, n) = \beta(p, n), \quad 0 \leq p \leq n, \quad n \geq 0,$$

where

$$(13) \quad b(p, n) = \sum_{j=n-p}^n \frac{(-1)^{j+p-n}}{2j+1} \binom{n}{j} \binom{j}{n-p}.$$

The numerical procedure consists of fitting to $J(x)$ a least-squares approximation

$$(14) \quad \sum_{n=0}^k c_n W_n^*(x) \approx J(x),$$

where the FOURIER coefficients

$$(14') \quad c_n = \frac{4n+3}{(2n+1)(2n+2)} \int_0^1 J(x) W_n^*(x) dx$$

are computed by the SIMPSON integration formula. The measurements are to be taken at $x_j = j/N$, $0 \leq j \leq N$, and N must be an even integer.

By this method we obtained sufficiently smooth functions $i(r)$ in actual cases, with $20 \leq N \leq 30$, $k \approx 5$.

Unfortunately $i(r)$ occasionally had intervals of negativity, especially in such cases where $J(x)$ had a local minimum at $x=0$. This was due, either to excessive measurement errors or to errors introduced by truncation of the series, or both. Of course k should not be too large in the presence of noise. If k is too

large, the unphysical high-frequency oscillations are fitted too well. However, if k is just sufficiently small that the noise is suppressed it may be so small that the truncation error is unacceptable.

3. Avoiding intervals in which $i(r) < 0$ by quadratic programming

The idea that the extra information $i(r) \geq 0$ can be used as a restriction on the fit (14) leads to a way out of the dilemma mentioned at the end of § 2.

Discretization yields a quadratic programming problem:

$$(15) \quad \begin{aligned} & \text{Choose } c_0, c_1, \dots, c_k \quad \text{in order to} \\ & \text{minimize } \sum_{j=0}^{N-1} \left\{ J(j/N) - \sum_{n=0}^k c_n W_n^*(j/N) \right\}^2 \\ & \text{subject to } \sum_{n=0}^k c_n P_n^*((j-1)/m) \geq 0, \quad j=1, 2, \dots, m. \end{aligned}$$

Implementing this idea a double-precision FORTRAN routine was written and successfully applied with values of N between 20 and 30, and $m \approx 20$, $k \approx 5$. Problem (15) was solved by the iterative method of HILDRETH and D'ESOPPO ([7], 73-79). Average computer (IBM 7090) running time depends on N and m ; for $N \approx 25$ and $k \approx 5$ it was approximately 1 minute.

In figure 4 an example is presented.

We give some details of transforming (15) into the standard form of a quadratic programming problem. With

$$(16) \quad J_j = J((j-1)/N), \quad w(n, j) = W_n^*((j-1)/N), \quad 1 \leq j \leq N,$$

we have to minimize

$$(17) \quad D(c) = \sum_{j=1}^k \left(J_j - \sum_{n=0}^k \omega(n,j) c_n \right)^2$$

subject to the restrictions

$$(18) \quad \sum_{n=0}^k c_n P_n^*(r_\lambda) \geq 0, \quad r_\lambda = (\lambda-1)/m, \quad \lambda=1, 2, \dots, m.$$

With $0 \leq i \leq k$, $0 \leq n \leq k$, $1 \leq \lambda \leq m$,

$$Q_{in} = Q_{ni} = \sum_{j=1}^k \omega(i,j) \omega(n,j),$$

$$b_n = -2 \sum_{j=1}^k \omega(n,j) J_j, \quad A_{\lambda n} = -P_n^*(r_\lambda)$$

we get the problem:-

$$(19) \quad \begin{array}{l} \text{Minimize} \quad D = c' Q c + b' c \\ \text{subject to} \quad A c \leq 0. \end{array}$$

Further details may be found in § 6 and § 7. In these paragraphs we have simply to replace T_j by J_j , $\delta(n,j)$ by $\omega(n,j)$.

The matrix Q is strictly positive definite, when $N \geq k+1$.

In order to prove it we have only to prove that it is regular, since by (17)

$$\sum Q_{in} c_i c_n = \left(\sum_{v=0}^k c_v \vec{w}_v \right)^2, \quad \text{where } \vec{w}_v = (\omega(v,1), \dots, \omega(v,N)),$$

so that Q is positive semi-definite.

By definition

$$Q = W \cdot W'.$$

We shall prove first that the rank of W , $r(W) = k+1$.

W is a matrix of $(k+1)$ rows and N columns.

$$\omega(n,j) = M_j \sum_{p=0}^n \beta(p,n) s_j^p \quad \left(\begin{array}{l} n=0, 1, \dots, k \\ j=1, 2, \dots, N \end{array} \right)$$

where $s_j = \left(\frac{j-1}{N}\right)^2$

and $M_j = \sqrt{1-s_j} \neq 0$.

Let λ_n ($n=0, 1, \dots, k$) be $(k+1)$ constants satisfying

$$(a) \sum_{n=0}^k \lambda_n M_j \sum_{p=0}^n \beta(p, n) s_j^p = 0 \quad \text{for } j=1, \dots, N.$$

Because $M_j \neq 0$ we get, by changing the order of summation,

$$(a') \sum_{p=0}^k \left(\sum_{n=p}^k \lambda_n \beta(p, n) \right) s_j^p = 0 \quad (j=1, 2, \dots, N).$$

The expression

$$p^*(s) = \sum_{p=0}^k \left(\sum_{n=p}^k \lambda_n \beta(p, n) \right) s^p$$

defines a polynomial of degree k , and therefore cannot have more than k zeros, unless $p^*(s) \equiv 0$.

The system (a') expresses the fact that $p^*(s)$ vanishes for N different values of s , and therefore it must be identically zero.

This leads us to the triangular system of $k+1$ homogeneous equations, with the $k+1$ unknowns λ_n ,

$$\sum_{n=p}^k \lambda_n \beta(p, n) = 0, \quad p=0, 1, \dots, k,$$

whose diagonal coefficients are, by (8"), $\beta(n, n) = 4^n \binom{-3/2}{2n} \neq 0$.

Therefore $\lambda_n = 0$ ($n=0, 1, \dots, k$) and all the $k+1$ rows of the matrix W are linearly independent, i.e. $r(W) = k+1$.

Now, in order to prove that $Q = WW'$ is regular, we compute its determinant, making use of the algebraic theorem (see e.g. [9], pp. 64-65): -

"If A is matrix of order (m, n) , and B is a matrix of order (n, m) , and if $m \leq n$, then $\det(AB)$ is equal to the sum of products of the corresponding minors of A and B ".

The corresponding minors of W and W' are determinants of order $k+1$, each minor of one being formed from the transposed matrix of the minor of the other.

$\det(Q)$ is therefore the sum of $\binom{N}{k+1}$ squares of determinants of order $k+1$, and therefore $\det(Q) \geq 0$. As $r(W) = k+1$, W has at least one minor which is not zero, and therefore $\det(Q) > 0$.

If $N \geq 20$, we should have, of course, $k \leq 10$, because otherwise the fit would be insufficiently smooth.

B. The case where only a few mean values of $J(x)$ are measured

4. The problem of light-pipes

In this and the following paragraphs we consider the following problem:

Determine in $r \leq 1$ a radial non-negative intensity $i(r)$ from N "side-on" measurements,

$$(20) \quad T_j = \int_{I_j} J(x) dx, \quad 1 \leq j \leq N.$$

For convenience we put

$$(21) \quad i(r) = 0 \quad \text{in} \quad r > 1, \quad J(x) = 0 \quad \text{in} \quad |x| \geq 1.$$

The intervals I_j are defined by

$$(22) \quad \begin{cases} I_j = \{x \mid e_j < x < e_{j+1}\}, \\ e_j = \alpha + (j-1)h. \end{cases}$$

h is a positive constant depending on the experimental configuration, $\alpha = e_1$ is the left-hand endpoint of the first interval. $\alpha = e_1$ and $e_{N+1} = \alpha + Nh$ may or may not lie in $-1 < x < 1$, but e_2 and e_N must lie in this interval. From these conditions we obtain

$$-1-h < \alpha < 1-(N-1)h \quad \text{and} \quad 0 < h < 2/(N-2).$$

In an actual series of experiments with cylindrical plasma discharges we had $N=7$ or $N=8$. One possible configuration is sketched in fig. 5.

Because of the complicated nature of the experiment and some uncontrollable disturbances α is not exactly known, and the intensity $i(x, y)$ in $x^2 + y^2 \leq 1$ is not exactly a radial intensity $i(r)$. However, in order to draw conclusions from the results T_j of the experiment, we assume i to be a function of radius only. As pointed out by BRACEWELL in [2], it is impossible to determine an arbitrary intensity $i(x, y)$ by side-on observation in one direction only.

The difficulty in finding an effective procedure for computing $i(r)$ is due also to the fact that T_j is not exactly equal to $\int_{I_j} J(x) dx$, but is disturbed by the "noise" of the measuring apparatus, which is estimated to be as great as 10 to 20% of the maximum of the T_j . We suppose the T_j to be disturbed independently of each other.

Two problems arise:

- 1) Determine α , or, equivalently, the centre of the circle in which $i(r)$ exists.

Primarily we have to determine the centre $r=0$ (that is the origin of our coordinate system), but for computational reasons we need the left-hand endpoint $\alpha = -\tilde{c} h$ of the first interval. \tilde{c} is the (not necessarily integral) number of intervals I_j to the left side of $x=0$ (fig. 6). For reasons to be explained in § 6 \tilde{c} should not be an integral multiple of $1/2$.

In the analysis of actual experimental results it was decided to fix the position of the centre on the basis of physical arguments, because of the lack of a satisfactory automatic method. This results in the introduction of some additional noise.

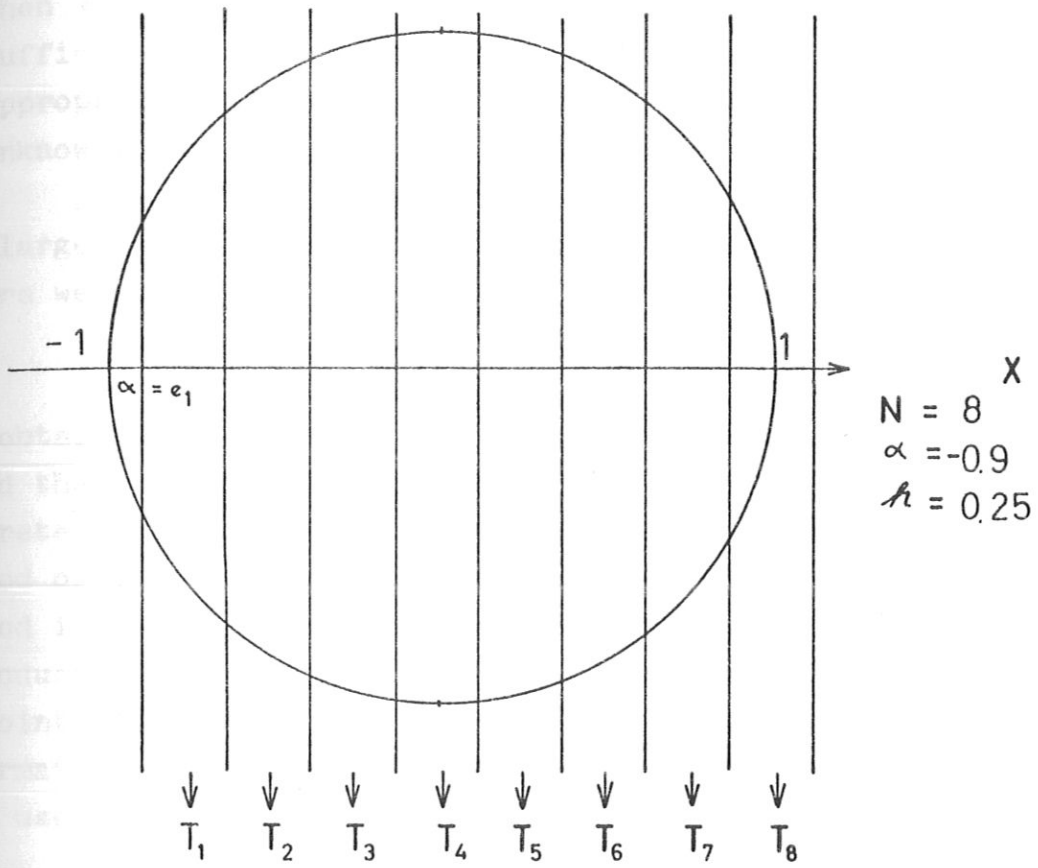
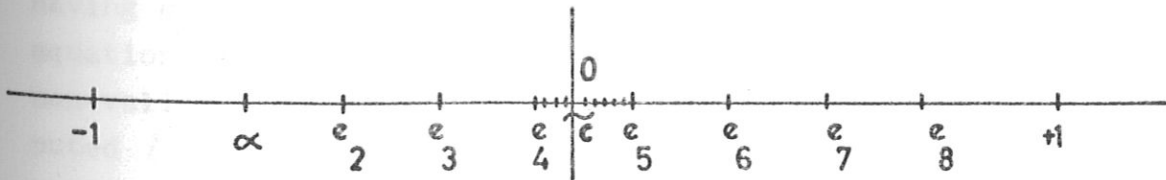


Fig. 5



$$\tilde{c} = 3,4$$

$$\alpha = -\tilde{c} h = -0,68$$

$$N = 7$$

$$h = 0.2$$

Fig. 6

- 2) When \tilde{c} , and therefore α , is given, determine a sufficiently smooth function $J(x)$, lying in an appropriate function space, and approximating the unknown, true function.

For large values of N and sufficiently small measurement errors we could put

$$J((e_j + e_{j+1})/2) \approx T_j/h$$

and obtain a smooth curve $J(x)$ by interpolation. $i(r)$ could then be computed, either by one of the usual discrete procedures, or by the quadratic programming method of § 3. However, for values of $N \lesssim 20$ this method is not to be recommended, because the error introduced by considering T_j/h as the value of J at the midpoint of the interval I_j is too large. We have the information expressed in (20), and it is wise to make full use of it in the computational procedure.

This naturally leads to the idea of approximating, in an appropriate function space S , by the least-squares principle, which in our case reads as follows:-

$$\text{Minimize } D = \sum_{j=1}^N (T_j - \int_{I_j} J(x) dx)^2,$$

(23)

where $J \in S$.

Having computed $J(x)$, we have to solve the integral equation (2) for $i(r)$. Because of the "noise", and the non-validity of our symmetry assumptions, the computed $i(r)$ may turn out to be negative in places. To avoid this difficulty we can impose upon our minimization problem the restriction

$$(24) \quad i(r) \geq 0 \quad \text{in } r \leq 1,$$

which may be approximated by

$$i(r_\lambda) \geq 0, \quad \lambda = 0, 1, 2, \dots, m,$$

where $\{r_\lambda\}$ is a sufficiently dense set in $r \leq 1$.

5. Approximation by a rational function

The three main qualitative forms of physically plausible functions $J(x)$ (arising in an actual series of cylindrical plasma discharges) are sketched in fig. 7. All of them are qualitatively obtainable by the rational function

$$(25) \quad J(x) = (1-x^2) \left\{ b_1 + \frac{b_2}{1 + b_3(x^2 - b_4)^2} \right\}$$

with suitable coefficients b_j . The problem now is to find that point $b = (b_1, b_2, b_3, b_4)$ in the four-dimensional coefficient space, which solves (23). The integrals $\int_{I_j} J(x) dx$ are expressible in terms of elementary functions, although the analytic calculation is tedious and the final expressions are very long.

We used two methods for minimizing $D = D(b)$.

α) Direct search (see [6]). This is a method based on the idea of evaluating sequential trial solutions, and comparing each solution with the best one previously obtained, together with a strategy for determining the best direction for the next trial.

β) SHARE-program SCOOP. This is a similar procedure for minimizing functions of n parameters, where a first "guess" is given.

It turned out that $D(b)$ has many local minima, and therefore it is very important to find a good first guess, i.e. a point which is sufficiently near to the absolute minimum (methods α) and β) yield a local minimum). We decided to try a random search for finding a suitable starting point, but it transpired that this required a very long computer running time (in the order of 30 minutes for a single curve). Furthermore, although the approximation of $J(x)$ was quite good, $i(r)$ was negative near $r=0$, in absolute value up to 25% of the maximum

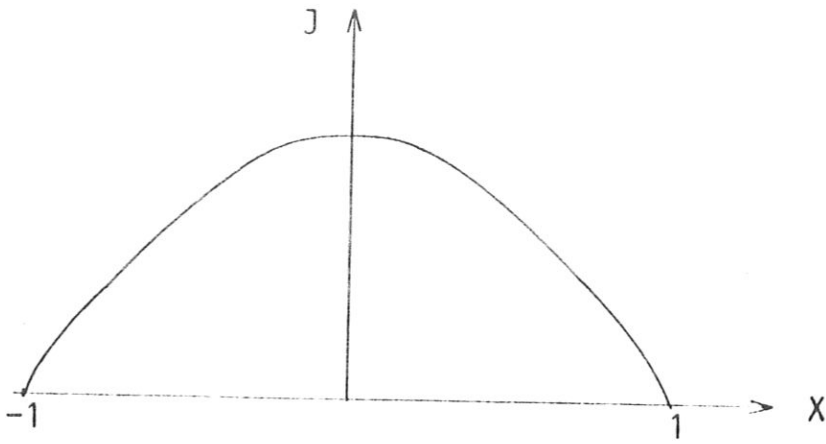
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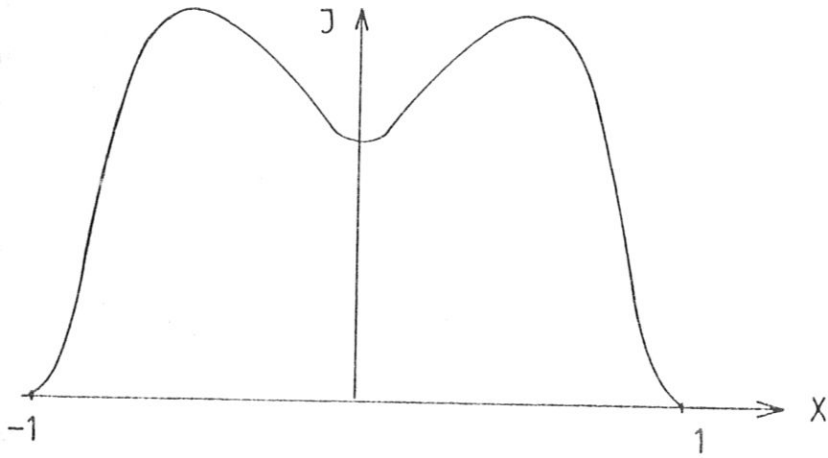
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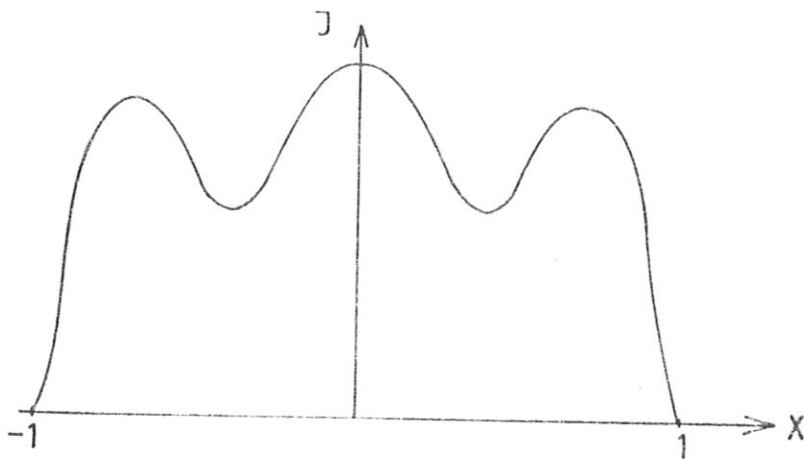


Fig. 7

positive value of $i(r)$. We do not have an efficient means for avoiding negativity by this method, for the function (25) is nonlinear in b .

We therefore consider these methods to be unsatisfactory.

6. Approximation using GEGENBAUER polynomials

We tried to solve (23) by minimizing D using the function

$$(10) \quad J(x) = \sum_{n=0}^k c_n W_n^*(x),$$

without the restriction $i(r) \geq 0$. The main properties of the W_n^* are given in § 2. For $N=8$ we tried the values of k from 3 to 8. There is no point in taking $k \geq N-1$, for then $D=0$, the approximation is too good, and for $k \geq N$ even not unique. Another reason why k should be not too large is that $J(x)$ has, as expected for physical reasons, only a few extrema and inflection points; so a too high k would lead to unwanted oscillations - a characteristic of polynomial approximation.

With

$$(26) \quad a_j = \begin{cases} e_j & \text{if } |e_j| \leq 1 \\ -1 & \text{if } e_j < -1 \\ 1 & \text{if } e_j > 1, \end{cases}$$

$$(27) \quad G_p(a, b) = \int_a^b \sqrt{1-x^2} x^{2p} dx,$$

and

$$(28) \quad \delta(n, j) = \int_{a_j}^{a_{j+1}} W_n^*(x) dx = \sum_{p=0}^n \beta(p, n) G_p(a_j, a_{j+1})$$

(the β are defined in § 2, formula (8"))
we have

$$(29) \quad \int_{T_j} J(x) dx = \sum_{n=0}^k c_n \delta(n, j).$$

The G_p may be recursively computed by the substitution $x = \sin t$ and using formula number 358 on page 321 of [3]. The results are

$$(30) \quad G_0(a, b) = \frac{1}{2} \left\{ \arcsin x \Big|_a^b + (x \sqrt{1-x^2}) \Big|_a^b \right\}$$

$$G_p(a, b) = \frac{1}{2p+2} \left\{ (x^{2p-1} (1-x^2)^{3/2}) \Big|_a^b + (2p-1) G_{p-1}(a, b) \right\}$$

for $p \geq 1$.

From (23) and (29) we get

$$(31) \quad D(c) = \sum_{j=1}^N \left\{ T_j - \sum_{n=0}^k \delta(n, j) c_n \right\}^2.$$

Clearly this quadratic form is positive semi-definite. The minimum of D may be found by solving the system of linear equations

$$\frac{\partial D}{\partial c_i} = 0, \quad 0 \leq i \leq k.$$

By differentiation and interchange of summation symbols we get

$$(32) \quad \sum_{n=0}^k Q(n, i) c_n = \sum_{j=1}^N \delta(i, j) T_j, \quad 0 \leq i \leq k,$$

where

$$(32') \quad Q(n, i) = \sum_{j=1}^N \delta(i, j) \delta(n, j)$$

are the elements of a symmetric matrix Q .

We consider the three cases

$$(') \quad k < N-1$$

$$('') \quad k = N-1 \quad (\text{number of coefficients } e_j \text{ equal to number of intervals } I_j).$$

$$('''') \quad k > N-1,$$

and assume Q to be non-singular for $k \leq N-1$. Of course, it will depend on the subdivision $\{e_j\}$ whether this condition is satisfied. Numerical computations showed that usually, but not always, Q is non-singular and therefore, because of (32'), positive definite. Of course, (32') is equivalent to

$$\sum_{n,i} Q(n,i) e_n e_i = \sum_j \left(\sum_i \delta(i,j) e_i \right)^2.$$

Analytical considerations also indicate that Q will usually be non-singular. Our procedure is a discrete analogue of method D of [5], but with the function system $\{W_n^*\}$, which is the orthogonal system of the continuous limit case as $N \rightarrow \infty$. Another argument is that the determinant of Q can be expressed explicitly in terms of elementary functions of α . If it does not vanish identically, it can only have isolated zeros.

There are two cases where it is quite evident that Q is singular.

(a) If there is an index λ such that $e_\lambda = 0$, then

$$e_{\lambda-1} = -e_{\lambda+1}, \quad \text{and by (27) and (28)}$$

$$G_p(e_{\lambda-1}, e_\lambda) = G_p(e_\lambda, e_{\lambda+1})$$

and

$$\delta(n, \lambda-1) = \delta(n, \lambda).$$

(b) If there is an index λ such that $(e_\lambda + e_{\lambda+1})/2 = 0$,

$$\text{then } e_\lambda = -e_{\lambda+1},$$

and

$$G_p(e_{\lambda-1}, e_\lambda) = G_p(e_{\lambda+1}, e_{\lambda+2}), \quad \delta(n, \lambda-1) = \delta(n, \lambda+1).$$

In both cases the matrix δ has two equal columns. Now, if the number of equal pairs of elements $\delta(n, j)$ is N_e (which is equal to the number of pairs of intervals symmetric with respect to $x=0$), then, if $k+1 > N - N_e$ the rank of W is less than $k+1$.

Using the theorem mentioned in paragraph 3 ([9], p.64-65), while computing the determinant $\det Q = \det (\delta \delta')$, we find that all the minors are zero, and therefore $\det Q = 0$.

These two cases may be formulated together by saying that when $\tilde{c} = -\alpha/h$ is equal to an integral multiple of $1/2$, and when the number of pairs of intervals symmetric with respect to $x=0$ is greater than $N - (k+1)$, then Q is singular.

One might think that this is a special disadvantage of approximating by a system of even orthogonal functions. But if one imagines the T_j/h as a step function to be fitted by a suitably smooth, even function, one sees, that information is lost if $x=0$ is equal to an endpoint or a midpoint of an interval. The reason is that $J(x)$ is supposed to be an even function, and disregarding the noise, one has less than N values T . If the intervals I_λ and I_μ lie symmetrically with respect to $x=0$, then $T_\lambda = T_\mu$.

Let us return to the three cases ('), (") and (").

Assuming Q to be regular, there is a unique solution in cases (') and ("). However, in case (") this solution is uninteresting, since it gives $D=0$. To show this, it is sufficient to see that if Q is regular, so also is the matrix δ , since δ is quadratic, and $Q = \delta \delta'$. Therefore there exists a unique solution to the system of linear equations

$$T_j - \sum_{n=0}^k \delta(n, j) c_n = 0, \quad j = 1, 2, \dots, k+1,$$

and by (31), $D = 0$.

In case (") Q is singular. Proof: The quadratic form

$$A = \sum_{\mu, \lambda=0}^k Q(\mu, \lambda) c_\mu c_\lambda$$

can be expressed as

$$A = \left(\sum_{\nu=0}^k \vec{\delta}_\nu c_\nu \right)^2$$

with the vectors $\vec{\delta}_i = (\delta(i, 1), \delta(i, 2), \dots, \delta(i, N))$.

The system $\sum_{v=0}^k \vec{\delta}_v c_v = 0$

does not have a unique solution. The $\vec{\delta}_v$ are N -dimensional vectors, and so more than N of them are always linearly dependent.

A double-precision FORTRAN-routine was written for case (') incidentally also applicable in case ("). Test runs show that in actual cases either the fit is not sufficiently good (if k is too small), or that $J(x)$ has large oscillations (if k is too large), even into negative values of $J(x)$, which is physically meaningless. In either case $i(r)$ may be significantly negative near $r=0$. It is possible, as also shown by test runs, to reproduce a polynomial $i(r)$ of degree k , given in advance, from which the values T_j are exactly calculated, but the method is very sensitive to superposition of noise on the T_j .

7. Reduction to a quadratic programming problem

We decided to take into account the information $i(r) \geq 0$. Then we have to find a vector

$$c' = (c_0, c_1, \dots, c_k)$$

(the prime denotes a row-vector, c denotes the corresponding column-vector), so that the quadratic form

$$(31) \quad D(c) = \sum_{j=1}^N \left\{ T_j - \sum_{n=0}^k \delta(n,j) c_n \right\}^2$$

is to be minimized under the restrictions

$$i(r_\lambda) = \sum_{n=0}^k c_n P_n^*(r_\lambda) \geq 0, \quad \lambda = 1, 2, \dots, m.$$

Here $k \leq N-1$, $\delta(n,j)$, $P_n^*(r_\lambda)$ are known constants, $P_n^*(r)$ is defined in § 2, formula (11). m , the number of restrictions, is chosen so as to give a sufficiently dense set of discrete values r_λ . After some trial runs we decided to take $k=5$ and $m=20$, $r_\lambda = (\lambda-1)/m$.

In order to transform the problem into the standard form of quadratic programming we compute

$$D(c) = \sum_{j=1}^N \left\{ \sum_{n=0}^k \delta(n,j) c_n \right\}^2 - 2 \sum_{n=0}^k \left\{ \sum_{j=1}^N \delta(n,j) T_j \right\} c_n + \sum_{j=1}^N T_j^2.$$

Furthermore, if we set

$$(32'') \quad Q_{ni} = Q_{in} = Q(n,i),$$

$$(33) \quad b_n = -2 \sum_{j=1}^N \delta(n,j) T_j, \quad A_{in} = -P_n^*(r_i),$$

we can state our problem in matrix form:-

By suitably choosing c

$$(34) \quad \begin{aligned} &\text{minimize } D(c) = c' Q c + b' c \\ &\text{subject to the restriction } A c \leq 0. \end{aligned}$$

As pointed out in § 6, we assume Q to be a positive definite $(k+1, k+1)$ matrix, ($\tilde{c} = -\alpha/h$ is not an integral multiple of $1/2$).

To solve this problem we used the iterative method of HILDRETH and D'ESOPPO ([7], 73-79). This method yields a sequence $c^{(j)}$ converging to the unique solution c if two conditions are fulfilled:

- (.) the matrix Q must be strictly positive definite,
- (..) there must exist a vector c such that

$$A_i c \leq -\varepsilon \quad \text{for } 1 \leq i \leq m \text{ and some } \varepsilon > 0.$$
 Here $A_i = (A_{i0}, A_{i1}, \dots, A_{ik})$.

(.) is assumed to be satisfied if $-\alpha/h \neq j/2$ for all integers j . (..) is satisfied if $c' = (-1, 0, 0, \dots, 0)$.

For then $A_i c = -\sum_{n=0}^k c_n P_n^*(r_i) = -P_0^*(r_i) = -1 < 0$.

As a stopping criterion for the double-precision FORTRAN-program we used:- "Stop iterating if

$$|c_i^{(p+1)} - c_i^{(p)}| / |c_i^{(p)}| < 10^{-4} \quad \text{and} \quad c_i^{(p)} \neq 0$$

for five values p in succession or after at most L iterations with $L = 200$."

As a measure of the quality of approximation the quantity

$$d_{rel} = \frac{\sqrt{D}}{\sum_{j=1}^N T_j}$$

was computed.

8. Numerical experience

We give statistics for the computing times of several runs. The program was organized so that in one computer run an arbitrary number of functions $i(r)$ could be computed, after reading the corresponding sets of data T_j

Number of functions $i(r)$ computed	running time in minutes	average running time for one function in minutes
56	43	0.77
86	60	0.70
59	63	1.07
66	56	0.85
56	70	1.27
123	120	0.98
446	412	0.92

The average number of iterations is another important statistic. For the 56 functions computed in 43 minutes the total number of iterations was 1754, and we have

1754/56 \approx 31.4 as an average number of iterations.

The quality of the approximation is illustrated by some sketches ^{*}), in which are plotted the T_j/h and the functions $J(x)$ and $i(r)$. If the measurements were accurate, in the intervals (e_j, e_{j+1}) the curve $y = J(x)$ should intersect the horizontals $y = T_j/h$. Of course, this could also be achieved with inaccurate T_j by choosing h sufficiently high, but then $J(x)$ and $i(r)$ would be physically meaningless, strongly oscillating functions. In case of measurement errors it is not wise to attempt too close an approximation. d_{rel} was usually about 10^{-2} .

On the other hand, it must be admitted that polynomial, least-squares approximations of sufficiently low degree yield functions which do not have too many extrema and inflection points, but at the same time are not quite satisfactory. This happened especially often in cases where $J(x)$ had a local minimum at $x = 0$. Sometimes $J(x)$ had a local maximum very near to $x = 1$, which was considered to be physically meaningless. The reason for it lies in the wild behaviour of the functions $W_n^*(x)$ at $x \approx 1$, which is illustrated in ([5], p. 19)

This computational effect was considered tolerable, because the more interesting region is near $x = 0$. Further details are mentioned in the comments under the sketches (see appendix).

We tried also to reconstruct an analytically given function $i(r)$ from the values T_j computed for a given subdivision (e_j, e_{j+1}) , $j = 1, 2, \dots, N$ of $-1 \leq x \leq 1$. We chose an even function $i(r)$, which has a deep local minimum at $r = 0$, and one maximum between 0 and 1.

^{*}) see appendix, p. 51

We took

$$i(r) = \frac{1}{2} (1 - r^2) (r^2 + \beta),$$

which is ≥ 0 in $0 \leq r \leq 1$, if $\beta \geq 0$.

If $0 \leq \beta < 1$, then $i(r)$ has a maximal value

$$i_{\max} = \frac{1}{8} (1 + \beta)^2 \quad \text{at } r = r_{\max} = \sqrt{(1 - \beta)/2}, \text{ and}$$

by varying β one can alter the ratio

$$\frac{i_{\max}}{i(0)} = \frac{\frac{1}{8} (1 + \beta)^2}{\beta/2}.$$

We choose $\beta = 0.1$, which yields $i(0) = 0.05$,

$$i_{\max} = 0.15125, \quad r_{\max} = \sqrt{0.45} \approx 0.67, \quad \frac{i_{\max}}{i(0)} = 3.025.$$

Expanding $i(r) = \frac{1}{2} \sum c_j P_j^*(r)$ yields $c_j = 0$ for $j \geq 3$,

$$c_0 \approx 0.22095, \quad c_1 \approx -0.002963, \quad c_2 \approx -0.013545,$$

from which $J(x)$ and the T_j can be computed.

The T_j were computed for $\alpha = -0.827$, $h = 0.23$, $N = 7$.

The lists, in appendix p. 50, give the exactly computed values of $J(x)$ and $i(r)$ (computed from the given coefficients c_j), and then the values of $J(x)$ and $i(r)$, computed from the values T_j . Subsequently the influence of inaccurate estimated α , and then the influence of disturbed T_j are investigated. The values of T_j are disturbed by adding normally distributed random numbers with mean 0 and standard deviation equal to

$$\frac{1}{10} \max_{j=1}^N T_j.$$

The authors would like to express the opinion, that studies in curve-fitting according to ideas proposed by HOOKE and JEEVES ([6], 222-224) should be undertaken. That is, one should have an alternative to polynomial approximation (or more generally: to approximation by special systems of smooth orthogonal functions) and instead, fit to given values \bar{y}_n (which are supposed to belong to a sufficiently dense set of equidistant x_n)

a discrete set of points (x_n, y_n) . In determining y_n , the number of maxima, minima and inflection points of the smooth curve $y = y(x)$, which we suppose to represent the proper physical process, should be used, in the form of linear difference inequality restrictions on the convexity and concavity of the fitted function. Under these restrictions a convenient deviation function

$D(y_1, \dots, y_N; \bar{y}_1, \dots, \bar{y}_N; x_1^*, \dots, x_M^*)$ is to be minimized, where the x_j^* are the special abscissae (of inflection points and extrema), and the y_j and x_λ^* are to be determined.

By this method it should be possible to avoid the previously mentioned disadvantages of approximation by a finite sum of special functions.

9. The case of an arbitrary radius R

Up to now we assumed the radius of the configuration to be equal to 1. For convenience we shall now give the essential transformation formulae for the case of an arbitrary, positive radius R . For this purpose let us assume that for the case $R=1$ we use the variables ξ and ρ , $-1 \leq \xi \leq 1$, $0 \leq \rho \leq 1$, and the intensities $J(\xi)$ and $i(\rho)$.

Then

$$(35) \quad T_j \approx \int_{e_j}^{e_{j+1}} J(\xi) d\xi \approx J(\bar{\xi}_j) h,$$

$$(36) \quad J(\xi) = 2 \int_{\rho=|\xi|}^1 \frac{i(\rho) \rho d\rho}{\sqrt{\rho^2 - \xi^2}}.$$

For the case of an arbitrary $R > 0$ we use the variables x and r , $-R \leq x \leq R$, $0 \leq r \leq R$, $J^*(x)$, $i^*(r)$.

a discrete set of points (x_n, y_n) . In determining y_n , the number of maxima, minima and inflection points of the smooth curve $y = y(x)$, which we suppose to represent the proper physical process, should be used, in the form of linear difference inequality restrictions on the convexity and concavity of the fitted function. Under these restrictions a convenient deviation function

$D(y_1, \dots, y_N; \bar{y}_1, \dots, \bar{y}_N; x_1^*, \dots, x_M^*)$ is to be minimized, where the x_j^* are the special abscissae (of inflection points and extrema), and the y_j and x_λ^* are to be determined.

By this method it should be possible to avoid the previously mentioned disadvantages of approximation by a finite sum of special functions.

9. The case of an arbitrary radius R

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Then

$$(35) \quad T_j \approx \int_{\xi_j}^{\xi_{j+1}} J(\xi) d\xi \approx J(\bar{\xi}_j) h,$$

$$(36) \quad J(\xi) = 2 \int_{\varrho=|\xi|}^1 \frac{i(\varrho) \varrho d\varrho}{\sqrt{\varrho^2 - \xi^2}}.$$

For the case of an arbitrary $R > 0$ we use the variables x and r , $-R \leq x \leq R$, $0 \leq r \leq R$, $J^*(x)$, $i^*(r)$.

Then

$$(35') \quad T_j \approx \int_{e_j^*}^{e_{j+1}^*} J^*(x) dx \approx J^*(\bar{x}_j) h^*,$$

$$(36') \quad J^*(x) = 2 \int_{r=|x|}^R \frac{i^*(r) r dr}{\sqrt{r^2 - x^2}}.$$

(The meaning of e_j^* , \bar{x}_j , \bar{x}_j , h^* is obvious).

$$\text{Then} \quad \xi = x/R, \quad \rho = r/R, \quad h = h^*/R.$$

By comparing (35) and (35') we obtain

$$(37) \quad J^*(x) = R^{-1} J(\xi).$$

From (36) we obtain, by changing the variables,

$$J(\xi) = 2 \int_{r=|x|}^R \frac{i(r/R) (r/R) (dr/R)}{R^{-1} \sqrt{r^2 - x^2}} = 2 \int_{r=|x|}^R \frac{R^{-1} i(r/R) r dr}{\sqrt{r^2 - x^2}}$$

which, by (37), $= R J^*(x)$.

Therefore (36') yields

$$(38) \quad i^*(r) = \frac{i(\rho)}{R^2} = \frac{i(r/R)}{R^2}.$$

Acknowledgements

The authors are indebted to a number of colleagues at the Institut für Plasmaphysik for helpful discussions, especially to Dr. A. Eberhagen, the evaluation of whose experimental results gave rise to the above work, and to F.M. Larkin from Culham Laboratory, who assisted us in styling the English text of this report.

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D. Appendixa) Description of the programs

We present here the listings of the two FORTRAN programs which implement the above methods. In order to use the programs they may have to be adjusted to suit individual requirements.

A short description of the programs is given and, together with the comment cards in the programs themselves, this should be enough to indicate how any necessary changes may be effected.

Because of limited programming time, program efficiency and elegance have, in places, been disregarded in the interests of simplicity. The programs are given as they have been run, together with an example of a printed output (of program 1) and some curves obtained from the output of program 2.

The programs consist of the following routines:

	program 1	program 2
standard type	INVERS	INVERS
	LGLSY	LGLSY
	BKOEFF	BKOEFF
	BETAKO	BETAKO
	GAMMAK	GAMMAK
similar type	MATCAL	MATCAL
	QUADRA	QUADRA
special type	MAIN routine	MAIN routine
	COMPUT	INOUT
		COMPUT
		GINDE
		DEFIN

Subroutines INVERS and LGLSY in the two programs differ in COMMON statements only.

MATCAL and QUADRA have the same function in both programs, but are a little different in each of them.

The routines BKOEFF, BETAKO, GAMMAK, which compute the coefficients $b(\rho, n)$, $\beta(\rho, n)$, $\gamma(\rho, n)$ are given in [5].

Description of program 1 - (continuous case)

The program uses a quadratic-programming method to approximate a continuous function with Gegenbauer polynomials of degrees from 4 to 7 successively.

In order to run the program it is necessary to give the following data deck:

Items on card	Format	explanation
NT, RGROSS	I10,E14.4	NT = Number of points including 0. RGROSS = scaling factor
T ₁ , ..., T ₆ T ₇ , ..., T _{NT}	6E12.4	T _i = The measured values, $i=1, \dots, NT$, of $J(x)$.

These data can be repeated as often as new measurements have been made.

Description of program 2 - (discrete case)

This program solves the above problem with polynomials of degree 5 only.

The results are written on tape as well as on off-line printer. The results on tape can then be used for further computations, or for graphical output. Channel 3 is used in this program.

The input data is divided into three classes:

- 1) Data for one machine run.
- 2) Data for a whole set of measurements, which consists of single measurements at given times.
(A measurement consists of N measured values which are to be fitted by one curve)
- 3) Data for a single measurement.

Class 1 consists of a single card, class 2 consists of a single card which is repeated every time a new set is started, and class 3 consists of 2 cards for every measurement. (see data sheet).

The program assumes that for each measurement there are 8 measured points, the first of which can be excluded in case it was wrongly measured. In that case, NST should be 7. (see data sheet). This limitation to 8 points can easily be removed by changing only the READ instructions in subroutine INOUT. In any case, NST cannot exceed the value of 20, because of DIMENSION statements. The program checks then if one of the intervals lies outside the segment $\langle -1, 1 \rangle$, and decides how many points will enter the computation (N8 - see glossary).

The comment cards in the various routines, in addition to the data sheet and the following glossary should be sufficient to explain the function and the use of each of them.

DATA - sheet

class	Items on card	Format	Explanation and comments
1	MISPAR	I 12	Mispar = running number of set of measurements on tape. It is equal to the number of sets which are already on tape.
2	NOB, HROH, DELAMB, EK, WELL, NST, NZAHL	I4,4F8.3, 2I3	NOB = Identification number of set. HROH = Original length of interval. DELAMB, EK, WELL = data irrelevant to this program, but which may be needed for further computations. NST = number of measured points. Usually NST = 8. If the first point was wrongly measured, NST = 7.
3	1)NOB,WALZT, $T_1, T_2, T_3, T_4,$ 2)NOB, T_5, \dots, T_8, FC	I4,5E12.4	WALZT = time T_i = the measured values ($i = 1, 2, \dots, 8$) FC = centre = \tilde{c} .

GLOSSARY

Routine	FORTTRAN Variable	Term in theoretical part	Explanation and comments
MAIN	N8	N	Number of measured values which are to enter computation (see NST in data sheet!)
	AL(I)	e_i	end points of computed intervals
	DELTA(J,N)	$\delta(n, j)$	$\delta_{n,j} = \sum_{p=0}^n \beta(p, n) \int_{I_j} x^{2p} \sqrt{1-x^2} dx$
	R(I)	r_λ	points in which the restrictions are computed ($r_\lambda = \lambda/m, \lambda = 0, 1, \dots, m-1$)
	P(I,N)	$A_{i,n}$	coefficients of restrictions
	E(I)	b_i	coefficients of the linear part of D
	QREC	Q	The positive definite matrix
	Q	Q^{-1}	
	V,G,U,W		Matrices and vectors needed for solving the quadratic-programming problem (see [7])
	NIT		number of iterations, not exceeding 200
	D,DREL	D,Drel	Absolute and relative deviation function

Routine	FORTTRAN Variable	Term in theoretical part	Explanation and comments
INOUT	RGROSS	R	Radius of Vessel in mm. This is a constant in the program.
COMPUT	F, RINT	$J(x), i(r)$	

SUBROUTINE MATCAL

SUBROUTINE MATCAL

```

C          COMPUTES MATRICES NECESSARY TO SOLVE QUADRATIC-PROGR.PROBLEM
D DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D DIMENSION QREC(11,11),G(11,11)
D DIMENSION RSL(11),FMATEL(11,11)
D DIMENSION P(20,20),E(20),T(40),X(40)
D DIMENSION WE(40,11)
COMMON B,BETA,NK,GAMMA,K
COMMON NT
COMMON QREC,G
COMMON RSL,FMATEL
COMMON P,E,T,X
COMMON WE
D DIMENSION R(20)
C CALCULATION OF WE.
D DO 3 J=1,NT
C DO 3 N=1,K
D WE(J,N)=0.
D DO 2 L=1,N
D 2 WE(J,N)=WE(J,N)+BETA(L,N)*X(J)**(2*L-2)
D 3 WE(J,N)=WE(J,N)*SQRTF(1.-X(J)**2)
D DO 4 I=1,20
C WW=I-1
C CALCULATION OF MATRIX P-RESTRICTIONS.
D 4 R(I)=WW*0.05
D DO 5 I=1,20
D DO 5 N=1,K
D P(I,N)=0.
D DO 5 L=1,N
D 5 P(I,N)=P(I,N)-GAMMA(L,N)*R(I)**(2*L-2)
C CALCULATION OF VECTOR E.
D DO 7 N=1,K
D E(N)=0.
D DO 6 J=1,NT
D 6 E(N)=E(N)-WE(J,N)*T(J)
D 7 E(N)=2.*E(N)
C CALCULATION OF MATRIX QREC.
D DO 8 M=1,K
D DO 8 L=1,K
D QREC(M,L)=0.
D DO 8 J=1,NT
D 8 QREC(M,L)=QREC(M,L)+WE(J,M)*WE(J,L)
RETURN
END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

SUBROUTINE COMPUT(C)

SUBROUTINE COMPUT(C)

COMPUTES AND PRINTS RESULTING FUNCTION AND INTENSITIES

```

DIMENSION S(40),RE(40),RE2(40)
DIMENSION WE(40,11)
DIMENSION V(20),PP(11),F(40),RINT(40)
DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
DIMENSION QREC(11,11),Q(11,11)
DIMENSION RSL(11),FMATEL(11,11)
DIMENSION P(20,20),E(20),T(40),X(40)
COMMON B,BETA,NK,GAMMA,K
COMMON NT
COMMON QREC,Q
COMMON RSL,FMATEL
COMMON P,E,T,X
COMMON WE,RGROSS
DIMENSION C(11)
WRITE OUTPUT TAPE 3,100,K
DO 52 J = 1,NT
  S(J) = 0.
DO 52 N = 1,K
52  S(J) = S(J) + C(N) * WE(J,N)
  D = 0.
DO 54 J = 1,NT
54  D = D + (T(J) - S(J))**2
WRITE OUTPUT TAPE 3,700,D
700 FORMAT( 6H D =E12.5)
WRITE OUTPUT TAPE 3,102
DO 5 I = 1,NT
  J = I - 1
DO 2 N = 1,K
  PP(N) = 0.
DO 2 L = 1,N
  EX = X(1)**(2*L - 2)
2  PP(N) = PP(N) + GAMMA(L,N) * EX
  F(I) = 0.
  RINT(I) = 0.
DO 4 N = 1,K
  F(I) = F(I) + C(N) * WE(I,N)
4  RINT(I) = RINT(I) + C(N) * PP(N)
  RINT(I) = RINT(I) / (2.*RGROSS)
  RE(I) = F(I) - T(I)
WRITE OUTPUT TAPE 3,101, J,X(I),T(I),F(I),RINT(I),RE(I)
5 CONTINUE
RETURN
100 FORMAT(1H1,10X,16HVALUES FOR K =I3////)
101 FORMAT(1H 14,F8.2,14X,5F16.4)
102 FORMAT(1H ,7X6HX OR R,23X4HT(X),10X9HCAPITAL I,9X7HSMALL I,9X10HDI
1FFERENCE//)
END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

Program 2

SOLVES THE APPROXIMATION PROBLEM BY CONVERTING IT

INTO A QUADRATIC-PROGRAMING PROBLEM
DISCRETE CASE.

```

C
C
D DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D DIMENSION T(20),ALPHA(1),H(1)
D DIMENSION QREC(11,11),Q(11,11)
D DIMENSION RSL(11),FMATEL(11,11),DELTG(1)
COMMON B,BETA,NK,GAMMA,K
COMMON T,ALPHA,N8
COMMON QREC,Q
COMMON RSL,FMATEL,DELTG
COMMON H,WALZT,NZAHL
COMMON MISPAR,RGROSS,FC,NINT
D DIMENSION P(20,20),E(11),DELTA(11,11)
COMMON P,E,DELTA
D DIMENSION CM(11,11),C(11),AL(21),ALI(21)
COMMON AL

READ INPUT TAPE 12,5111,MISPAR

REWIND 13
  NK = 11

D CALL BKOEFF
D CALL BETAKO
D CALL GAMMAK

  IF(MISPAR - 1) 8,509,509
509 DO 510 I = 1,MISPAR
510 CALL SKPFIL(13)
C READ AND WRITE DATA FOR A SET OF MEASUREMENTS.
  8 CALL INOUT(1)
  NL = 0
C READ AND WRITE DATA FOR SINGLE MEASUREMENTS.
  10 CALL INOUT(2)
  MI = 0
  NL = NL + 1
C TRANSFORM THE ORIGINAL INTERVAL INTO ONES INCLUDED
C IN (-1,1)
D 59 IF(ABSF(ALPHA)-1.) 1,1,2
D 1 AL(1) = ALPHA
  GO TO 5
D 2 IF(ALPHA) 4,1,3
D 3 AL(1) = 1.
  GO TO 5
D 4 AL(1) = -1.
D 5 DO 15 L = 1,N8
  J = L+1
D ALI(J) = FLOATF(L)*H + ALPHA
D IF(ABSF(ALI(J)) - 1.) 11,11,12
D 11 AL(J) = ALI(J)
  GO TO 15
D 12 IF(ALI(J)) 14,11,13
D 13 AL(J) = 1.

```


SOLVES THE APPROXIMATION PROBLEM BY CONVERTING IT

```

D 14 AL(J) = -1.
15 CONTINUE
    NNN = N8+1
WRITE OUTPUT TAPE 3,150,(AL(J),J=1,NNN)

    K = 6

C          COMPUTE MATRICES AND VECTORS
D  CALL MATCAL

WRITE OUTPUT TAPE 3,138,DETLG
C          CHECK QREC FOR SINGULARITY.
    DO 93 I = 1,K
      DO 93 J = 1,K
D  CM(I,J) = 0.
      DO 93 L = 1,K
D 93 CM(I,J) = CM(I,J) + QREC(I,L)*Q(L,J)

C          IS LOG OF DETERMINANT TOO SMALL.
    IF(DETLG + 9.)87,87,88
C          IT IS.CHANGE THE CENTRE.
87  MI = MI + 1
D  FMI = MI
    FC = FC + (-1.)**MI*FMI*0.05
CALL INOUT(3)
GO TO 59

C          PRINT THE UNIT MATRIX QREC*QREC**(-1)
88 DO 33 I = 1,K
33 WRITE OUTPUT TAPE 3,161,(CM(I,J),J=1,K)

C          SOLVE THE QUADRATIC-PROGRAMING PROBLEM
D  CALL QUADRA(C)
C          COMPUTE FUNCTION VALUES AND INTENSITIES
    CALL COMPUT(C)
402 IF(NL - NZAHL)10,7,7
    7 ENDFILE 13
    GO TO 8
1000 FORMAT(1HJ,3(E20.5))
139 FORMAT(1HJ,(8(E12.4,4X)))
190 FORMAT(I6,F4.1)
150 FORMAT(18HJTHE INTERVALS ARE/1HJ,9E13.4////)
160 FORMAT(28HLCHECKING THE INVERSE MATRIX//)
161 FORMAT(1H 8E20.5)
170 FORMAT(12HLO U T P U T//)
137 FORMAT(32H SINGULAR MATRIX, NO COMPUTATION////////)
138 FORMAT(21H LOG10 OF DETERMINANT E12.3)
5111 FORMAT(I12)
    END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

SUBROUTINE INOUT(I)

SUBROUTINE INOUT(I)

READS AND WRITES DATA

DEFINES ALPHA AND N8.

```

C      DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
C      DIMENSION T(20),ALPHA(1),H(1)
D      DIMENSION QREC(11,11),Q(11,11)
D      DIMENSION RSL(11),FMATEL(11,11),DELTA(1)
D      COMMON B,BETA,NK,GAMMA,K
      COMMON T,ALPHA,N8
      COMMON QREC,Q
      COMMON RSL,FMATEL,DELTA
      COMMON H,WALZT,NZAHL
      COMMON MISPAR,RGROSS,FC,NINT
      GO TO (7,8,20),I
7  READ INPUT TAPE 12,103,NOB,HROH,DELAMB,EK,WELL,NST,NZAHL
   WRITE OUTPUT TAPE 3,101
   WRITE OUTPUT TAPE 3,1000,NOB,DELAMB,EK,WELL,NZAHL,NST
   NINT = 20
   RGROSS = 23.
10  MISPAR=MISPAR+1
   WRITE OUTPUT TAPE 3,106,MISPAR
106 FORMAT(1HL,15,16H SCHUSSE ON TAPE/)
   WRITE TAPE 13,NOB,WELL,NZAHL,NINT
   GO TO 25
8  IF(NST-8) 45,46,46
45  READ INPUT TAPE 12,105,NOB,WALZT,DUMMY,(T(I),I=1,3),NOB,(T(I),I=
14,7),FC
   GO TO 19
46  READ INPUT TAPE 12,105,NOB,WALZT,(T(I),I=1,4),NOB,(T(I),I=5,8),FC
19  IF(HROH - 4.5) 11,11,12
11  IF(FC - 4.) 13,13,14
13  FNST = NST
   RGROSS = (FNST - FC) * HROH
   GO TO 12
14  RGROSS = FC * HROH
12  H = HROH/RGROSS
   WRITE OUTPUT TAPE 3,102,HROH,RGROSS,H
20  ALPHA = -FC*H
   ALPHA = ALPHA
D   WRITE OUTPUT TAPE 3,104,FC,ALPHA
   ST = NST
   E1 = ALPHA + (ST - 0.5) * H
   IF(E1 - 1.) 22,22,23
22  N8 = NST
   GO TO 24
23  N8 = NST - 1
24  WRITE OUTPUT TAPE 3,130,WALZT,N8
   WRITE OUTPUT TAPE 3,82
   WRITE OUTPUT TAPE 3,139,(T(L),L=1,N8)
25  RETURN
100 FORMAT(19HINPUT DATA/)
102 FORMAT(1HL,20X,6HHROH =E12.4/1H ,18X,3H RGROSS =E12.4/1H 23X,3HH =E
112.4)
105 FORMAT(14,5E12.4)
103 FORMAT(14,4F8.3,2I3)
139 FORMAT(1HJ,(8(E12.4,4X)))

```

SUBROUTINE INOUT(I)

```
82 FORMAT(1HJ,24H THE MEASURED VALUES ARE)
104 FORMAT(1H ,22X,4HFC =E12.4/1H ,19X,7HALPHA =E12.4///)
130 FORMAT(1H ,30X,5HTIME E12.4,3X,9HMICROSEC.,10X,12,8H LIGHTS/)
1000 FORMAT(1HJ,20X,10HSCHUSS NR.15//1H ,18X,8HDELAMB = E12.4/1H ,22X,
14HEK =E12.4/1H ,20X,6HWELL = E12.4//1H ,22X,14,11H ZEITPUNKTE/
21H ,22X,14,16H ORIGINAL LIGHTS//)
END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)
```

SUBROUTINE COMPUT (C)

SUBROUTINE COMPUT (C)

COMPUTES FUNCTION VALUES AND INTENSITIES

C
D DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D DIMENSION T(23),ALPHA(1),H(1)
D DIMENSION QREC(11,11),Q(11,11)
D DIMENSION RSL(11),FMATEL(11,11),DELTA(1)
COMMON B,BETA,NK,GAMMA,K
COMMON T,ALPHA,NB
COMMON QREC,Q
COMMON RSL,FMATEL,DELTA
COMMON H,WALZT,NZAHL
COMMON MISPAR,RGROSS,FC,NINT
DIMENSION AL(21)
D DIMENSION P(11),C(11),V(11),X(101),F(101),RINT(101),XX(21)
FNINT = NINT
R = RGROSS / 10.
NNN = NB + 1
HROH = H * R

TRANSFORM BACK TO ORIGINAL INTERVALS

C
AL(1) = ALPHA * R
DO 6 I = 2,NNN
6 AL(I) = AL(I-1) + HROH
WRITE OUTPUT TAPE 3,102
WRITE OUTPUT TAPE 3,103,(AL(I),I=1,NNN)
DO 5 I = 1,NINT
D X(I) = FLOAT(I-1)/FNINT
D XX(I) = X(I) * R
DO 2 N = 1,K
D V(N) = 0.
D 1 P(N) = 0.
DO 2 L = 1,N
D V(N) = V(N) + BETA(L,N) * X(I)**(2*L-2)
D 2 P(N) = P(N) + GAMMA(L,N) * X(I)**(2*L-2)
D F(I) = 0.
D 3 RINT(I) = 0.
DO 4 N = 1,K

D F(I) = F(I) + V(N) * C(N)
D 4 RINT(I) = RINT(I) + C(N) * P(N)

TRANSFORM BACK TO ORIGINAL VALUES OF FUNCTION AND INTENSITIES

D RINT(I) = RINT(I) * H / (2. * R)
D F(I) = F(I) * SQRTF(1. - X(I) ** 2) * H

5 CONTINUE

WRITE OUTPUT TAPE 3,100
WRITE OUTPUT TAPE 3,101,(XX(I),F(I),RINT(I),I = 1,NINT)
WRITE TAPE 13,WALZT, NB, (XX(I),RINT(I),I=1,NINT)
WRITE TAPE 13,(AL(I),I=1,NNN),(T(I),I=1,NB)
WRITE TAPE 13,(F(I),I = 1,NINT)
RETURN
100 FORMAT(1HJ 2H X,16X,1HF,18X,4HRINT//)
101 FORMAT(1H ,F4.2,E25.8,E25.6)
102 FORMAT(18HJTHE INTERVALS ARE//)
103 FORMAT(1H ,9E14.4/1H ,8E15.4)
END(1,0,0,0,0,0,1,0,0,1,0,0,0,0)

SUBROUTINE MATCAL

```

SUBROUTINE MATCAL
D   DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D   DIMENSION I(20),ALPHA(1),H(1)
D   DIMENSION QREC(11,11),Q(11,11)
D   DIMENSION RSL(11),FMATEL(11,11),DELTA(1)
D   COMMON B,BETA,NK,GAMMA,K
D   COMMON T,ALPHA,N8
D   COMMON QREC,Q
D   COMMON RSL,FMATEL,DELTA
D   COMMON H,WALZT,NZAHN
D   COMMON MISPAR,RGROSS,FC,NINT
D   DIMENSION P(20,20),E(11),DELTA(11,11)
D   COMMON P,E,DELTA
D   DIMENSION AL(21)
D   COMMON AL
D   DIMENSION R(20)
D   DO 31 J=1,N8
D   DO 31 N = 1,K
D   DELTA(J,N) = 0.
D   DO 31 L = 1,N
D   J = J
D   F = AL(J)
D   B = AL(J+1)
D   L = L
D   31 DELTA(J,N) = DELTA(J,N) + BETA(L,N)*DEFIN(L,F,B)

D   60 DO 61 I=1,20
D   WW = I-1
D   RW = 0.05
D   61 R(I) = WW*RW
D   DO 63 I=1,20
D   DO 63 N=1,K
D   62 P(I,N)=0.
D   DO 63 L=1,N
D   63 P(I,N) = P(I,N) - GAMMA(L,N)*R(I)**(2*L-2)
D   DO 65 N=1,K
D   64 E(N) = 0.
D   DO 65 J=1,N8
D   65 E(N)=E(N)-(DELTA(J,N)*I(J))*2.

D   DO 94 M=1,K
D   DO 94 L = 1,K
D   QREC(M,L) = 0.
D   DO 94 J = 1,N8
D   94 QREC(M,L) = QREC(M,L) + DELTA(J,K) * DELTA(J,L)

C   CALCULATE QREC**=-1.
D   CALL INVERS

D   RETURN
D   END(1, , , , ,1, , ,1, , , , , )

```

FUNCTION GINDE(L,X)

FUNCTION GINDE(L,X)

```

C          COMPUTES THE INDEFINITE INTEGRALS
D  DIMENSION G(11)
D      Z = 1. - X**2
D      IF(Z) 2,1,2
D  2      Y = X / SQRTF(Z)
D      IF DIVIDE CHECK 1,6
D  6      IF QUOTIENT OVERFLOW 1 , 3
D  1      IF(X) 7,7,10
D  7      Y = -1.E38
D      GO TO 3
D 10      Y = 1. E 38
D  3      G(1) = 0.5*(ATANF(Y)+X*SQRTF(Z))
D      IF(L-1)4,4,5
D  4      GINDE=G(1)
D      GO TO9
D  5      DO 8 I=2,L
D          I = I
D      W = FLOATF(I)
D  8      G(I)=1./(2.*W)*(-(X**(2*I-3)*Z*SQRTF(Z))+(2.*W-3.)*G(I-1))
D      GINDE = G(L)
D  9      RETURN
D      END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

FUNCTION DEFIN(L,A,B)

FUNCTION DEFIN(L,A,B)

```

C          COMPUTES THE DEFINITE INTEGRALS
D  1      G1 = GINDE(L,A)
D  2      G2 = GINDE(L,B)
D      DEFIN = G2-G1
D      RETURN
D      END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

General routines

SUBROUTINE QUADRA(C)

SUBROUTINE QUADRA(C)

```

C          SOLVES THE QUADRATIC-PROGRAMMING PROBLEM
D  DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D  DIMENSION QREC(11,11),Q(11,11)
D  DIMENSION RSL(11),FMATEL(11,11)
D  DIMENSION P(20,20),E(20),T(40),X(40)
D  DIMENSION WE(40,11)
COMMON B,BETA,NK,GAMMA,K
COMMON NT
COMMON QREC,Q
COMMON RSL,FMATEL
COMMON P,E,T,X
COMMON WE
D  DIMENSION C(11),CC(11)
D  DIMENSION U(20),UPSI(20),R(20),V(20),G(20,20),W(20)
D  DIMENSION A(20),Z(20)
D  DO 500 I = 1,K
D 500 UPSI(I) = 0.
D  DO 68 I = 1,20
D 85 V(I) = 0.
D  DO 68 M = 1,K
D  DO 68 N = 1,K
D 68 V(I) = V(I) + P(I,N) * Q(N,M) * E(M) * 0.5
D  DO 70 I = 1,20
D  DO 70 J = 1,20
D  G(I,J)=0.
D  DO 70 L = 1,K
D  DO 70 N = 1,K
D 70 G(I,J) = G(I,J) + 0.25 * P(I,N) * Q(N,L) * P(J,L)
D  DO 71 I=1,20
D 71 U(I) = 0.
D  NIT = 0.
D 76 DO 6 I = 1,20
D 6 W(I) = 0.
D  DO 58 J = 2,20
D 58 W(1) = G(1,J)*U(J) + W(1)
D  W(1) = -1./G(1,1)*(W(1) + V(1)*0.5)
D  IF DIVIDE CHECK 1,2
D 2 IF QUOTIENT OVERFLOW 1,3
D 1 W(1) = -1.E37 * (W(1) + V(1) * 0.5)
D 3 IF(W(1))66,66,67
D 66 U(1) = 0.
D  GO TO 55
D 67 U(1) = W(1)
D 55 DO 73 I = 2,20
D  A(I) = 0.
D  Z(I) = 0.
D  L = I-1
D  DO 56 J = 1,L
D 56 A(I) = G(I,J)*U(J) + A(I)
D  M = I + 1
D  DO 57 J = M,20
D 57 Z(I) = G(I,J)*U(J) + Z(I)

```

SUBROUTINE QUADRA(C)

```

D      W(I) = -1./G(I,I)*(A(I) + Z(I) + 0.5*V(I))
      IF DIVIDE CHECK 4,5
      5 IF QUOTIENT OVERFLOW 4,10
D      4 W(I) = -1.E37 * (A(I) + Z(I) + 0.5*V(I))
D      10 IF(W(I)) 91,91,74
D      91 U(I) = 0.
      GO TO 73
D      74 U(I) = W(I)
      73 CONTINUE

      DO 40 J = 1,10
D      40 W(20) = W(20) + G(20,J) * U(J)
D      W(20) = -1./G(20,20)*(W(20) + V(20)*0.5)
      IF DIVIDE CHECK 7,8
      8 IF QUOTIENT OVERFLOW 7,9
D      7 W(20) = -1.E37 * (W(20) + V(20)*0.5)
D      9 IF(W(20)) 41,41,42
D      41 U(20) = 0.
      GO TO 43
D      42 U(20) = W(20)
      43 DO 25 L=1,K
D      CC(L) = 0.
      DO 25 N = 1,20
D      25 CC(L) = CC(L)+P(N,L)*U(N)
      DO 84 I = 1,K
D      C(I) = 0.
      DO 84 L = 1,K
D      84 C(I) = C(I)-Q(I,L)*(CC(L)+E(L))*0.5
      NIT = NIT + 1
      DO 501 I = 1,K
D      UPI = ABSF(UPSI(I) - C(I))
D      UP = UPI /ABSF( UPSI(I))
      IF DIVIDE CHECK 1001,1002
      1002 IF QUOTIENT OVERFLOW 503,1003
      1001 IF(UPI - 0.0001) 501,501,503
      1003 IF(UP - 0.0001) 501,501,503
      501 CONTINUE
      GO TO 90
      503 DO 502 I = 1,K
D      502 UPSI(I) = C(I)
      IF(NIT - 200) 76,89,89
      90 IF(NN-4) 83,83,89
      83 NN = NN + 1
      DO 504 I = 1,K
D      504 UPSI(I) = C(I)
      IF(NIT - 200) 76,89,89
      89 WRITE OUTPUT TAPE 3,133
      80 WRITE OUTPUT TAPE 3,145,NIT
      133 FORMAT(33HLNUMBER OF ITERATIONS EXCEEDS 200/150 NO CONVERGENCE)
      104 FORMAT(1H ,6E20.4)
      145 FORMAT(1HL,15,10HITERATIONS/)
      RETURN
      END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```


SUBROUTINE LGLSY

SUBROUTINE LGLSY

SOLVES A SYSTEM OF LINEAR EQUATIONS.

PROGRAMMED BY W.LUENOW.

```

C
C
C
D   DIMENSION DETLG(1),FMATEL(11,11),RSL(11)
D   DIMENSION T(20),ALPHA(1)
D   DIMENSION DUM(11,11),DUM1(11,11)
D   DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D   DIMENSION LIST(11)
COMMON B,BETA,NK,GAMMA,JM
COMMON T,ALPHA,N8
COMMON DUM,DUM1
COMMON RSL,FMATEL
COMMON DETLG
DO 1 J = 1,JM
1 LIST(J) = 0
C
C
D   DO 2 J=1,JM
D   GE = 0.0
C
C
D   DO 3 K=1,JM
C
D   DO 5 L=1,JM
D   IF(K-LIST(L)) 5,3,5
5 CONTINUE
C
D   IF(ABSF(FMATEL(J,K))-GE) 3,4,4
D   4 GE = ABSF(FMATEL(J,K))
D   LIST(J) = K
3 CONTINUE
C
C
D   IF(GE-1.0E-35) 1000,1001,1001
1000 WRITE OUTPUT TAPE 3, 1100
1100 FORMAT (49H1ABSF DES GROESSTEN ELEMENTES KLEINER ALS 1.0E-35)
CALL EXIT
C
C
D1001 IF(J-1) 1002,1002,1003
D1002 DETLG = 0.43429*LOGF(GE)
GO TO 10
C
D1003 DETLG = DETLG + 0.43429*LOGF(GE)
C
C
10 M = LIST(J)
JP = J+1
IF(JP-JM) 11,11,800
C
11 DO 12 K=JP,JM
D   FA = FMATEL(K,M)/FMATEL(J,M)
C
IF DIVIDE CHECK 5101,5150
5150 IF QUOTIENT OVERFLOW 5101,5100

```

SUBROUTINE LGLSY

```

D5101 IF(FMATEL(K,M)) 5160,5160,5161
D5160 B1 = -1.
      GO TO 5170
D5161 B1 = +1.
C
D5170 IF(FMATEL(J,M)) 5180,5180,5181
D5180 B2 = -1.
      GO TO 5190
D5181 B2 = +1.
C
D5190 FA = 1.0E+37*B1*B2
C
      5100 DO 13 L=1,JM
C
          DO 14 L2=1,JM
            IF(L-LIST(L2)) 14,13,14
          14 CONTINUE
C
D      FMATEL(K,L) = FMATEL(K,L) - FA*FMATEL(J,L)
      13 CONTINUE
C
D      RSL(K) = RSL(K) - FA*RSL(J)
      12 CONTINUE
      2 CONTINUE
C
C
      800 DO 15 J=1,JM
          JR = JM-J+1
          M = LIST(JR)
D      FA = RSL(JR)
          KL = JR+1
          IF(KL-JM) 32,32,16
C
          32 DO 20 K=KL,JM
              M1 = LIST(K)
D      20 FA = FA - FMATEL(JR,M1)*RSL(K)
C
D      16 RSL(JR) = FA/FMATEL(JR,M)
C
          IF DIVIDE CHECK 5200,5250
          5250 IF QUOTIENT OVERFLOW 5200,15
D5200 IF(FA) 5201,5201,5202
D5201 B1 = -1.
      GO TO 5210
D5202 B1 = +1.
C
D5210 IF(FMATEL(JR,M)) 5220,5220,5221
D5220 B2 = -1.
      GO TO 5230
D5221 B2 = +1.
C
D5230 RSL(JR) = 1.0E+37*B1*B2
C
      15 CONTINUE
C
C

```

SUBROUTINE LGLSY

```

      DO 50 J=1,JM
34  IF(J-LIST(J)) 33,50,33
33  M = LIST(J)
D   FA = RSL(M)
D   RSL(M) = RSL(J)
D   RSL(J) = FA
      JA = LIST(M)
      LIST(M) = LIST(J)
      LIST(J) = JA
      GO TO 34
50  CONTINUE
C
      RETURN
      END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

SUBROUTINE INVERS

SUBROUTINE INVERS

```

C           COMPUTES THE INVERSE OF A GIVEN MATRIX
D   DIMENSION C(11,11),CM1(11,11)
D   DIMENSION RSL(11),FMATEL(11,11)
D   DIMENSION I(20),ALPHA(1)
D   DIMENSION B(11,11),BETA(11,11),GAMMA(11,11)
D   COMMON B,BETA,NK,GAMMA,K
D   COMMON I,ALPHA,N8
D   COMMON C,CM1
D   COMMON RSL,FMATEL
D   DO 5 N = 1,K
D   DO 1 I = 1,K
D   DO 1 J = 1,K
D   1 FMATEL(I,J) = C(I,J)
      DO 2 I = 1,K
      IF(N-I) 4,3,4
D   3 RSL(I) = 1.
      GO TO 2
D   4 RSL(I) = 0.
      2 CONTINUE
D   CALL LGLSY
D   DO 5 I = 1,K
D   CM1(I,N) = RSL(I)
D   5 CONTINUE
D   RETURN
      END(1,0,0,0,0,0,1,0,0,1,0,0,0,0,0)

```

c) Example of an Output

VALUES FOR K = 5

D = 0.86775E 02 X OR R		T(X)	CAPITAL I	SMALL I	DIFFERENCE
0	-0.	29.8000	29.5995	-0.	-0.2005
1	0.04	29.9000	29.7110	0.0046	-0.1890
2	0.08	30.0000	30.0469	0.0190	0.0469
3	0.12	30.4000	30.6095	0.0444	0.2095
4	0.16	30.7000	31.3993	0.0827	0.6993
5	0.20	31.3000	32.4091	0.1360	1.1091
6	0.24	32.2000	33.6182	0.2065	1.4182
7	0.28	33.2000	34.9852	0.2957	1.7852
8	0.32	34.7000	36.4415	0.4040	1.7415
9	0.36	36.8000	37.8872	0.5301	1.0872
10	0.40	39.5000	39.1888	0.6708	-0.3112
11	0.44	42.4000	40.1818	0.8203	-2.2182
12	0.48	44.8000	40.6785	0.9702	-4.1215
13	0.52	44.7000	40.4817	1.1095	-4.2183
14	0.56	40.7000	39.4054	1.2254	-1.2946
15	0.60	36.4000	37.3031	1.3032	0.9031
16	0.64	32.0000	34.1012	1.3288	2.1012
17	0.68	27.5000	29.8379	1.2897	2.3379
18	0.72	23.0000	24.7020	1.1783	1.7020
19	0.76	18.5000	19.0657	0.9948	0.5657
20	0.80	13.8000	13.4993	0.7525	-0.3007
21	0.84	9.7000	8.7475	0.4827	-0.9525
22	0.88	6.3000	5.6253	0.2420	-0.6747
23	0.92	3.6000	4.7355	0.1201	1.1355
24	0.96	1.5000	5.6704	0.2502	4.1704
25	1.00	0.	-0.	0.8199	-0.

d) Reconstruction of a given Function

F AND I ARE THE GIVEN FUNCTIONS
 FS AND IS ARE THE RECONSTRUCTED FUNCTIONS
 F1 AND I1 ARE THE FUNCTIONS COMPUTED AFTER THE CENTRE HAS BEEN SHIFTED BY 0.013 TO THE RIGHT
 F2 AND I2 ARE THE FUNCTIONS COMPUTED FROM DISTURBED VALUES I(J)

X	F(X)	FS(X)	F1(X)	F2(X)
0.05	0.20000000	0.20000000	0.20015321	0.20722602
0.10	0.20057879	0.20057879	0.20087960	0.20418192
0.15	0.20226105	0.20226105	0.20288976	0.19604078
0.20	0.20488539	0.20488539	0.20573109	0.18587942
0.25	0.20818703	0.20818703	0.20891410	0.1775276
0.30	0.21180376	0.21180376	0.21160392	0.17562102
0.35	0.21528500	0.21528500	0.21378866	0.19203066
0.40	0.21810342	0.21810342	0.21535192	0.19704839
0.45	0.21967035	0.21967035	0.21648601	0.21772173
0.50	0.21935460	0.21935461	0.21732882	0.23832522
0.55	0.21650635	0.21650635	0.21755257	0.25151677
0.60	0.21048655	0.21048654	0.21592062	0.25033666
0.65	0.20070399	0.20070399	0.21000920	0.23073363
0.70	0.18666264	0.18666264	0.19637251	0.19401580
0.75	0.16802353	0.16802353	0.17146560	0.14834788
0.80	0.14468952	0.14468955	0.13359335	0.10816022
0.85	0.11692800	0.11692799	0.08590218	0.09031732
0.90	0.08556531	0.08556516	0.03964715	0.16042350
0.95	0.05234165	0.05234121	0.01531593	0.14831163
1.00	0.02074276	0.02074187	0.03271820	0.17471065

R	I(R)	IS(R)	I1(R)	I2(R)
0.05	0.04999999	0.05000003	0.04948173	0.15745167
0.10	0.05112187	0.05112191	0.05135211	0.14777616
0.15	0.05444999	0.05445004	0.05660973	0.12100688
0.20	0.05987187	0.05987190	0.06429449	0.03347224
0.25	0.06719999	0.06720003	0.07311862	0.04430833
0.30	0.07817188	0.07817190	0.08186080	0.01356678
0.35	0.08644999	0.08645003	0.08977760	-0.00000419
0.40	0.09762187	0.09762192	0.09691401	0.00867304
0.45	0.10919999	0.10920005	0.10419234	0.04035231
0.50	0.12062187	0.12062193	0.11317902	0.08864064
0.55	0.13124999	0.13125005	0.12547867	0.14253467
0.60	0.14037187	0.14037193	0.14178772	0.18735798
0.65	0.14720000	0.14720005	0.16076186	0.20839927
0.70	0.15087187	0.15087193	0.17801835	0.13549729
0.75	0.15045000	0.15045010	0.18580905	0.14803363
0.80	0.14492187	0.14492209	0.17417128	0.07909641
0.85	0.13320000	0.13320035	0.13469267	0.01677017
0.90	0.11412187	0.11412232	0.06842185	-0.
0.95	0.08644999	0.08645026	-0.00007860	0.05516473
1.00	0.04887187	0.04887119	0.	0.21867814

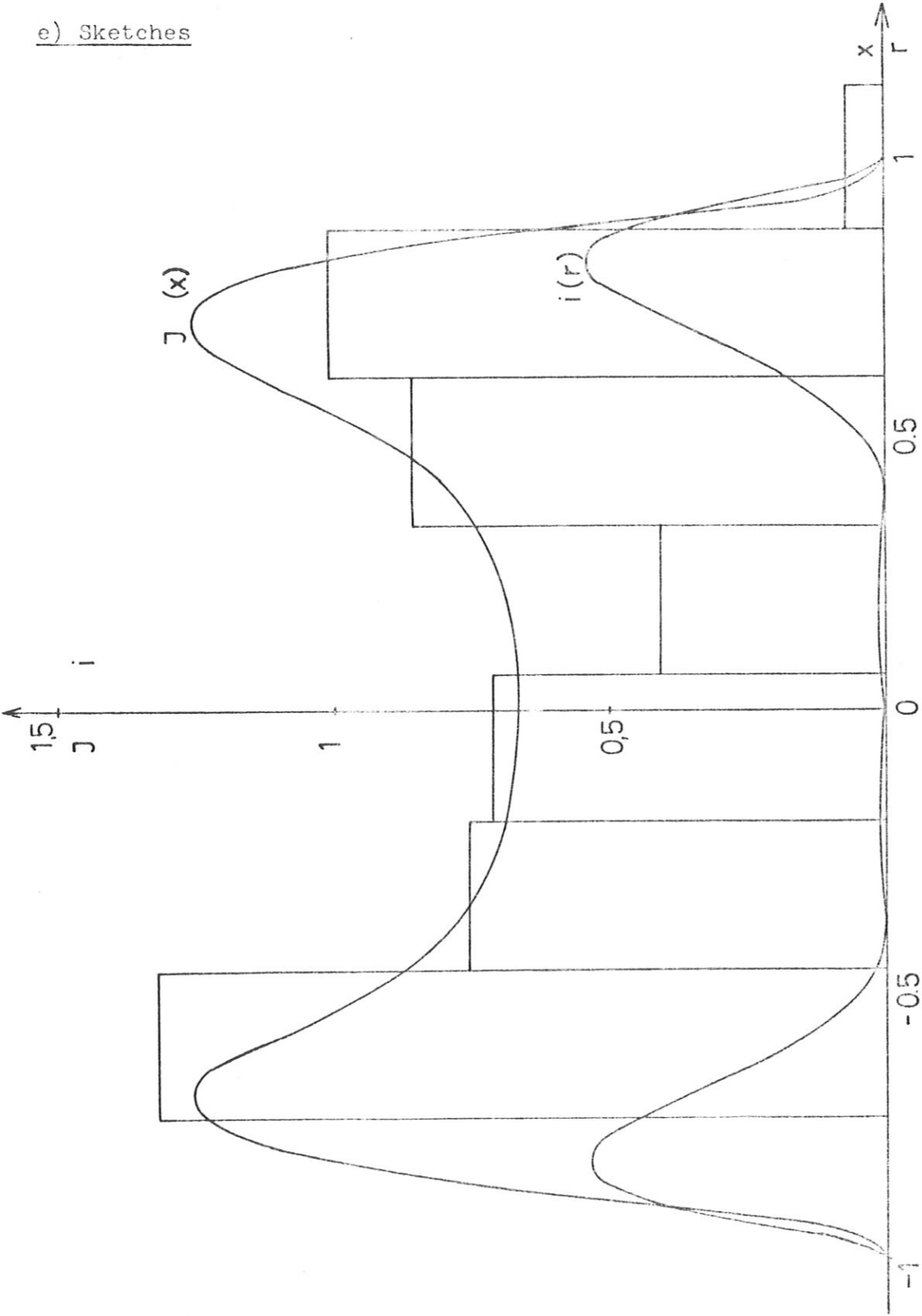


Fig.: 8

$J(x)$ = the fit to the values T_j/h of the step-function.
 $x = 5$ as in most of the cases which have been measured.

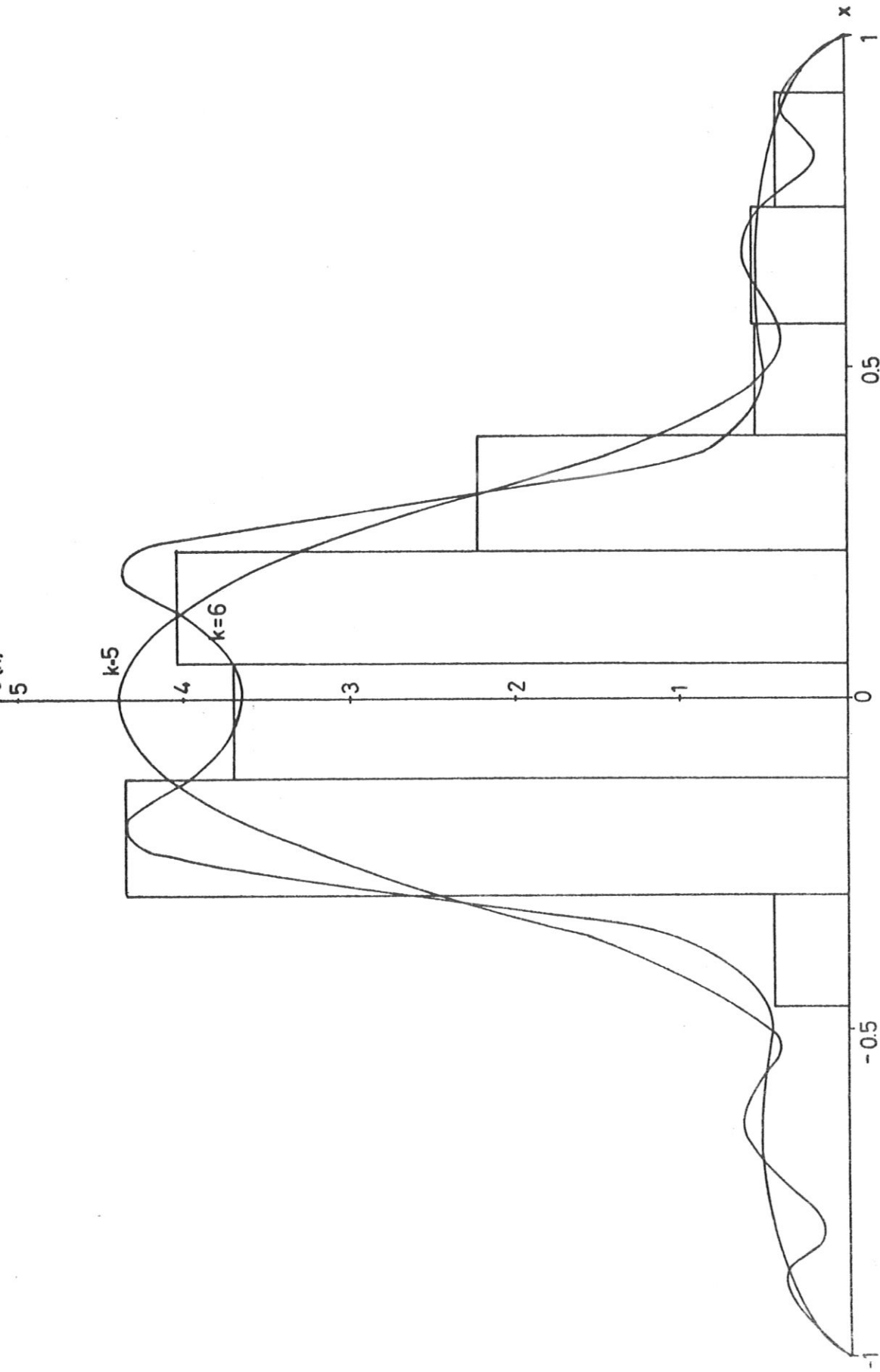


Fig.: 9

$k = 5$ is not large enough in order to approximate $J(x)$; $k = 6$ is much better in this case.

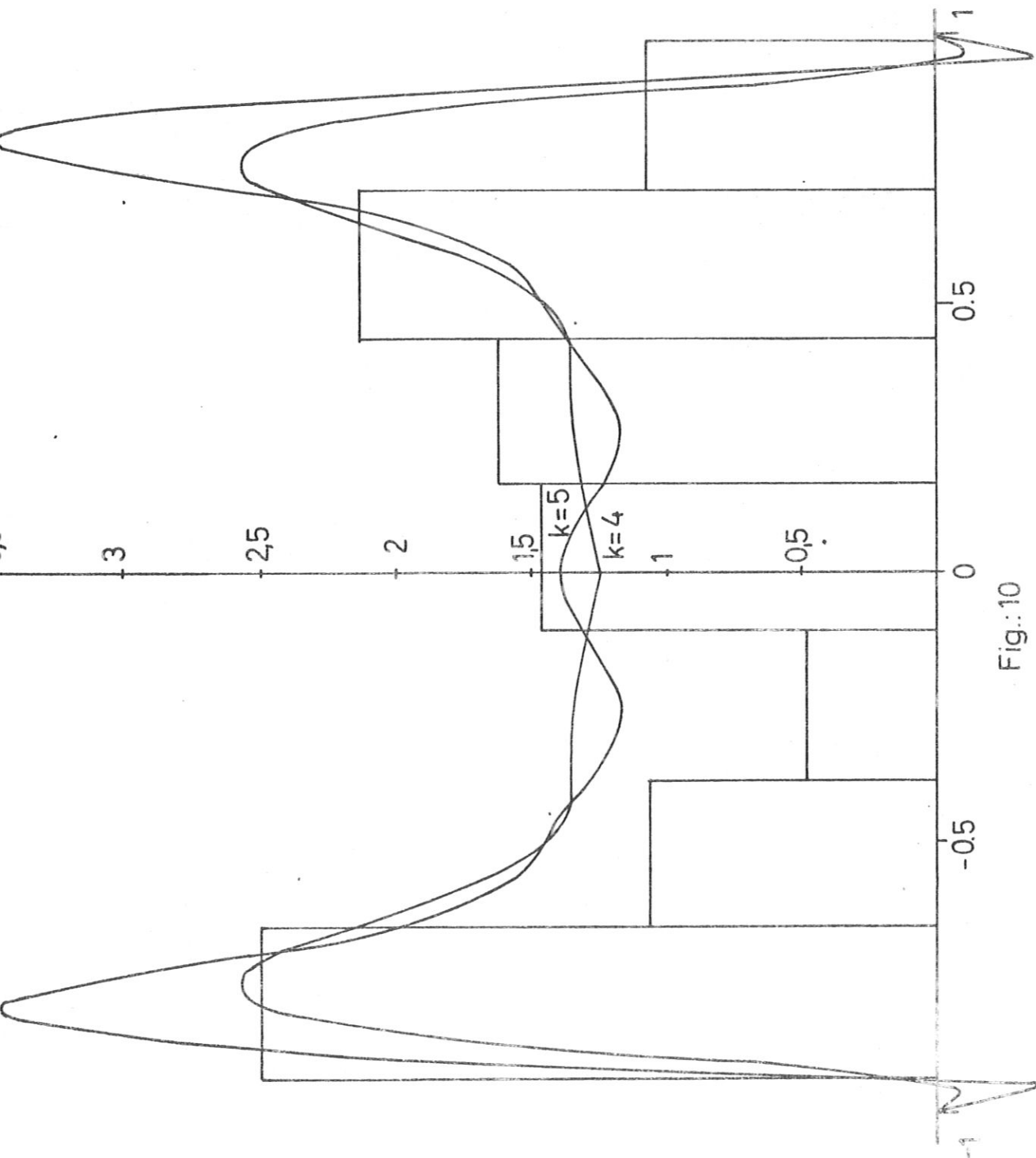


Fig.:10

The fit with $k = 4$ is better than with $k = 5$. Negative values are in the neighbourhood of $|x| = 1$.

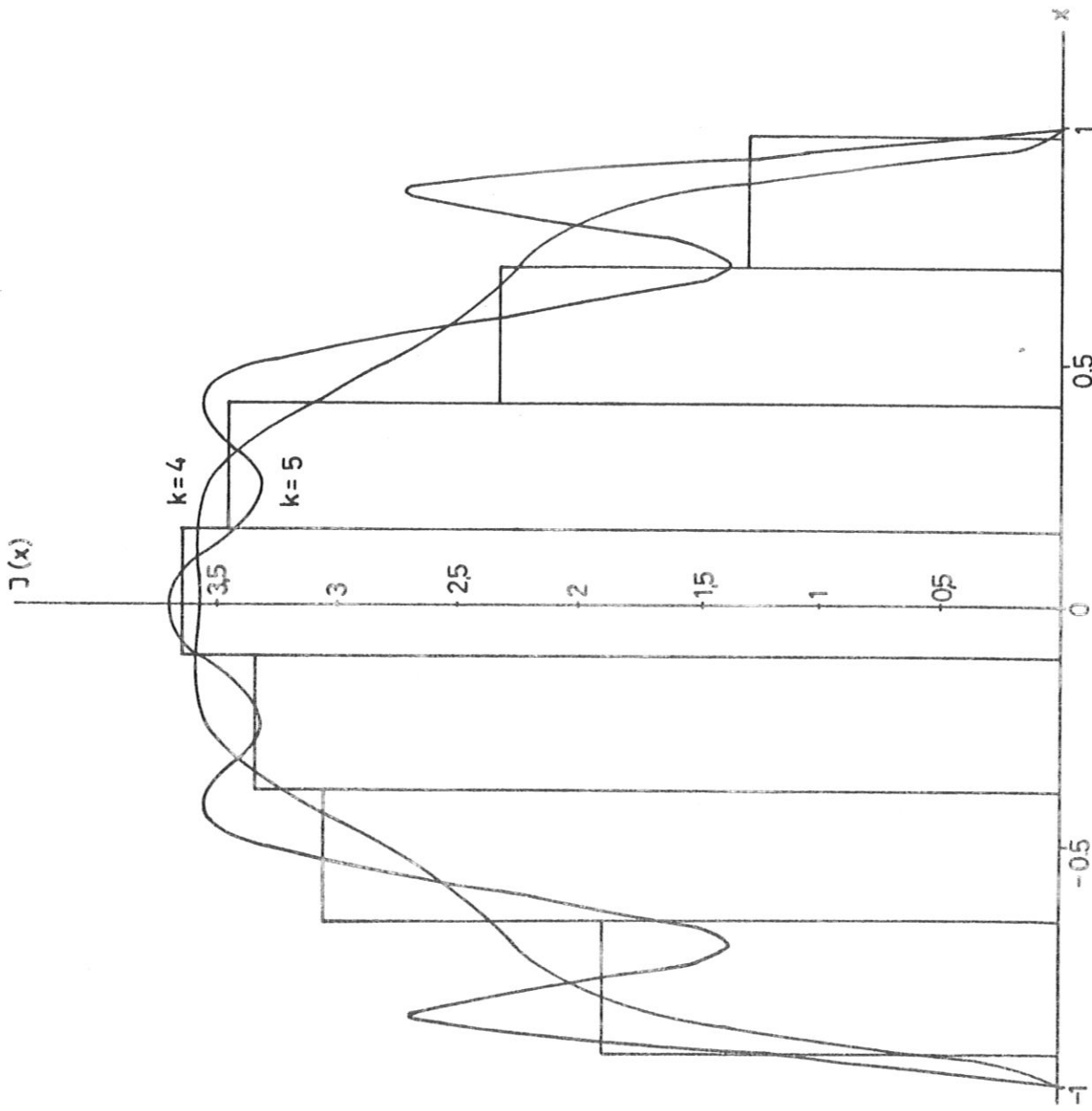


Fig.: 11

The fit with $k = 5$ is too close. $k = 4$ gives a more reasonable result.

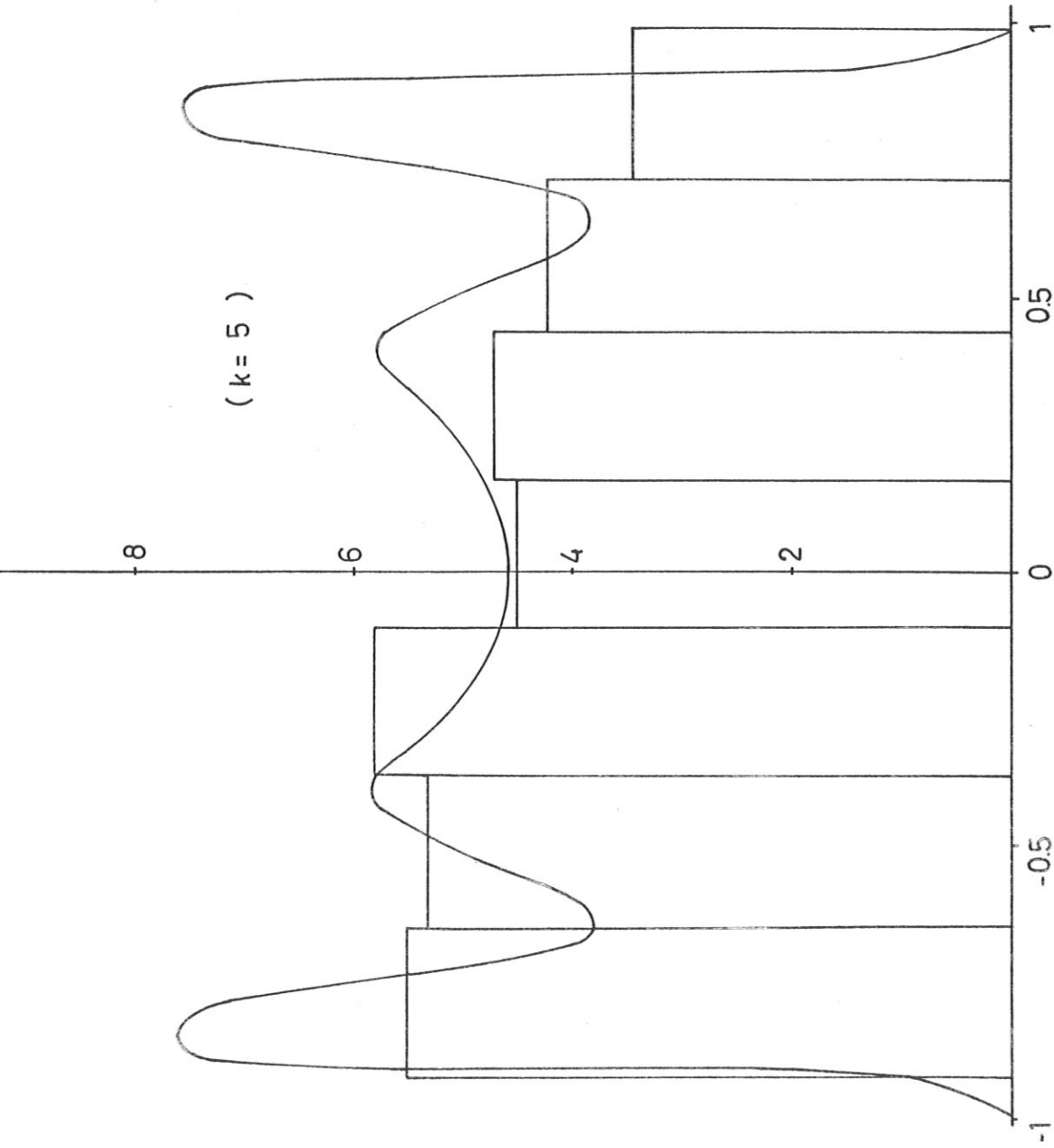


Fig.: 12

An example for the possible behaviour of $J(x)$ in the neighbourhood of $|x|=1$.