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# The dynamics of within-group and between-group interaction

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## Abstract

The article strengthens and provides a dynamic extension of the theory on collective rent seeking and private provision of a public good. Each individual agent within each group chooses in continuous or discrete time a continuous or discrete effort level. The combined effort within each group provides within-group public goods which are used as an input in the between-group  $n$ -group competition for an external prize. Intergroup mobility and intergroup warfare are allowed for. Each group and each individual agent within each group get a fraction of the prize based on a linear combination of equity and relative effort. A model/algorithm is developed generating analytical results and simulations illustrating how the interaction within and between groups proceeds through time.

*JEL classification:* C72; C73; D72

*Keywords:* Non-cooperative games; Evolutionary game dynamics; Rent seeking; Within-group strategy selection; Between-group competition; Micro-macro link

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## 1. Introduction

The article strengthens and provides a dynamic extension of the theory on rent seeking, pioneered by Tullock (1967, 1980), supplemented by Katz et al. (1990) and Nitzan (1991). A model/algorithm is developed generating analytical results and simulations illustrating the interaction within and between groups of generally different and arbitrary size, allowing for intergroup mobility and intergroup warfare. The means employed in the between-group  $n$ -group competition are public goods generated by within-group interaction in each group.

Each individual agent  $A_{ij}$  within each group  $A_i, i = 1, \dots, n$ , chooses at time  $t$  a cost  $S_{ij}$  of effort,  $0 \leq S_{ij} \leq c_i$ . The combined cost  $S_i$  of effort within group  $A_i$  provides within-group public goods which are used in the between-group  $n$ -group competition for an external prize  $E$ . Each group  $A_i$  and each individual agent  $A_{ij}$  within each group get a fraction of the prize  $E$  based on a linear combination of equity and relative effort.

Section 2 presents the static model involving within-group and between-group interaction. Section 3 provides static equilibrium analysis. Section 4 develops an eight-step continuous and discrete time algorithm for an arbitrary number  $n$  of groups  $A_i$  each consisting of an arbitrary number  $n_i(t)$  of individual agents  $A_{ij}$ . Section 5 presents the dynamic model assuming fixed group sizes  $n_i(t)$ , and employs it to a continuous and discrete strategy within-group  $n_i(t)$ -person prisoner's dilemma in each group  $A_i$ , accounting for between-group  $n$ -group competition for an external prize  $E$ . A simulation allowing the two strategies 'always cooperate' and 'always defect' for two groups is performed. Section 6 presents the dynamic model assuming intergroup mobility, and performs simulations. Section 7 presents the dynamic model assuming intergroup warfare, illustrating with simulations. Section 8 considers the combined operation of intergroup mobility and intergroup warfare, and provides simulations. Section 9 introduces the  $n_i(t)$ -person reactive strategies 'tit-for-tat' and 'bully' and gives simulations for one group in isolation.

2. Static model

Assume  $n$  groups  $A_i, i = 1, \dots, n$ , with  $n_i(t) = n_i$  individual agents  $A_{ij}, j = 1, \dots, n_i$ , in each group (Fig. 1). Each individual agent  $A_{ij}$  chooses a strategy  $S_{ij}$ , by which we mean that he has a cost  $S_{ij}$  of effort,  $0 \leq S_{ij} \leq c_i$ , so that each group  $A_i$  has a cost of effort

$$S_i \stackrel{\text{def}}{=} \sum_{j=1}^{n_i} S_{ij}, \tag{2.1}$$

where  $0 \leq S_i \leq n_i c_i$ . All the  $n$  groups  $A_i$  have a cost of effort

$$S \stackrel{\text{def}}{=} \sum_{i=1}^n S_i = \sum_{i=1}^n \sum_{j=1}^{n_i} S_{ij}, \tag{2.2}$$

where  $0 \leq S \leq \sum_{i=1}^n n_i c_i$ .

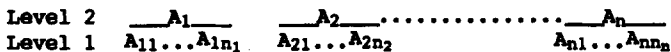


Fig. 1. Hierarchy represented by matrix A.

Each group  $A_i$  is involved in a within-group game which determines the provision of a public good that is used as an input in the competition for a between-group prize,  $E$ , against the other groups. Let an effort of  $c_i$  in group  $A_i$  produce within-group public goods  $B_i(t) = B_i$ . An effort of  $S_{ij}$  by an individual agent  $A_{ij}$  thus generates within-group public goods  $B_i S_{ij}/c_i$ , and an effort  $S_i$  by group  $A_i$  generates within-group public goods  $B_i S_i/c_i$ . Interesting cases of  $B_i$  are when  $B_i$  is constant and when  $B_i$  is proportional to  $n_i$ .  $B_i$  being constant means that the public goods produced by a cooperative act  $c$  is divided between the group members  $A_{ij}$ , giving a smaller share to each as the group size  $n_i$  increases.  $B_i$  being proportional to  $n_i$  means that the benefits reaped by one individual group member do not reduce the benefits received by another group member.

The between-group  $n$ -group competition for the prize  $E$  gives a fraction  $f_i$ , i.e.  $f_i E$  to group  $A_i$ , and a fraction  $f_{ij}$  of this, i.e.  $f_{ij} f_i E$  to individual agent  $A_{ij}$ . The one-period (static) payoff  $P_i$  to each group  $A_i$  is thus

$$P_i^{\text{mod}} = \frac{S_i}{c_i} B_i + f_i E - S_i, \tag{2.3}$$

where mod means that this is a model of nature. (2.3) says that group  $A_i$  has a total cost  $S_i$  of effort, which generates a fraction  $S_i/n_i c_i$  of the within-group public goods  $n_i B_i$ , and a fraction  $f_i$  of the between-group prize  $E$ . Similarly, the one-period payoff  $P_{ij}$  to each individual agent  $A_{ij}$  is

$$P_{ij}^{\text{mod}} = f_{ij} \left[ \frac{S_i}{c_i} B_i + f_i E \right] - S_{ij}. \tag{2.4}$$

A variety of different distribution rules  $f_{ij}$  and  $f_i$  can be considered. Nitzan (1991, p. 1524) uses the rule

$$f_{ij}^{\text{mod}} = \frac{(1 - a_i) S_{ij}}{S_i} + \frac{a_i}{n_i}, \tag{2.5}$$

which means that a proportion  $a_i$  is distributed within group  $A_i$  on egalitarian grounds (i.e. equal portion to each) and the remaining proportion  $1 - a_i$  is distributed according to the relative amount of public goods produced by each individual agent  $A_{ij}$ , which corresponds to the relative effort of each individual agent  $A_{ij}$ . Analogously for the between-group distribution, set

$$f_i^{\text{mod}} = \frac{(1 - a) \frac{B_i}{c_i} S_i}{\sum_{i=1}^n \frac{B_i}{c_i} S_i} + \frac{a}{n}, \tag{2.6}$$

where a proportion  $a$  is distributed between the  $n$  groups  $A_i$  on egalitarian grounds and the remaining proportion  $1 - a$  is distributed according to the relative amount of public goods produced by each group  $A_i$ . An alternative to (2.6) is

$$f_1 = \dots = f_m^{\text{mod}} = 1/m, f_{m+1} = \dots = f_n^{\text{mod}} = 0 \tag{2.7}$$

where  $S_1 = \dots = S_m > S_{m+1} \geq \dots \geq S_n$  and  $1 \leq m \leq n$ , which means that the group (assume  $A_1$  without loss of generality) producing the highest effort  $S_1$  gets the entire between-group prize  $E$ , or, if  $m$  groups equally produce the highest effort, that these equally share the prize  $E$  among themselves. In the remainder of the article I assume (2.6). Inserting (2.5) and (2.6) into (2.4) gives

$$\begin{aligned}
 P_{ij} &= \left[ \frac{(1 - a_i)S_{ij}}{S_i} + \frac{a_i}{n_i} \right] \left[ \frac{S_i}{c_i} B_i + \left( \frac{(1 - a) \frac{B_i}{c_i} S_i}{\sum_{i=1}^n \frac{B_i}{c_i} S_i} + \frac{a}{n} \right) E \right] - S_{ij} \\
 &= \frac{(1 - a_i) B_i}{c_i} S_{ij} + \frac{a_i B_i}{n_i c_i} S_i + (1 - a_i) E \frac{(1 - a) \frac{B_i}{c_i} S_{ij}}{\sum_{i=1}^n \frac{B_i}{c_i} S_i} \\
 &\quad + (1 - a_i) E \frac{a}{n} \frac{S_{ij}}{S_i} + \frac{a_i E}{n_i} \frac{(1 - a) \frac{B_i}{c_i} S_i}{\sum_{i=1}^n \frac{B_i}{c_i} S_i} + \frac{a_i E}{n_i} \frac{a}{n} - S_{ij}. \tag{2.8}
 \end{aligned}$$

**3. Static equilibrium analysis**

The first order condition for a global interior maximum of  $P_{ij}$  for individual agent  $A_{ij}$  is

$$\begin{aligned}
 \frac{\partial P_{ij}}{\partial S_{ij}} &= \frac{(1 - a_i) B_i}{c_i} + \frac{a_i B_i}{n_i c_i} + (1 - a_i) E (1 - a) \frac{B_i}{c_i} \left( \frac{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right) - \frac{B_i}{c_i} S_{ij}}{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right)^2} \right) \\
 &\quad + (1 - a_i) E \frac{a}{n} \left( \frac{S_i - S_{ij}}{S_i^2} \right) + \frac{a_i E}{n_i} (1 - a) \frac{B_i}{c_i} \left( \frac{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right) - \frac{B_i}{c_i} S_i}{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right)^2} \right) \\
 - 1 &= 0. \tag{3.1}
 \end{aligned}$$

In the remainder of the article I set  $a_i = 1$  and  $a = 0$ , which means that the payoff available to each group  $A_i$  is divided equally among its members  $A_{ij}$ , and that each group  $A_i$  gets a fraction of the between-group prize  $E$  proportional to the relative amount of public goods it produces compared with the other groups. (3.1) then becomes

$$E \frac{B_i}{c_i} \left( \frac{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right) - \frac{B_i}{c_i} S_i}{\left( \sum_{i=1}^n \frac{B_i}{c_i} S_i \right)^2} \right) = n_i - \frac{B_i}{c_i}. \tag{3.2}$$

Setting  $B_i = B$ ,  $c_i = c$ , and summing both sides in (3.2) from  $i = 1$  to  $i = n$ , i.e.  $\sum_{i=1}^n (\cdot)$ , gives

$$E \frac{S(n-1)}{S^2} = \left( \sum_{i=1}^n n_i \right) - \frac{nB}{c}, \tag{3.3}$$

i.e.

$$S = \frac{(n-1)E}{\left( \sum_{i=1}^n n_i \right) - \frac{nB}{c}}. \tag{3.4}$$

If we assume that all the groups  $A_i$  are equally large, i.e.  $n_1 = n_2 = \dots = n_n$ , and that identical individual agents  $A_{ij}$  in a symmetric Nash equilibrium choose the same effort level  $S_{ij}$ , which implies  $S = nS_i = nn_iS_{ij}$ , (3.4) gives that

$$S_{ij} = \frac{(n-1)E}{nn_i \left( nn_i - \frac{nB}{c} \right)} = \frac{(n-1)Ec}{n^2 n_i^2 \left( c - \frac{B}{n} \right)}. \tag{3.5}$$

Checking the border conditions  $0 \leq S_{ij} \leq c$ , (3.5) shows that the former implies  $c > B/n_i$  and the latter implies

$$E \leq \frac{n^2 n_i^2}{(n-1)} \left( c - \frac{B}{n_i} \right). \tag{3.6}$$

Inserting  $S_{ij} = c$  and (3.5) into (2.8) with  $a_i = 1$  and  $a = 0$  shows that the former maximizes  $P_{ij}$  when (3.6) is not satisfied. That is, when  $E$  is above that level given by (3.6), it is beneficial for an individual agent  $A_{ij}$  to limit his effort to  $S_{ij} = c$  rather than choose the higher  $S_{ij}$  according to (3.5).

If we assume  $0 < B_i/n_i < c_i < B_i$ , the within-group game is a ‘continuous’ version of an  $n_i$ -person prisoner’s dilemma, in the sense that each individual agent  $A_{ij}$  can choose any effort level  $S_{ij}$ , i.e. degree of cooperation, on a continuous scale from 0 to  $c_i$ .

If we assume the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ), we get the standard (discrete)  $n_i$ -person prisoner's dilemma. Define  $h_{ic} - 1$  as the number of cooperating individual agents aside from individual agent  $A_{ij}$  in group  $A_i$ . If individual agent  $A_{ij}$  chooses  $S_{ij} = c_i$ , giving  $h_{ic}$  cooperators in group  $A_i$ , he gets according to (2.8) with  $a_i = 1$  and  $a = 0$ , a payoff

$$P_{ijc} = \frac{1}{n_i} \left[ h_{ic} B_i + \frac{h_{ic} B_i}{\sum_{i=1}^n h_{ic} B_i} E \right] - c_i \text{ if } S_{ij} = c_i. \quad (3.7)$$

If individual agent  $A_{ij}$  chooses  $S_{ij} = 0$ , giving  $h_{ic} - 1$  cooperators in group  $A_i$ , he gets according to (2.8) with  $a_i = 1$  and  $a = 0$ , a payoff

$$P_{ijd} = \frac{1}{n_i} \left[ (h_{ic} - 1) B_i + \frac{(h_{ic} - 1) B_i}{\left( \sum_{i=1}^n h_{ic} B_i \right) - B_i} E \right] \text{ if } S_{ij} = 0. \quad (3.8)$$

When the between-group prize  $E = 0$ , the unique Nash equilibrium for each individual agent  $A_{ij}$  in group  $A_i$  is the well-known mutual defection  $S_{ij} = 0$ . As  $E$  increases, a minimum value of  $E$  can be determined which guarantees cooperation  $S_{ij} = c_i$  by all the  $n_i$  individual agents  $A_{ij}$  in the  $n$  groups  $A_i$ . The most restrictive requirement for  $E$ , to ensure cooperation by all the individual agents  $A_{ij}$ , is derived by assuming there are exclusively cooperators in the  $n$  groups, and that one individual agent  $A_{ij}$  in group  $A_i$  contemplates whether to defect. Hence set  $h_{ic} = n_i$  in (3.7) and (3.8). Cooperation by the individual agent  $A_i$  follows if (3.7) is larger than (3.8), i.e. if

$$\frac{1}{n_i} \left[ n_i B_i + \frac{n_i B_i}{\sum_{i=1}^n n_i B_i} E \right] - c_i > \frac{1}{n_i} \left[ (n_i - 1) B_i + \frac{(n_i - 1) B_i}{\left[ \left( \sum_{i=1}^n n_i B_i \right) - B_i \right]} E \right], \quad (3.9)$$

i.e. if

$$E > \frac{\left( \sum_{i=1}^n n_i B_i \right) \left[ \left( \sum_{i=1}^n n_i B_i \right) - B_i \right]}{\left[ \left( \frac{1}{n_i} \sum_{i=1}^n n_i B_i \right) - B_i \right] B_i} \left[ c_i - \frac{B_i}{n_i} \right], \quad (3.10)$$

where  $n_i > 0 \forall i = 1, \dots, n$ , to avoid possible division with 0. That is, groups  $A_i$  without members  $A_{ij}$  are not included in (3.10), which is consistent with these not

getting any fraction of the between-group prize  $E$ . Note from (3.10) that the requirement for  $E$  when  $B_i = B$  is most strict for the largest group and least strict for the smallest group. Calculating the right-hand side of (3.10)  $\forall i = 1, \dots, n$ , choose that  $i$  that generates the strictest requirement for  $E$ .

#### 4. The algorithm

One common algorithm to determine the relative fitnesses of several strategies, when there is only one group, is employed by Martinez-Coll and Hirshleifer (1991), Maynard Smith (1982), Schuster and Sigmund (1985), Taylor and Jonker (1978), and Zeeman (1981). That algorithm makes five major simplifications:

1. It assumes only one group, in isolation/vacuum, not several groups, nor a hierarchy.
2. It assumes for each individual agent  $A_{ij}$  a finite number of discrete strategies  $S_{ij}$  rather than a continuous effort level  $S_{ij}$ , say  $0 \leq S_{ij} \leq c_i$ .
3. It does not specify group size, nor specific numbers of individual agents playing the various strategies. It specifies instead what fractions of the group play the various strategies. These fractions change through time dependent upon what strategies are most fit.
4. It assumes that the fractions play two-person games, not  $n_i(t)$ -person games.
5. It assumes that the reactive strategies, tit-for-tat and bully, have instantaneous response capability. For a two-person game this implies that a tit-for-tat player will be able to distinguish between pure cooperators and pure defectors after having met one pure defector once. This is because the tit-for-tat player will cooperate and be exploited the first time it meets a defector, the defector enjoying a free ride, while, in subsequent time periods, whenever a tit-for-tat player meets a defector, the tit-for-tat player will also defect.

Generalizing that algorithm, this section rectifies these five simplifications:

1. It assumes a two-level hierarchy  $A$  and an arbitrary number  $n$  of groups  $A_i$ .
2. It assumes a continuous and discrete time model. It posits for each individual agent  $A_{ij}$  a continuous strategy model with a continuous effort level  $S_{ij}$ , where  $0 \leq S_{ij} \leq c_i$ , and a discrete strategy model allowing for a finite number of discrete strategies  $S_{ij}$ .
3. It specifies the exact number  $n_i(t)$  of individual agents  $A_{ij}$  in each group  $A_i$  at time  $t$ , and the exact fraction  $p_i(s, t)$  or  $p_{is}(t)$ , and exact number  $h_{is}(t)$  playing each strategy  $s = S_{ij}$ , in group  $A_i$  at time  $t$ .
4. It allows  $n_i(t)$  to change through intergroup mobility, intergroup warfare, or analogs of reproduction and death of individual agents  $A_{ij}$  in each group  $A_i$ . The

changes  $dn_i(t)/dt$  and  $dp_i(s, t)/dt$  in  $n_i(t)$  and  $p_i(s, t)$  (or  $p_{is}(t)$ ) depend upon the within-group  $n_i(t)$ -person games, and between-group  $n$ -group games.

5. It acknowledges in a discrete time multi-period game that the  $n_i(t)$ -person reactive strategies, tit-for-tat and bully, account for the conditions and strategies  $S_{ij}$  played both in the current time period  $t$  and the previous time period  $t - 1$ .

The eight-step algorithm is as follows:

1. Set initial conditions: specify the number  $n$  of groups  $A_i$ , and the initial number  $n_i(0)$  of individual agents  $A_{ij}$  in each group  $A_i$  at time  $t = 0$ . Specify the strategy set  $S_{ij}$  of available strategies  $S_{ij}$  for each individual agent  $A_{ij}$ . For the continuous strategy model, define the function  $p_i(s, t)$  as the fraction of individual agents  $A_{ij}$  that play the strategy  $s = S_{ij}$ ,  $0 \leq s = S_{ij} \leq c_i$  at time  $t$ , where  $t$  is continuous or discrete,  $t = 0, 1, 2, \dots$ . For the discrete strategy model, define the initial number  $h_{is}(0)$  of individual agents  $A_{ij}$  playing each strategy  $S_{ij}$  at time  $t = 0$ . Define

$$p_{is}(t) \stackrel{\text{def}}{=} \frac{h_{is}(t)}{n_i(t)} \tag{4.1}$$

as the fraction of individual agents  $A_{ij}$  in group  $A_i$  that chooses strategy  $s = S_{ij}$  at time  $t$ . Observe that

$$\int_0^{c_i} p_i(s, t) ds = 1, \quad \sum_{S_{ij}} p_{is}(t) = 1, \quad \text{and} \quad \sum_{S_{ij}} h_{is}(t) = n_i(t). \tag{4.2}$$

2. Determine the mean yield  $Y_i(s, t)$  (for continuous strategies  $s = S_{ij}$ ,  $0 \leq s = S_{ij} \leq c_i$ ) or  $Y_{is}(t)$  (for discrete strategies  $S_{ij}$ ) that individual agent  $A_{ij}$  in group  $A_i$  can expect to receive by playing strategy  $S_{ij}$  at time  $t$ , using (2.8) (continuous strategy version), (3.7) and (3.8) (discrete strategy version), or the model to be analyzed.

3. The overall mean yield  $Y_i(t)$  for an individual agent  $A_{ij}$  in group  $A_i$  is

$$Y_i(t) \stackrel{\text{mod}}{=} \int_0^{c_i} p_i(s, t) Y_i(s, t) ds \tag{4.3}$$

for the continuous strategy model and

$$Y_i(t) \stackrel{\text{mod}}{=} \sum_{S_{ij}} p_{is}(t) Y_{is}(t) \tag{4.4}$$

for the discrete strategy model.



The overall mean yield  $Y(t)$  for an individual agent  $A_{ij}$  in any group  $A_i$ ,  $i = 1, \dots, n$ , is

$$Y(t) \stackrel{\text{mod}}{=} \frac{1}{N(t)} \sum_{i=1}^n n_i(t) Y_i(t), \tag{4.5}$$

where  $Y_i(t)$  is given by (4.3) or (4.4) and

$$N(t) \stackrel{\text{def}}{=} \sum_{i=1}^n n_i(t) \tag{4.6}$$

is the total number of individual agents  $A_{ij}$  in all the groups  $A_i$ ,  $i = 1, \dots, n$ .

4. The relative fitness  $F_i(s, t)$  for strategy  $S_{ij}$  in group  $A_i$  is

$$F_i(s, t) \stackrel{\text{def}}{=} Y_i(s, t) - Y_i(t) \tag{4.7}$$

for the continuous strategy model, substituting  $F_i(s, t) \rightarrow F_{is}(t)$  and  $Y_i(s, t) \rightarrow Y_{is}(t)$  for the discrete strategy model.

5. The change  $\Delta p_i(s, t) = p_i(s, t + 1) - p_i(s, t)$  in the fraction of individual agents  $A_{ij}$  in group  $A_i$  that play strategy  $S_{ij}$  is

$$\begin{aligned} p_i(s, t + 1) - p_i(s, t) &\stackrel{\text{mod}}{=} k_i p_i(s, t) F_i(s, t) \\ &= k_i p_i(s, t) [Y_i(s, t) - Y_i(t)] \end{aligned} \tag{4.8}$$

for the continuous strategy model, substituting  $p_i(s, t) \rightarrow p_{is}(t)$ ,  $F_i(s, t) \rightarrow F_{is}(t)$ , and  $Y_i(s, t) \rightarrow Y_{is}(t)$  for the discrete strategy model. The change  $\Delta p_i(s, t)$  is positive if the mean yield  $Y_i(s, t)$  that individual agent  $A_{ij}$  can expect to receive by playing strategy  $s = S_{ij}$  is larger than the overall mean yield  $Y_i(t)$  that an individual agent  $A_{ij}$  in group  $A_i$  can expect to receive, and negative if the reverse is the case.  $k_i$  is the rapidity of change, or ‘sensitivity’ of the process. If  $k_i$  is too large we may get irregular movement of the process for group  $A_i$ . To ensure stability,  $k_i$  must be sufficiently small. If  $k_i$  is too small, there will be no change in the fraction. (4.8) is a difference equation. Assuming  $\Delta t = 1$ , it can be written as a differential equation, viz.

$$\frac{d p_i(s, t)}{d t} \stackrel{\text{mod}}{=} k_i p_i(s, t) [Y_i(s, t) - Y_i(t)] \tag{4.9}$$

for the continuous strategy model, with the above substitutions for the discrete strategy model.

The change  $dn_i(t)/dt$  in the number  $n_i(t)$  of individual agents  $A_{ij}$  in group  $A_i$  is given by

$$\frac{d n_i(t)}{d t} \stackrel{\text{mod}}{=} g(r_i, n_i(t), p_i(s, t), Y_i(s, t), Y_i(t), Y(t)), \tag{4.10}$$

for the continuous strategy model, with the above substitutions for the discrete strategy model, where  $r_i$  is the ‘sensitivity’ of the process for group  $A_i$ . Examples of functions  $g(\cdot)$  are considered in Sections 6 (intergroup mobility) and 7 (intergroup warfare).

6. For the discrete time model, let  $t + 1 \rightarrow t$ . For the discrete strategy model, calculate  $h_{is}(t) = p_{is}(t)/n_i(t)$ , and round off the specific number  $h_{is}(t)$  of individual agents  $A_{ij}$  playing strategy  $S_{ij}$  to the nearest integer value:

$$h_{is}(t) \leftarrow \text{round}(h_{is}(t)). \tag{4.11}$$

7. Print out the relevant curves;  $Y_i(s, t)$ ,  $p_i(s, t)$ ,  $h_i(s, t)$ ,  $Y_{is}(t)$ ,  $p_{is}(t)$ ,  $h_{is}(t)$ , etc.

8. Go to 2.

**5. Dynamic model with fixed group sizes  $n_i(t)$**

Using the dynamic version of (2.8) to express (4.9) involves writing the effort level  $S_i$  for group  $A_i$  in terms of  $p_i(s, t)$ , accounting on the one hand for the effort level of the other  $n_i(t) - 1$  individual agents in group  $A_i$ , and on the other hand for the specific effort level  $S_{ij} = s$  chosen by the particular individual agent  $A_{ij}$ . Hence, referring to (2.1),  $S_i$  can be written as

$$S_i \stackrel{\text{def}}{=} (n_i(t) - 1) \int_0^{c_i} p_i(s, t) s \, ds + s. \tag{5.1}$$

Inserting (5.1) into (2.8) with  $a_i = 1$  and  $a = 0$  gives

$$\begin{aligned} Y_i(s, t) = & \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) \\ & + \left\{ \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) E \right\} \\ & \times \left\{ \left[ \sum_{i=1}^n \frac{B_i(t)}{c_i} n_i(t) \int_0^{c_i} p_i(s, t) s \, ds \right] \right. \\ & \left. + \left( \frac{s}{c_i} - \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds \right) B_i(t) \right\}^{-1} - s. \tag{5.2} \end{aligned}$$

Inserting (4.3) and (5.2) into (4.9) gives

$$\begin{aligned}
 & \frac{dp_i(s, t)}{dt} \\
 &= k_i p_i(s, t) \left[ \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) \right. \\
 & \quad + \left. \left\{ \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) E \right\} \right. \\
 & \quad \times \left. \left[ \sum_{i=1}^n \frac{B_i(t)}{c_i} n_i(t) \int_0^{c_i} p_i(s, t) s \, ds \right] \right. \\
 & \quad + \left. \left( \frac{s}{c_i} - \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds \right) B_i(t) \right\}^{-1} - s \\
 & \quad - \int_0^{c_i} p_i(s, t) \left[ \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) \right. \\
 & \quad + \left. \left\{ \left[ \left( 1 - \frac{1}{n_i(t)} \right) \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds + \frac{s}{n_i(t)c_i} \right] B_i(t) E \right\} \right. \\
 & \quad \times \left. \left[ \sum_{i=1}^n \frac{B_i(t)}{c_i} n_i(t) \int_0^{c_i} p_i(s, t) s \, ds \right] \right. \\
 & \quad + \left. \left( \frac{s}{c_i} - \frac{1}{c_i} \int_0^{c_i} p_i(s, t) s \, ds \right) B_i(t) \right\}^{-1} - s \Big] ds, \tag{5.3}
 \end{aligned}$$

where  $i = 1, 2, \dots, n$  and  $0 \leq s \leq c_i$ . (5.3) is an integro-differential equation which is most easily solved using the explicit method (i.e. rewriting with (4.8)) for the time integration, and a standard numerical method for the strategy integrals.

To derive the analog of (5.3) for discrete strategies  $S_{ij}$ , set

$$p_i(s, t) \stackrel{\text{mod}}{=} \sum_{S_{ij}} p_{is}(t) \delta(s - S_{ij}) \stackrel{S_{ij}=c_i}{=} p_{ic}(t) \delta(s - c_i) + p_{id}(t) \delta(s), \tag{5.4}$$

where the rightmost two terms give the special case where only the two strategies  $S_{ij} = c_i$  and  $S_{ij} = 0$  are allowed, and where  $\delta$  is Dirac's delta 'function' defined by

$$\int_{-\infty}^{\infty} p_i(s, t) \delta(s - S_{ij}) \, ds \stackrel{\text{def}}{=} p_i(S_{ij}, t), \text{ and } \delta(s - S_{ij}) \stackrel{\text{def}}{=} 0 \text{ for } s \neq S_{ij}. \tag{5.5}$$

With the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ), where  $0 < B_i(t)/n_i(t) < c_i < B_i(t)$ , the within-group game

is the standard (discrete)  $n_i(t)$ -person prisoner's dilemma. Inserting (5.4) into (5.3) and simplifying gives the dynamic model for discrete strategies  $S_{ij}$ , i.e.

$$\begin{aligned}
 & \frac{dp_{ic}(t)}{dt} \\
 &= k_i p_{ic}(t)(1 - p_{ic}(t)) \left( \frac{B_i(t)}{n_i(t)} - c_i \right. \\
 & \quad + \left. \left\{ \left( \left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) + \frac{1}{n_i(t)} \right) B_i(t) E \right\} \right. \\
 & \quad \times \left. \left. \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) + (1 - p_{ic}(t)) B_i(t) \right]^{-1} \right\} \right. \\
 & \quad - \left. \left\{ \left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) B_i(t) E \right\} \right. \\
 & \quad \times \left. \left. \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right]^{-1} \right\} \right) \\
 &= k_i p_{ic}(t)(1 - p_{ic}(t)) \left( \frac{B_i(t)}{n_i(t)} - c_i \right. \\
 & \quad + \left. \left\{ \left[ \left( \frac{1}{n_i(t)} \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right] B_i(t) E \right\} \right. \\
 & \quad \times \left. \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) + (1 - p_{ic}(t)) B_i(t) \right] \right. \\
 & \quad \times \left. \left. \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right]^{-1} \right\} \right), \tag{5.6}
 \end{aligned}$$

where  $i = 1, 2, \dots, n$ , and  $p_{id}(t) = p_{ic}(t) - 1$ . (5.6) has a very nice interpretation. Within the rightmost parentheses, the first two terms  $B_i(t)/n_i(t) - c_i$  describe the influence of the within-group  $n_i(t)$ -person prisoner's dilemma, and have a negative effect upon  $dp_{ic}(t)/dt$  to a degree corresponding to the 'severity' of the within-group  $n_i(t)$ -person prisoner's dilemma. Note that if  $c_i < B_i(t)/n_i(t)$ , the within-group game is a pure cooperation game. If  $c_i > B_i(t)/n_i(t)$ , the within-group game is a pure defection game. The range of  $c_i$  from  $B_i(t)/n_i(t)$  to  $B_i(t)$  thus specifies how 'severe' the within-group  $n_i(t)$ -person prisoner's dilemma is. The last two terms within these parentheses, simplified to one term in the lower part of the equation, describe the influence of the between-group  $n$ -group competition for

the external prize  $E$ , and have a positive effect upon  $dp_{ic}(t)/dt$  to a degree corresponding to the size of  $E$ . (5.6) is a non-linear coupled ordinary differential equation of the first order. Although it has no straightforward analytical solution, it is richly endowed with information.

Given  $0 \leq p_{ic}(t) \leq 1$ , (5.6) specifies three qualitatively different stationary values for  $p_{ic}(t)$ , i.e. the values that  $p_{ic}(t)$  eventually tends towards as  $t$  approaches infinity. The stationary values of  $p_{ic}(t)$  are found by setting the left-hand side  $dp_{ic}(t)/dt$  equal to zero. The first stationary value,  $p_{ic}(t) = 1$ , follows from setting the second term on the right-hand side of (5.6) equal to zero, i.e.  $1 - p_{ic}(t) = 0$ , and requiring that the term within the rightmost parentheses is strictly positive, which ensures that  $p_{ic}(t)$  is strictly monotonically increasing towards its highest possible value  $p_{ic}(t) = 1$ . Expressing this requirement in terms of  $E$  gives

$$\begin{aligned}
 E > & \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) + (1 - p_{ic}(t)) B_i(t) \right] \\
 & \times \left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right] \\
 & \times \left\{ \left[ \left( \frac{1}{n_i(t)} \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right] B_i(t) \right\}^{-1} \\
 & \times \left[ c_i - \frac{B_i(t)}{n_i(t)} \right], \tag{5.7}
 \end{aligned}$$

which is equivalent to the static requirement (3.10) when  $p_{ic}(t) = 1 \forall i = 1, \dots, n$ . In other words, if (5.7) is satisfied due to a sufficiently large between-group prize  $E$ ,  $p_{ic}(t)$  will eventually and inevitably be driven towards one as  $t$  approaches infinity, ensuring that all the groups  $A_i$  will consist exclusively of cooperating individual agents  $A_{ij}$ .

If the rightmost parenthesis in (5.6) is negative,  $p_{ic}(t)$  is monotonically decreasing. The most extreme form for monotonic decrease involves setting  $E = 0$ , which means exclusively within-group and not between-group interaction. In this case (5.6) has an analytical solution given by

$$p_{ic}^{E=0}(t) = \frac{1}{1 - \left( 1 - \frac{1}{p_{ic}(0)} \right) e^{(c_i - B_i(t)/n_i(t))k_i t}}, \tag{5.8}$$

where  $p_{ic}(0)$  is the initial value of  $p_{ic}(t)$  when  $t = 0$ . We see that  $p_{ic}^{E=0}(t) \rightarrow 0$  when  $t \rightarrow \infty$  given that  $c_i > B_i(t)/n_i(t)$ , which gives the second stationary value

$p_{ic}(t) = 0$ . This means that the number of cooperating individual agents  $A_{ij}$  in each group  $A_i$  will eventually vanish, as is implied by an  $n_i(t)$ -person prisoner's dilemma specified by  $0 < B_i(t)/n_i(t) < c_i < B_i(t)$ . (We also see from (5.8) that  $p_{ic}^{E=0}(t) \rightarrow 1$  when  $t \rightarrow \infty$  given that  $c_i < B_i(t)/n_i(t)$ , since the within-group game then is a pure cooperation game.)

The third stationary value of  $p_{ic}(t)$ , for those  $i, i = 1, \dots, n$ , where  $E > 0$  and (5.6) is not satisfied, is found by setting the term within the rightmost parentheses in (5.6) equal to zero for those  $i$ , and then solving the equations to find  $p_{ic}(t) = p_{ic}(\infty)$ . These values  $p_{ic}(\infty)$  may be larger, equivalent, or smaller than their corresponding initial values  $p_{ic}(0)$ , meaning that  $p_{ic}(t)$  may be monotonically increasing, constant, or monotonically decreasing for different  $i, i = 1, \dots, n$ , as this follows from the rightmost term being positive, zero, or negative for different  $i$ , given the parameter values and how the various  $p_{ic}(t)$  proceed over time.

(5.7) contains ample information of how the parameters  $n, n_i(t), B_i(t), c_i$ , and  $p_{ic}(t)$  affect the requirement for  $E$  to ensure cooperation by the individual agents  $A_{ij}$  in group  $A_i$ . The requirement increases (becomes more strict) with  $c_i$  and decreases with  $B_i(t)$ , as is consistent with the within-group  $n_i(t)$ -person prisoner's dilemma becoming more 'severe' as  $c_i - B_i(t)/n_i(t)$  increases. The requirement also increases with  $n_i(t)$ , as this gives more individual agents  $A_{ij}$  among which to divide the fraction of the prize  $E$  that group  $A_i$  receives.

Based on (5.6) and (5.7) we can specify how  $p_{ic}(t)$  depends upon various characterizations of  $E$ . Assume  $B_i(t)$  is proportional to  $n_i(t)$  (i.e.  $B_i(t) \sim n_i(t)$ ), and set for simplicity  $B_i(t) = B_i$ . If  $E$  is constant or  $E \sim n_i(t)$  then  $p_{ic}(t) \rightarrow 0$  when any  $n_u(t) \rightarrow \infty, i, u = 1, \dots, n$ . If  $E \sim n_i^2(t)$  then  $p_{ic}(t) \rightarrow 0$  when  $n_i(t) \rightarrow \infty$ . If  $E \sim n_u(t)$  then  $0 \leq p_{ic}(t) \leq 1$  when  $n_u(t) \rightarrow \infty, u \neq i, i, u = 1, \dots, n$ . If  $E \sim n_i^3(t)$  then  $0 \leq p_{ic}(t) \leq 1$  when  $n_i(t) \rightarrow \infty$ . If  $E \sim n_u^2(t)$  then  $p_{ic}(t) \rightarrow 1$  when  $n_u(t) \rightarrow \infty, u \neq i$ . If  $E \sim n_i^4(t)$  then  $p_{ic}(t) \rightarrow 1$  when  $n_i(t) \rightarrow \infty$ .

Simulating (5.3) and/or (5.6) enables us to predict a solution for generally all parameter values  $n, n_i(t), B_i(t), c_i$ , and  $E$ , accounting for the combined within-group  $n_i(t)$ -person prisoner's dilemma and between-group  $n$ -group competition for a prize  $E$  over time. Simulate (5.6) allowing for the two strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ). Set  $n = 2$  and posit fixed group sizes  $n_i(t) = n_i = 1000, i = 1, 2$ . Considering the case with  $B_i(t)$  proportional to  $n_i(t)$ , set  $B_i(t) = n_i(t)$ .  $c_i$  has to be chosen so that  $0 < B_i(t)/n_i(t) < c_i < B_i(t)$ . Assuming a low degree of 'severity' for the within-group  $n_i(t)$ -person prisoner's dilemma, set  $c_i = 2$ . With these parameter values, the requirement (5.7) for  $E$  to guarantee cooperation by all the individual agents  $A_{ij}$  in the static case is  $E > 3,998,000$ . If  $E = 0$ , the fraction  $p_{ic}(t)$  of cooperating individual agents  $A_{ij}$  will go exponentially towards 0 as described by (5.8). If  $E > 3,998,000$ , the fraction  $p_{ic}(t)$  of cooperating individual agents  $A_{ij}$  will go exponentially towards one. Interesting cases of  $E$  are thus  $0 < E < 3,998,000$ , which gives an intermediate degree of cooperation in each group  $A_i$ .

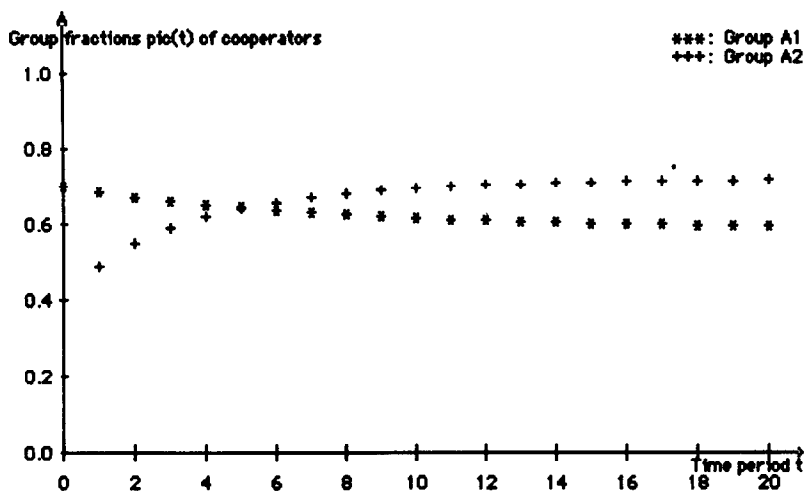


Fig. 2. Evolution of group fractions  $p_{ic}(t)$  of cooperators,  $E = 2,700,000$ ,  $k_i = 0.5$ .

Choose  $E = 2,700,000$ . With the above parameter values, the fraction  $p_{ic}(t)$  of cooperating individual agents  $A_{ij}$  in the two groups  $A_i$ ,  $i = 1, 2$ , will eventually be equal, i.e.  $p_{1c}(\infty) = p_{2c}(\infty)$ , regardless of initial conditions  $p_{ic}(0)$ . I will change this by setting  $B_1(t) = 0.9n_1(t)$  and  $B_2(t) = 1.1n_2(t)$ , which means that a given degree of cooperation, ceteris paribus, gives fewer public goods in group  $A_1$  than in group  $A_2$ , which implies that eventually  $p_{1c}(\infty) < p_{2c}(\infty)$ . Note that by setting  $B_1(t) = 0.9n_1(t)$  and  $B_2(t) = 1.1n_2(t)$ , the requirement (5.7) for group  $A_1$  increases to  $E > 4,442,444.4$ , while for group  $A_2$  it decreases to  $E > 3,634,363.6$ . Fig. 2 sets the initial fractions of cooperators  $c$  to  $p_{1c}(0) = 0.7$  and  $p_{2c}(0) = 0.4$  (i.e.  $h_{1c}(0) = 700$  and  $h_{2c}(0) = 400$ ) and simulates over 20 time periods with 'sensitivity' factor  $k_i = 0.5$ .

We see that the defectors in the least cooperative group  $A_2$  switch to cooperation while the cooperators in the most cooperative group  $A_1$  switch to defection. As the number of cooperating individual agents  $A_{ij}$  in the two groups becomes more equal, an equilibrium arises with a mixture of cooperators and defectors in each group  $A_i$ . This equilibrium is such that no single individual agent  $A_{ij}$  has any incentive to deviate unilaterally. Consistent with (5.7),  $p_{ic}(\infty)$ ,  $i = 1, 2$ , increases with  $E$  and  $B_i(t)$ , and decreases with  $c_i$  and  $n_i(t)$ .

### 6. Dynamic model with intergroup mobility

Allowing for intergroup mobility means positing a model (4.10) specifying a change  $dn_i(t)/dt$  in the number  $n_i(t)$  of individual agents  $A_{ij}$  in each group  $A_i$

over time. This change  $dn_i(t)/dt$  is such that individual agents  $A_{ij}$  flow between the groups  $A_i$ ,  $i = 1, \dots, n$ , preserving the total number  $N(t)$ , as given by (4.6), of individual agents  $A_{ij}$  in all the  $n$  groups  $A_i$ ,  $i = 1, \dots, n$ , i.e.  $N(t) = N = \text{constant}$ . I consider the model

$$\begin{aligned} \frac{dn_i(t)}{dt} \text{ mod} &= (\gamma_i - \mu_i)n_i(t)[Y_i(t) - Y(t)] \\ &= r_i n_i(t) \left[ Y_i(t) - \frac{1}{N} \sum_{i=1}^n n_i(t) Y_i(t) \right], \end{aligned} \quad (6.1)$$

using (4.5), where  $Y_i(t)$  is given by (4.3) inserting (5.2) for the continuous strategy model, and (4.4) inserting (5.2) and (5.4) for the discrete strategy model.  $r_i = \gamma_i - \mu_i > 0$  is the ‘compound’ sensitivity of the process, in the sense that  $\gamma_i$  is the sensitivity without switching costs, and  $\mu_i$  is the cost of switching group  $A_i$ . Switching costs  $\mu_i$  account for both entry costs (if  $Y_i(t) - Y(t)$  is positive) and exit costs (if  $Y_i(t) - Y(t)$  is negative) for individual agents  $A_{ij}$  switching to and from, respectively, group  $A_i$ . Switching costs  $\mu_i$  slow down the rapidity of change  $dn_i(t)/dt$ , or sensitivity, of the process, and thus have a negative sign in (6.1). The change  $dn_i(t)/dt$  is positive if the overall mean yield  $Y_i(t)$  for individual agent  $A_{ij}$  in group  $A_i$  is larger than the overall mean yield  $Y(t)$  for an individual agent  $A_{ij}$  in any group  $A_i$ ,  $i = 1, \dots, n$ , and negative if the reverse is the case. Aside from the economic interpretation, (6.1) can be interpreted as the Malthusian law of population growth, where the sign of the proportionality factor  $r_i[Y_i(t) - Y(t)]$  specifies whether the size  $n_i(t)$  of group  $A_i$  (the population) increases or decreases.

To derive the requirement for preserving the total number  $N(t) = N$  of individual agents  $A_{ij}$  in all the  $n$  groups  $A_i$ ,  $i = 1, \dots, n$ , sum both sides in (6.1) from  $i = 1$  to  $i = n$ , i.e.  $\sum_{i=1}^n (\cdot)$ . Inserting (4.6) gives

$$\begin{aligned} \sum_{i=1}^n \frac{dn_i(t)}{dt} &= \frac{d}{dt} \sum_{i=1}^n n_i(t) \\ &= \frac{d}{dt} N(t) = \frac{d}{dt} N = \sum_{i=1}^n r_i n_i(t) [Y_i(t) - Y(t)] = 0. \end{aligned} \quad (6.2)$$

(6.2) can be simplified by setting the ‘sensitivity’ factor  $r_i$  constant across the groups  $A_i$ ,  $i = 1, \dots, n$ , i.e.  $r_i = r$ . Using (4.5) gives

$$\begin{aligned} \sum_{i=1}^n r_i n_i(t) [Y_i(t) - Y(t)] &= r \left( \sum_{i=1}^n n_i(t) Y_i(t) \right) - rY(t) \sum_{i=1}^n n_i(t) \\ &= rNY(t) - rY(t)N = 0. \end{aligned} \quad (6.3)$$



which means that the requirement for preserving the total number  $N(t) = N$  of individual agents  $A_{ij}$  in all the  $n$  groups  $A_i$ ,  $i = 1, \dots, n$ , is automatically satisfied if the ‘sensitivity’ factor  $r_i = r = \text{constant}$ .

The final requirement for (6.1) is that  $0 \leq n_i(t) \leq N \ \forall i = 1, \dots, n$ . I set  $n$  constant so that  $n_i(t) = 0$  means that group  $A_i$  exists but has no members  $A_{ij}$ . If  $n_i(t) = N$  then all the groups  $A_u$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ , except group  $A_i$  has no members  $A_{uj}$ .

Inserting (5.2) and (5.4) into (4.4) allowing for the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ), where  $0 < B_i(t)/n_i(t) < c_i < B_i(t)$ , gives

$$\begin{aligned}
 Y_i(t) &= p_{ic}(t)(B_i(t) - c_i) \\
 &+ \left\{ \left[ \left( \sum_{i=1}^n p_{ic}(t)n_i(t)B_i(t) \right) + \left( 1 - \frac{1}{n_i(t)} + \frac{p_{ic}(t)}{n_i(t)} - 2p_{ic}(t) \right) \right. \right. \\
 &\times B_i(t) \left. \right] p_{ic}(t)n_i(t)B_i(t)E \left. \right\} \\
 &\times \left\{ \left[ \left( \sum_{i=1}^n p_{ic}(t)n_i(t)B_i(t) \right) + (1 - p_{ic}(t))B_i(t) \right] \right. \\
 &\times \left. \left[ \left( \sum_{i=1}^n p_{ic}(t)n_i(t)B_i(t) \right) - p_{ic}(t)B_i(t) \right] \right\}^{-1}, \tag{6.4}
 \end{aligned}$$

defined for  $n_i(t) > 0$ ,  $i = 1, \dots, n$ . (6.1) inserting (6.4) has no straightforward analytical solution. With one available strategy cooperation  $c$  (effort level  $S_{ij} = c_i$ ), giving  $p_{ic}(t) = 1$ , (6.1) inserting (6.4) becomes

$$\begin{aligned}
 \frac{dn_i(t)}{dt} &= r_i \frac{n_i(t)}{N(t)} \\
 &\times \sum_{\substack{u=1 \\ u \neq i}}^n n_u(t) \left[ (B_i(t) - B_u(t) - c_i + c_u) + \frac{(B_i(t) - B_u(t))E}{\sum_{i=1}^n n_i(t)B_i(t)} \right]. \tag{6.5}
 \end{aligned}$$

(6.5) illustrates that with one available strategy cooperation  $c$  there is no inter-group mobility if the parameters in the  $n$  groups  $A_i$ ,  $i = 1, \dots, n$  are equivalent, i.e.  $B_i(t) = B(t)$  and  $c_i = c \ \forall i = 1, \dots, n$ . If  $B_i(t) = B(t)$  then all the individual agents  $A_{ij}$  will, compatible with the principle of competitive exclusion in population biology, eventually seek that group  $A_i$  having the lowest cost  $c_i$  of effort for a cooperative act  $c$ , i.e.  $c_i < c_u \ \forall i, u = 1, \dots, n$ ,  $i \neq u$ , the other groups  $A_u$  becoming extinct in the sense of having no members  $A_{uj}$ . If  $c_i = c$  then all

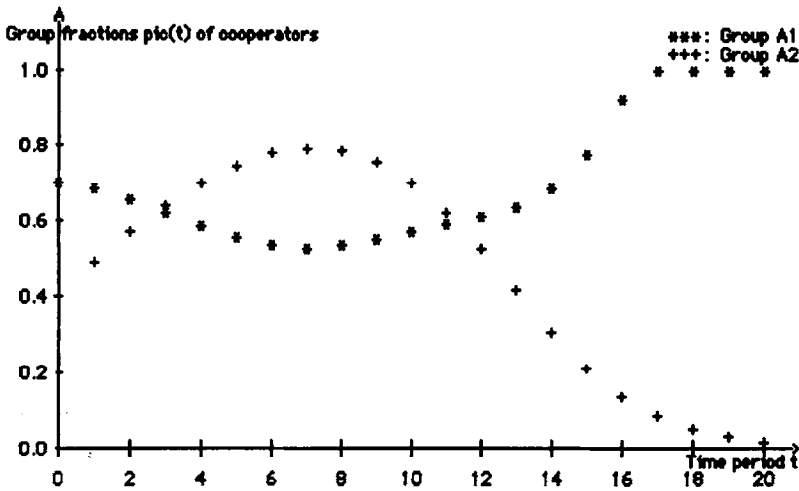


Fig. 3. Evolution of fractions  $p_{ic}(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $r_i = 0.00021$ .

individual agents  $A_{ij}$  will eventually seek that group  $A_i$  that can generate the highest amount of within-group public goods  $B_i(t)$  from a cooperative act  $c$ , the other groups  $A_u$  becoming extinct. Note that this latter process is quicker if  $E$  is large. If  $B_i(t) \neq B(t)$  and  $c_i \neq c$  then all the individual agents  $A_{ij}$  will eventually seek that group  $A_i$  where the bracketed term in (6.5) is positive when two groups  $A_i$  and  $A_u$  remain,  $i, u = 1, \dots, n$ ,  $i \neq u$ . Observe that in this case  $B_i(t) < B_u(t)$  that is not beneficial for individual agent  $A_{ij}$  can be offset by a beneficial  $c_i \ll c_u$ , and vice versa.

Simulate (6.1) inserting (6.4) allowing for the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ). Figs. 3 and 4 assume the same parameters as in Fig. 2, i.e.  $n = 2$ ,  $n_1(0) = 1000$ ,  $B_1(t) = 0.9n_1(t)$ ,  $B_2(t) = 1.1n_2(t)$ ,  $E = 2,700,000$ , initial fractions  $p_{1c}(0) = 0.7$  and  $p_{2c}(0) = 0.4$ , and ‘sensitivity’ factors  $k_i = 0.5$  and  $r_i = 0.00021$ .

Fig. 3 should be compared with Fig. 2. During the first time periods the defectors in the least cooperative group  $A_2$  switch to cooperation at a high rate,  $B_2(t) = 1.1n_2(t)$ , while the cooperators in the most cooperative group  $A_1$  switch to defection at a slightly slower rate,  $B_1(t) = 0.9n_1(t)$ . Because of the higher group fraction  $p_{1c}(t)$  of cooperators in group  $A_1$  initially, group  $A_1$  is considered more attractive, some individual agents  $A_{2j}$  switching from group  $A_2$  to group  $A_1$ , giving  $n_1(t) > n_2(t)$  (cf. Fig. 4). From time period four, the fraction  $p_{2c}(t)$  of cooperators in group  $A_2$  is larger than the fraction  $p_{1c}(t)$  of cooperators in group  $A_1$ , which makes group  $A_2$  more attractive. Hence from time periods four to eight,  $n_1(t)$  decreases and  $n_2(t)$  increases. From time period eight,  $n_2(t) > n_1(t)$ , and the between-group prize  $E$  is no longer large enough to sustain cooperation within the

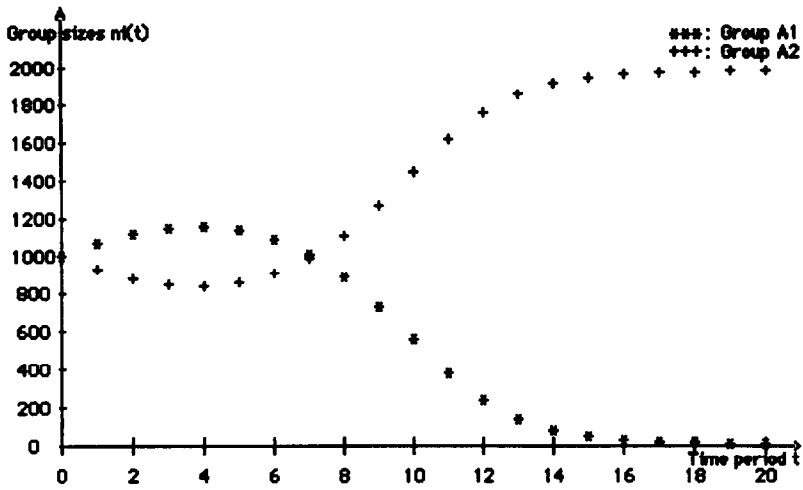


Fig. 4. Evolution of group sizes  $n_i(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $r_i = 0.00021$ .

larger group  $A_2$ . Observe from (5.7) and (3.10) that the requirement for the between-group prize  $E$  gets more restrictive for the larger group  $A_2$  and less restrictive for the smaller group  $A_1$  as the group size  $n_2(t)$  gets increasingly larger than  $n_1(t)$ . The individual agents  $A_{1j}$  in group  $A_1$  thus have an incentive to cooperate, but the small size  $n_1(t)$  makes group  $A_1$  command a relatively small proportion of the between-group prize  $E$ . Hence group  $A_2$  can function more as if in isolation. When group  $A_1$  eventually becomes extinct, no between-group prize  $E$  is large enough to counteract the logic of the within-group  $n_2(t)$ -person prisoner's dilemma, which dictates defection as the strategy that will flourish.

### 7. Dynamic model with intergroup warfare

Describing intergroup warfare means positing a model (4.10) specifying a change  $dn_i(t)/dt$  based on how the groups  $A_i$ ,  $i = 1, \dots, n$ , wage war against each other over time. There is no preservation of the total number  $N(t)$  of individual agents  $A_{ij}$  in the  $n$  groups  $A_i$ . A generalized  $n$ -group version of Lanchester's square law for conventional warfare is

$$\frac{dn_i(t)}{dt} \text{ mod} = -\lambda_i \left[ \sum_{\substack{u=1 \\ u \neq i}}^n n_u(t) Y_u(t) \right] + h_i(t), \tag{7.1}$$

where  $\lambda_i$  is the 'sensitivity' factor for the process,  $h_i(t)$  is the reinforcement rate, and  $0 \leq n_i(t) \leq N(t) \forall i = 1, \dots, n$ . Operational loss due to non-combat factors is

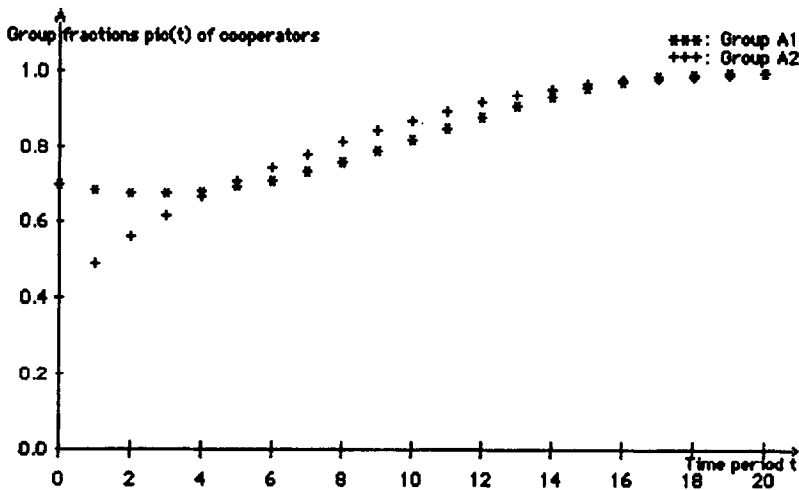


Fig. 5. Evolution of fractions  $p_{ic}(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $\lambda_i = 0.000014$ .

assumed to be negligible.  $Y_u(t)$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ , is given by (4.3) inserting (5.2) for the continuous strategy model and (4.4) inserting (5.2) and (5.4) for the discrete strategy model. Consider  $n$  groups (forces)  $A_i$ ,  $i = 1, \dots, n$ , out in the open, each within kill range of its  $n - 1$  enemies  $A_u$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ . If one group  $A_i$  suffers a loss, fire is concentrated on the remaining individual agents  $A_{ij}$  in group  $A_i$ . This gives the combat loss rate for group  $A_i$  as described by the first term on the right-hand side of (7.1). That is, each group  $A_i$  suffers a loss proportional to the sum of the 'strengths' of the opposing  $n - 1$  groups  $A_u$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ . Note that 'strength' in (7.1) is defined as  $n_u(t)Y_u(t)$  and not  $n_u(t)$ , where  $Y_u(t)$  is given by (4.3). This means that  $dn_i(t)/dt$  is affected not only by the numbers  $n_u(t)$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ , of individual agents  $A_{uj}$  in the other  $n - 1$  groups  $A_u$ , as is commonly done in the simulation of Lanchester's square law, but also by the mean yields  $Y_u(t)$  for these individual agents, as these depend upon the parameters for the groups  $A_u$  and the various strategies  $S_{uj}$  that the individual agents  $A_{uj}$  choose. Hence the different capacities of the various individual agents  $A_{uj}$  are accounted for, as these capacities depend upon the parameters and strategy sets  $S_{uj}$  for the groups  $A_u$ .

(7.1) inserting (6.4) has no straightforward analytical solution. Simulate (7.1) inserting (6.4) allowing for the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ). Figs. 5 and 6 assume the same parameters as in Figs. 2, 3, and 4, i.e.  $n = 2$ ,  $n_i(0) = 1000$ ,  $B_1(t) = 0.9n_1(t)$ ,  $B_2(t) = 1.1n_2(t)$ ,  $E = 2,700,000$ , initial fractions  $p_{1c}(0) = 0.7$  and  $p_{2c}(0) = 0.4$ , reinforcement rate  $h_i(t) = 0$ , and 'sensitivity' factors  $k_i = 0.5$  and  $\lambda_i = 0.000014$ .

Fig. 5 should be compared with Fig. 2. The between-group competition is now more fierce, in the sense that the two groups  $A_i$  and  $A_u$  compete not only for the

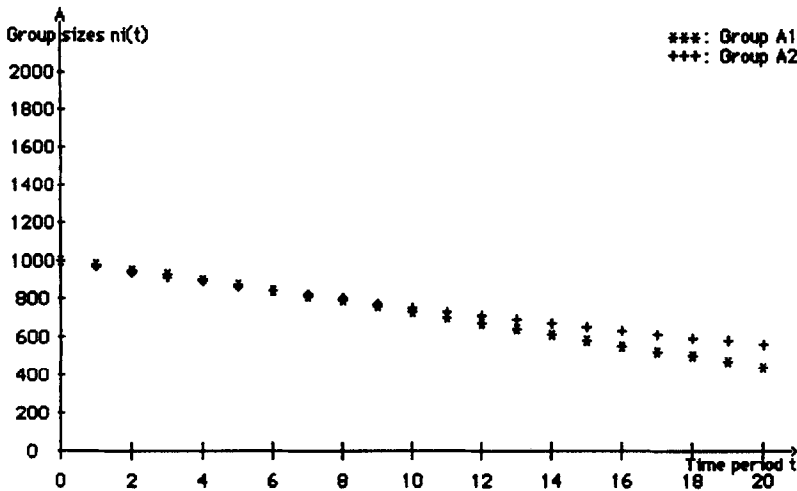


Fig. 6. Evolution of group sizes  $n_i(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $\lambda_i = 0.000014$ .

between-group prize  $E$ ; they also compete to exterminate each other's individual agents  $A_{ij}$  and  $A_{uj}$ . This provides an additional incentive for each individual agent  $A_{ij}$  and  $A_{uj}$ ,  $i \neq u$ ,  $i, u = 1, \dots, n$ , within each group  $A_i$  and  $A_u$  to choose cooperation  $c$ , which increases the fractions  $P_{ic}(t)$  and  $p_{uc}(t)$ , respectively, of cooperators. Conversely, by comparison with Fig. 2, a given fraction  $p_{ic}(t)$  of cooperation within each group can be obtained with a lower between-group prize  $E$ . These results are compatible with the common claim by Hart (1983) and others that more fierce competition in the product-market (i.e. between organizations) reduces managerial slack, i.e. increases managerial cooperation. I will return to this issue in a future article. Fig. 6 illustrates the evolution of the group sizes  $n_i(t)$ , where eventually  $n_1(t) < n_2(t)$  because the parameters  $B_1(t) = 0.9n_1(t)$  and  $B_2(t) = 1.1n_2(t)$  make group  $A_2$  more favorable, more cooperative, and stronger than group  $A_1$ .

Generalizing (7.1) to guerilla warfare, set

$$\frac{dn_i(t)}{dt} \text{ mod} = -\lambda_i \left[ \sum_{\substack{u=1 \\ u \neq i}}^n n_i(t)n_u(t)Y_u(t) \right] + h_i(t). \tag{7.2}$$

Assume that each group  $A_i$ ,  $i = 1, \dots, n$ , occupies a certain geographical region and is invisible to all its opposing  $n - 1$  groups  $A_u$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ . The groups  $A_u$  fire into the region but cannot know when a kill of an individual agent  $A_{ij}$  has been made. As for conventional warfare, each group  $A_i$  suffers a loss proportional to the sum of the 'strengths' of the opposing  $n - 1$  groups  $A_u$ ,  $u \neq i$ ,  $i, u = 1, \dots, n$ . Guerilla warfare assumes in addition that each group  $A_i$  suffers a

loss proportional to the number  $n_i(t)$  of individual agents  $A_{ij}$  in its own group  $A_i$ , since the larger  $n_i(t)$  is, the greater is the probability that a shot from the other  $n - 1$  groups  $A_u$  will kill a member from group  $A_i$ .

### 8. Dynamic model with intergroup mobility and intergroup warfare

The combined operation of intergroup mobility and intergroup warfare can be modeled combining (6.1) and e.g. a generalized version of (7.2), giving

$$\frac{dn_i(t)}{dt} \text{ mod} = r_i n_i(t) \left[ Y_i(t) - \frac{1}{N(t)} \sum_{i=1}^n n_i(t) Y_i(t) \right] - \lambda_i \left[ \sum_{\substack{u=1 \\ u \neq i}}^n \sum_{x=0}^{x_{\max}} \sum_{y=1}^{y_{\max}} n_u^x(t) n_u^y(t) Y_u^y(t) \right] + h_i(t), \tag{8.1}$$

where  $n_i(t)n_u(t)Y_u(t)$  in (7.2) is substituted with a Taylor series, and where  $N(t)$  is no longer a constant. The ‘sensitivity’ factors  $r_i$  and  $\lambda_i$  determine the relative weight of intergroup mobility and intergroup warfare.

(8.1) inserting (6.4) has no straightforward analytical solution. Simulate (8.1) inserting (6.4) allowing for the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ). Figs. 7 and 8 assume the same parameters as in Figs. 2, 3, and 4, i.e.  $n = 2$ ,  $n_i(0) = 1000$ ,  $B_1(t) = 0.9n_1(t)$ ,  $B_2(t) = 1.1n_2(t)$ ,  $E = 2,700,000$ ,  $x_{\max} = 0$ ,  $y_{\max} = 1$ , initial fractions  $p_{1c}(0) = 0.7$

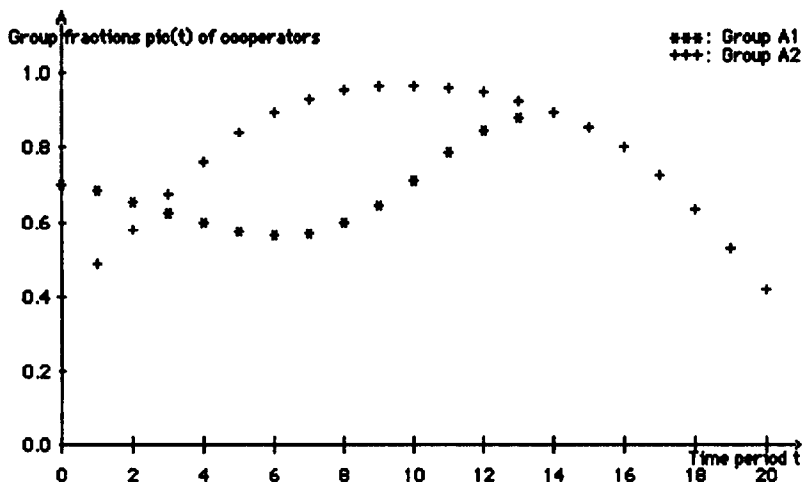


Fig. 7. Evolution of  $p_{1c}(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $r_i = 0.00021$ ,  $\lambda_i = 0.000014$ .

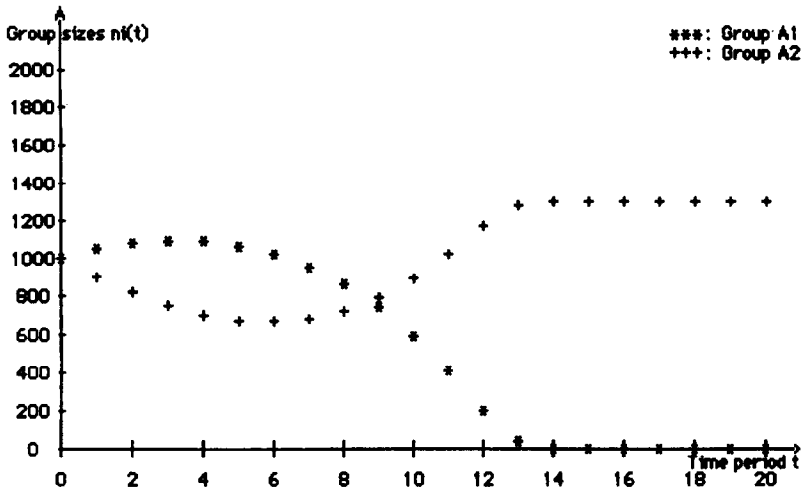


Fig. 8. Evolution of  $n_i(t)$ ,  $E = 2,700,000$ ,  $k_i = 0.5$ ,  $r_i = 0.00021$ ,  $\lambda_i = 0.000014$ .

and  $p_{2c}(0) = 0.4$ , reinforcement rate  $h_i(t) = 0$ , and ‘sensitivity’ factors  $k_i = 0.5$ ,  $r_i = 0.00021$ , and  $\lambda_i = 0.000014$ .

Figs. 7 and 8 illustrate the combined operation of intergroup mobility from Section 6 and intergroup warfare. The intergroup warfare makes  $n_1(t)$  approach zero faster than when exclusively assuming intergroup mobility. As  $n_1(t)$  approaches zero, group  $A_2$  can function more in isolation, implying a decline in  $p_{2c}(t)$  eventually giving  $p_{2c}(\infty) = 0$ .  $n_1(14) = 0$  makes  $p_{1c}(t)$  undefined for  $n_1 \geq 14$ , while  $n_2(t) = 1304$  for  $t \geq 14$ .

### 9. Simulation of the algorithm with four strategies in one group

This section shows that the introduction of the  $n_i(t)$ -person reactive strategies, tit-for-tat and bully, into the within-group one-group  $n_i(t)$ -person prisoner’s dilemma does not facilitate within-group cooperation in the way the between-group competition for an external prize  $E$  does. Define the two reactive strategies as follows:

1.  $T_e$ ; cooperate in the initial time period  $t = 0$ , and then cooperate in each subsequent time period  $t$ ,  $t = 1, 2, \dots$  if  $e$  or more of the other  $n_i(t) - 1$  individual agents in group  $A_i$  chose cooperation  $c$  during the previous time period  $t - 1$  (Taylor, 1976).
2.  $G_e$ ; defect in the initial time period  $t = 0$ , and then defect in each subsequent time period  $t$ ,  $t = 1, 2, \dots$  if  $e$  or more of the other  $n_i(t) - 1$  individual agents in group  $A_i$  chose cooperation  $c$  during the previous time period  $t - 1$ .

Inserting (5.4) into (5.2) allowing for the two discrete strategies cooperation  $c$  (effort level  $S_{ij} = c_i$ ) and defection  $d$  (effort level  $S_{ij} = 0$ ), where  $0 < B_i(t)/n_i(t) < c_i < B_i(t)$ , gives

$$\begin{aligned}
 Y_{ic}(t) = & \left( \left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) + \frac{1}{n_i(t)} \right) B_i(t) \\
 & + \frac{\left( \left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) + \frac{1}{n_i(t)} \right) B_i(t) E}{\left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) + (1 - p_{ic}(t)) B_i(t) \right]} - c_i \quad (9.1)
 \end{aligned}$$

and

$$\begin{aligned}
 Y_{id}(t) = & \left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) B_i(t) \\
 & + \frac{\left( 1 - \frac{1}{n_i(t)} \right) p_{ic}(t) B_i(t) E}{\left[ \left( \sum_{i=1}^n p_{ic}(t) n_i(t) B_i(t) \right) - p_{ic}(t) B_i(t) \right]}. \quad (9.2)
 \end{aligned}$$

The payoffs  $Y_{iT}(t)$  and  $Y_{iG}(t)$  for playing the reactive strategies  $T_e$  and  $G_e$  are determined from (9.1) and (9.2) based on the definitions of  $T_e$  and  $G_e$  given above, i.e. depending upon the parameter  $e$  related to how many of the other  $n_i(t) - 1$  individual agents in group  $A_i$  chose cooperation  $c$  during the previous time period. Set  $n_i(t) = 1000$ ,  $B_i(t) = n_i(t)$ ,  $c_i = 2$ , ‘sensitivity’ factor  $k_i = 0.5$ , and simulate over 20 time periods with initial conditions  $p_{ic}(0) = 0.7$ ,  $p_{id}(0) = p_{iT}(0) = p_{iG}(0) = 0.1$ . Fig. 9 sets  $e = n_i(t) - 1$ . In the initial time period  $t = 0$ ,  $d$  and  $G_e$  go to prey upon  $c$  and  $T_e$ . In time period  $t = 1$ ,  $G_e$  switches to cooperation and  $T_e$  switches to defection. Hence thereafter  $d$  and  $T_e$  prey upon  $c$  and  $G_e$ , the latter two becoming extinct. If  $T_e$  defects in the initial time period  $t = 0$ ,  $T_e$  and  $d$  coincide throughout giving  $p_{id}(\infty) = p_{iT}(\infty) = 0.5$ .

Fig. 10 decreases  $e$  to  $e = 0.7n_i(t)$ .  $G_e$  defects and  $T_e$  cooperates through time period  $t = 3$ , after which they switch roles, giving a similar evolution as Fig. 9. Fig. 11 decreases  $e$  further to  $e = 0.1n_i(t)$ .  $G_e$  defects and  $T_e$  cooperates through time period  $t = 8$ , after which they switch roles. Since  $T_e$  is almost extinct at  $t = 8$ ,  $d$  thereafter preys upon  $G_e$ , the latter becoming extinct,  $d$  approaching  $p_{id}(\infty) = 1$ . With  $e = 0$ ,  $d$  and  $G_e$  coincide throughout giving  $p_{id}(\infty) = p_{iG}(\infty) = 0.5$ .

Figs. 9–11 show the superiority of defection  $d$ . The  $n_i(t)$ -person tit-for-tat strategy  $T_e$  is not as successful as the two-person tit-for-tat strategy (Axelrod, 1984), mainly because of the constellation of one against many in  $n_i(t)$ -person



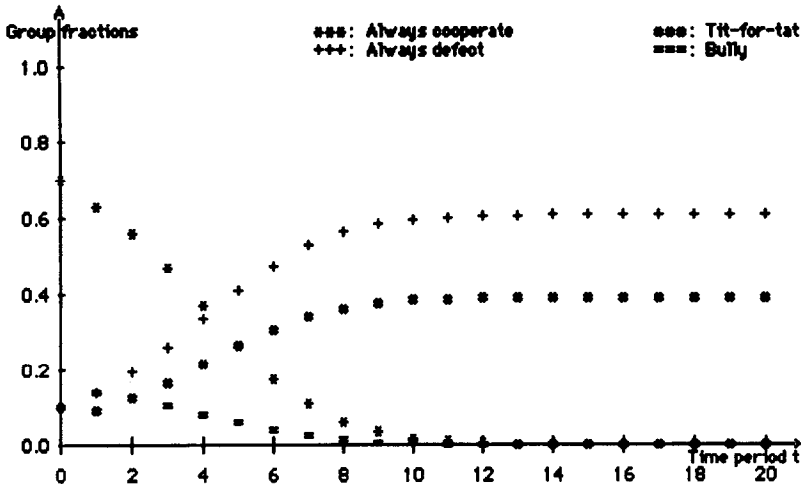


Fig. 9. Evolution of fractions  $p_{i_s}(t)$  for  $e = n_i(t) - 1$  for one group ( $E = 0$ ),  $k_i = 0.5$ .

games, not present in two-person games, and the rigidity of the parameter  $e$ . This can be partly remedied by assuming a ‘continuous’ version of an  $n_i(t)$ -person prisoner’s dilemma which, in addition to complicating the model, allows individual agent  $A_{ij}$  to choose a degree  $S_{ij}$ ,  $0 \leq S_{ij} \leq c_j$ , of cooperation based upon the overall degree  $S_i$  of cooperation in group  $A_i$  in the previous time period. This remedy is insufficient because defectors are not punished with sufficient degree of specificity, the punishment being spread out over all the other  $n_i(t) - 1$

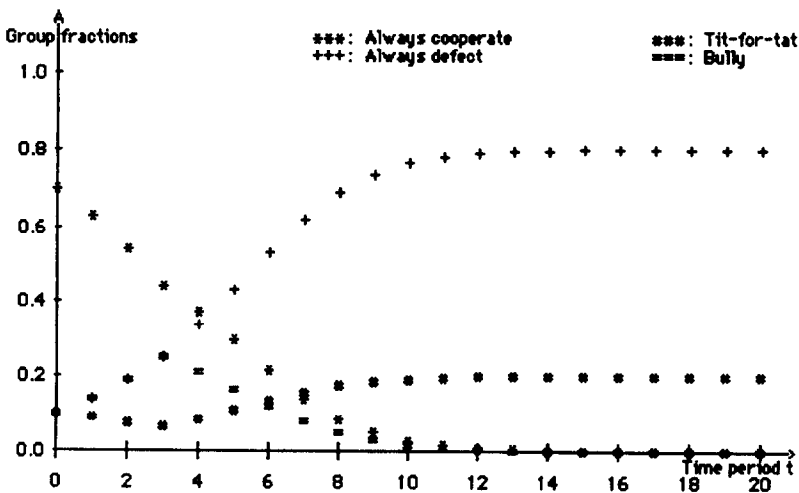


Fig. 10. Evolution of fractions  $p_{i_s}(t)$  for  $e = 0.7n_i(t)$  for one group ( $E = 0$ ),  $k_i = 0.5$ .

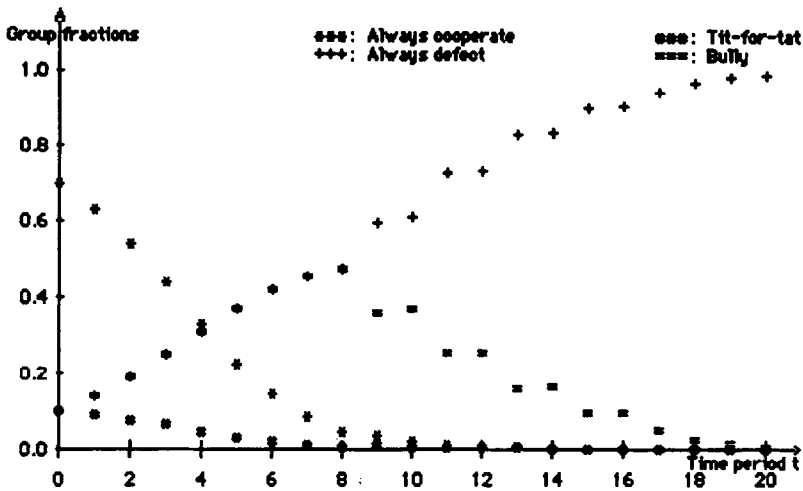


Fig. 11. Evolution of fractions  $p_{is}(t)$  for  $e = 0.1n_i(t)$  for one group ( $E = 0$ ),  $k_i = 0.5$ .

individual agents in group  $A_i$ . A more appropriate remedy considered in this article involves introducing an external structure in the form of between-group  $n$ -group competition for a specific prize  $E$ .

**Appendix: Between-group competition where the more cooperative group  $A_i$  gets the entire between-group prize  $E$**

This article has assumed between-group competition where each group  $A_i$  gets a fraction of the between-group prize  $E$  proportional to the relative amount of public goods it produces compared with the other groups. An alternative is to assume between-group competition where the group  $A_i$  producing the highest amount of public goods gets the entire between-group prize  $E$ , i.e. ‘winner takes all’, and that if  $m$  of the  $n$  groups produce an equally high amount of public goods, that these share the prize  $E$  equally. Assume for expositional convenience  $B_i = B_1$ ,  $c_i = c_1$ , and  $a_i = 1$ . Assume without loss of generality that  $n_1 = \dots = n_m > \dots \geq n_n$ , where  $1 \leq m \leq n \geq 2$ .

Assume the within-group  $n_i$ -person prisoner’s dilemma from Section 3 where  $0 < B_1/n_i < c_1 < B_1$ . Define  $h_{1c} - 1$  as the number of cooperating individual agents aside from individual agent  $A_{1j}$  in group  $A_1$ . If individual agent  $A_{1j}$  chooses  $S_{1j} = c_1$ , giving  $h_{1c}$  cooperators in group  $A_1$ , he gets, analogous to (3.7), a payoff

$$P_{1jc} = \frac{1}{n_1} [h_{1c}B_1 + E] - c_1 \text{ if } h_{1c} > h_{2c} \geq \dots \geq h_{nc}$$

$$\begin{aligned}
 &= \frac{1}{n_1} \left[ h_{1c} B_1 + \frac{E}{m} \right] - c_1 \text{ if } h_{1c} = \dots = h_{mc} > \dots \geq h_{nc} \\
 &= \frac{1}{n_1} \left[ h_{1c} B_1 + \frac{E}{n} \right] - c_1 \text{ if } h_{1c} = \dots = h_{nc} \\
 &= \frac{1}{n_1} h_{1c} B_1 - c_1 \text{ if } h_{1c} < h_{kc} \text{ for at least one } k, k = 2, \dots, n.
 \end{aligned}
 \tag{A.1}$$

If individual agent  $A_{ij}$  chooses  $S_{ij} = 0$ , there are  $h_{ic} - 1$  cooperators in group  $A_j$ . Assume that  $h_{uc}$ ,  $u \neq i$ ,  $u = 1, \dots, n$ , is the number of cooperators in the other  $n - 1$  groups  $A_u$ . Individual agent  $A_{ij}$  then gets, analogous to (3.8), a payoff

$$\begin{aligned}
 P_{1jd} &= \frac{1}{n_1} [(h_{1c} - 1) B_1 + E] \text{ if } h_{1c} - 1 > h_{2c} \geq \dots \geq h_{nc} \\
 &= \frac{1}{n_1} \left[ (h_{1c} - 1) B_1 + \frac{E}{m} \right] \text{ if } h_{1c} - 1 = \dots = h_{mc} > \dots \geq h_{nc} \\
 &= \frac{1}{n_1} \left[ (h_{1c} - 1) B_1 + \frac{E}{n} \right] \text{ if } h_{1c} - 1 = \dots = h_{nc} \\
 &= \frac{1}{n_1} (h_{1c} - 1) B_1 \text{ if } h_{1c} - 1 < h_{kc} \text{ for at least one } k, k = 2, \dots, n.
 \end{aligned}
 \tag{A.2}$$

When  $m$  of the  $n$  groups are equally large, where  $2 \leq m \leq n$ , then the most restrictive requirement for  $E$  to ensure cooperation by all the individual agents in the  $m$  groups arises when everyone in the  $m$  groups cooperates, and one individual agent  $A_{1j}$  in group  $A_1$  contemplates whether to defect. To derive this requirement, set  $h_{1c} = n_1$  in (A.1) and (A.2). Cooperation by the individual agent  $A_{1j}$  is guaranteed if (A.1) is larger than (A.2), i.e. if

$$\frac{1}{n_1} \left[ n_1 B_1 + \frac{E}{m} \right] - c_1 > \frac{1}{n_1} (n_1 - 1) B_1,
 \tag{A.3}$$

i.e. if

$$E > mn_1 \left( c_1 - \frac{B_1}{n_1} \right) \text{ if } 2 \leq m \leq n.
 \tag{A.4}$$

If  $m = 1$ , which means that  $n_1 > n_2$ , then the individual agents in group  $A_1$  will coordinate such that  $h_{1c} = n_2 + 1$ , given that the between-group prize  $E$  is sufficiently large, which precludes any other group getting access to  $E$ . To derive the most restrictive requirement for  $E$ , assume  $n_2 > n_3$ , set  $h_{1c} = n_2 + 1$  in (A.1)

and (A.2). Cooperation by the individual agent  $A_{1j}$  is guaranteed if (A.1) is larger than (A.2), i.e. if

$$\frac{1}{n_1} [(n_2 + 1)B_1 + E] - c_1 > \frac{1}{n_1} \left[ n_2 B_1 + \frac{E}{2} \right], \quad (\text{A.5})$$

i.e. if

$$E > 2n_1 \left( c_1 - \frac{B_1}{n_1} \right) \text{ if } m = 1. \quad (\text{A.6})$$

(A.6) has a very nice interpretation. The requirement for having a within-group prisoner's dilemma is  $0 < B_1/n_i < c_1 < B_1$ . For the largest group  $A_1$  the expression  $c_1 - B_1/n_1$  thus gives the extent to which the cost  $c_1$  of cooperation for individual agent  $A_{1j}$  in the largest group  $A_1$  exceeds the lower limit  $B_1/n_1$  where the within-group game  $A_1$  goes from being a within-group prisoner's dilemma to becoming a within-group pure cooperation game. This expression  $c_1 - B_1/n_1$  is then multiplied with twice the number of individual agents  $A_{1j}$  in the largest group  $A_1$ . This gives the minimum value of  $E$  required to ensure cooperation by  $h_{1c} = n_2 + 1$  individual agents  $A_{1j}$  in the largest group  $A_1$ .

Observe that inserting  $n_1 = 2$  into (A.6) gives

$$E > 4(c_1 - B_1/2), \quad (\text{A.7})$$

which is three times less restrictive than

$$E > 12(c_1 - B_1/2), \quad (\text{A.8})$$

which follows from inserting  $n = n_i = 2$ ,  $B_i = B_1$ , and  $c_i = c_1$  into (3.10), for the proportional allocation rule.

(A.4) gets more restrictive as the number  $m$  of those largest groups having equal size increases, because the fraction of  $E$  available to each group then gets smaller, and may get too small to facilitate within-group cooperation. If  $E$  represents a fixed amount of resources in the environment, (A.4) may be realistic. However, if the nature of the between-group prize  $E$  is such that it increases with the number  $m$ , (A.4) may be too restrictive. More specifically, if the  $m$  groups compete for a between-group prize  $E$  that is proportional to  $m$ , and if we, in order to illustrate compatibility with the two-group case, choose a proportionality factor such that the  $m$  groups compete for a between-group prize  $mE/2$ , then (A.4) becomes (A.6).

Using (A.4) and (A.6) we can specify six different intervals for the between-group prize  $E$ , and what degree of cooperation within the  $n$  groups corresponds:

1.  $E > mn_1(c_1 - B_1/n_1)$  and  $m \geq 2$ :  $n_1$  cooperators in groups  $A_1$  to  $A_m$  (all remaining are defectors).

Table A1  
Game with two defectors in the other group

		Individual agent 2	
		c	d
Individual agent 1	c	$B_1 - c_1 + E/2, B_1 - c_1 + E/2$	$B_1/2 - c_1 + E/2, B_1/2 + E/2$
	d	$B_1/2 + E/2, B_1/2 - c_1 + E/2$	$E/4, E/4$

2.  $(m - q)n_1(c_1 - B_1/n_1) < E \leq (m - q + 1)(c_1 - B_1/n_1)$  and  $m \geq 3$ :  $n_1$  cooperators in groups  $A_1$  to  $A_{m-q}$ ,  $q = 1, \dots, m - 2$ .
3.  $E > 2n_1(c_1 - B_1/n_1)$  and  $m = 1$ :  $n_2 + 1$  cooperators in group  $A_1$ .
4.  $2n_{m+1}(c_1 - B_1/n_{m+1}) < E \leq 2n_1(c_1 - B_1/n_1)$ :  $n_{m+2} + 1$  cooperators in group  $A_{m+1}$ .
5.  $2n_n(c_1 - B_1/n_n) < E \leq 2n_{n-1}(c_1 - B_1/n_{n-1})$ : 1 cooperator in group  $A_1$ .
6.  $E \leq 2n_n(c_1 - B_1/n_n)$ : no cooperators.

In the remainder of the appendix I analyze in detail the case of two groups  $A_i$  with two individual agents  $A_{ij}$  in each group. An individual agent  $A_{ij}$  gets from between-group competition a payoff  $E/2$  if he is part of the more cooperative group, and a payoff  $E/4$  if the two groups are equally cooperative. The payoff matrix for the individual agents in each group, given that the two individual agents in the other group defect, is as in Table A1.

Table A1 is derived from the basic two-person prisoner’s dilemma by adding the payoff  $E/2$  to each entry except the two entries in the lower right corner, where the payoff  $E/4$  is added. Table A1 shows a chicken game and not a prisoner’s dilemma when

$$B_1/2 - c_1 + E/2 > E/4. \tag{A.9}$$

Given this, and given that there are four defectors in the other group, any defector has an incentive to switch to cooperation to increase his payoff.

Given that there are one cooperator and one defector in the other group, the payoff matrix for a group is as in Table A2.

Table A2 is derived from the basic two-person prisoner’s dilemma by adding the payoff  $E/2$  to the upper left entries, adding the payoff  $E/4$  to the entries on

Table A2  
Game with one cooperator and one defector in the other group

		Individual agent 2	
		c	d
Individual agent 1	c	$B_1 - c_1 + E/2, B_1 - c_1 + E/2$	$B_1/2 - c_1 + E/4, B_1/2 + E/4$
	d	$B_1/2 + E/4, B_1/2 - c_1 + E/4$	0, 0

Table A3

Game with two cooperators in the other group

		Individual agent 2	
		c	d
Individual agent 1	c	$B_1 - c_1 + E/4, B_1 - c_1 + E/4$	$B_1/2 - c_1, B_1/2$
	d	$B_1/2, B_1/2 - c_1$	0, 0

the diagonal from lower left to upper right, and keeping the lower right entries unchanged. Table A2 shows a mutual cooperation game when

$$B_1 - c_1 + E/2 > B_1/2 + E/4. \quad (\text{A.10})$$

Given this, if cooperation gets started in one group, it will be perpetuated to mutual cooperation in the other group.

Given that there are two cooperators in the other group, the payoff matrix for a group is as in Table A3. Table A3 is derived from the basic two-person prisoner's dilemma by adding the payoff  $E/4$  to the upper left entries, and keeping the other entries unchanged. Table A3 shows a cooperation game when (A.10) is satisfied. Given this, the individual agents can be expected to coordinate their actions so that the unique mutual cooperation equilibrium  $(B_1 - c_1 + E/4, B_1 - c_1 + E/4)$  is attained.

Both (A.9) and (A.10) imply that

$$E > 4(c_1 - B_1/2), \quad (\text{A.11})$$

which is thus the requirement for the between-group prize  $E$  to ensure the unique mutual cooperation equilibrium developed in this appendix.

(A.11) has a very nice interpretation. The requirement for having a within-group prisoner's dilemma is  $0 < B_1/2 < c_1 < B_1$ . The expression  $c_1 - B_1/2$  gives the extent to which the cost  $c_1$  of cooperation for an individual agent exceeds the lower limit  $B_1/2$  where the within-group game goes from being a within-group prisoner's dilemma to becoming a within-group pure cooperation game. This expression  $c_1 - B_1/2$  is then multiplied with the total number, four, of individual agents in the two groups. This gives the minimum value of  $E$  required to ensure cooperation by all four individual agents in the two groups.

Tables A1-A3 can be illustrated with the four-dimensional hypercube in Fig. A1, where  $c_1 = 3$  and  $B_1 = 4$  dictate that we have a within-group prisoner's dilemma, while  $E = 8$  nevertheless guarantees that the unique mutual cooperation equilibrium is attained.

Fig. A1 illustrates how the between-group assumption of a value  $E$  to be competed for between the two groups turns the basic within-group prisoner's

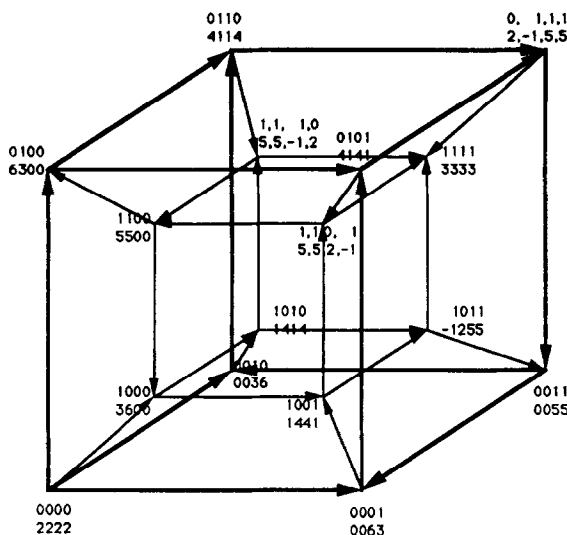


Fig. A1. Four-dimensional hypercube for payoffs to four individual agents divided into two groups.

dilemma into a game with mutual cooperation as the unique equilibrium. The hypercube in Fig. A1 has 16 nodes, corresponding to the 16 possible combinations of cooperation and defection by the four individual agents in the two groups. Each node has eight numbers attached to it, which are to be read as follows. The four numbers in the upper row are zero or one, corresponding to defection and cooperation, respectively, by the four individual agents. The two leftmost numbers in the upper row are for the two individual agents in what we may call group  $A_1$ , and the two rightmost numbers in the upper row are for the two individual agents in group  $A_2$ . The four numbers in the lower row are payoffs to the four individual agents, when the combination of cooperation and defection by the four individual agents is as specified in the upper row.

All nodes adjacent to each other in Fig. A1 are such that only one of the four individual agents switches strategy, either from cooperation to defection, or from defection to cooperation. A switch of strategy will occur if a higher payoff is attained by switching strategy, which is illustrated in Fig. A1 with an arrow in the direction of the higher payoff. There are thus 32 arrows in Fig. A1.

We observe in Fig. A1 that there is only one node that has four arrows coming into it, corresponding to the unique mutual cooperation equilibrium (1, 1, 1, 1) with payoff (3, 3, 3, 3). There is also only one node with four arrows going out from it, corresponding to mutual defection (0, 0, 0, 0) with payoff (2, 2, 2, 2), which is thus as far from the equilibrium as one can get. The quickest route from mutual defection (0, 0, 0, 0) to mutual cooperation (1, 1, 1, 1) goes through four steps, and involves a zigzagging movement, i.e. a switch of strategy alternately by an individual agent from each of the two groups.

We also observe that regardless of the node in Fig. A1 in which one is located initially, one will eventually end up in the unique mutual cooperation equilibrium  $(1, 1, 1, 1)$ , with one exception. There are five interconnected circles in Fig. A1, all going through four nodes, and all going through the two nodes  $(1, 1, 0, 0)$  and  $(0, 0, 1, 1)$  and thereafter one of the four chicken equilibria in Table A1. Two of the circles are on the inner cube, one is on the outer cube, and two are between the two cubes. These circles open up the possibility of infinite circular movement. No node along any of the five circles corresponds to a final resting point. The circular movements along all five circles can be broken at any point in time when at least three of the four individual agents cooperate, which happens at one of the four nodes in each of the five circles. Once any of the circles is broken, the unique mutual cooperation equilibrium is attained, from which no single individual agent has any incentive to deviate unilaterally.

One may impose a simple restriction to avoid the possibility of circular movement altogether. Fig. A1 shows that when three of the four individual agents cooperate, the single defector improves his payoff from two to three by switching to cooperation, while his companion cooperator, if he has a shorter reaction time to switch, improves his payoff from minus one to zero by switching to defection. These two gains are of equal magnitude one. However, the first switch is preferable from a utilitarian point of view, because it involves an increase in the overall payoff to the two groups from 11 to 12, whereas the second switch involves a decrease in the overall payoff to the two groups from 11 to 10. Hence if indifference is involved, and the cultural characteristics are such that switches from cooperation to defection occur faster and more often than switches from defection to cooperation, a utilitarian argument resolves the indifference. This restriction is not consistent with methodological individualism, and may well be abandoned.

The conclusion of this example is that simple competition between two groups for a between-group prize of value  $E$  turns a static within-group prisoner's dilemma into a game, where mutual cooperation by all the individual agents in the two groups is a unique equilibrium, from which no single individual agent thus has any incentive to deviate unilaterally. That is, a within-group mixed-motive game is turned into a total pure coordination game for the four individual agents. This unique mutual cooperation equilibrium is reachable from any combination of strategies initially played by the individual agents in the two groups.

Let us make some considerations of the size of the between-group prize  $E$  compared with the parameters  $B_1$  and  $c_1$  of the within-group prisoner's dilemmas. If the prisoner's dilemmas are 'intermediately severe', say  $c_1 = 3B_1/4$ , then (A.11) becomes  $E > B_1$ , which means that pure cooperation by the four individual agents  $A_{ij}$  is a unique equilibrium if the between-group prize  $E$  is larger than the public goods,  $B_1$ , produced by one cooperative act  $c$  by one individual agent  $A_{ij}$ . If the prisoner's dilemmas are 'maximally severe' in the sense of bordering towards pure defection games, i.e.  $c_1 = B_1$ , then (A.11) becomes  $E > 2B_1$ , which



means that pure cooperation by the four individual agents  $A_{ij}$  is a unique equilibrium if the between-group prize  $E$  is larger than the public goods,  $2B_1$ , produced by two individual agents  $A_{ij}$  choosing cooperative acts  $c$ . If the prisoner's dilemmas are 'minimally severe' in the sense of bordering towards pure cooperation games, i.e.  $c_1 = B_1/2$ , then (A.11), of course, becomes  $E > 0$ , which means a minuscule requirement for the between-group prize  $E$  to ensure a unique cooperation equilibrium for the four individual agents  $A_{ij}$ .

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