

Fast detection and automatic parameter estimation of a gravitational wave signal with a novel method

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The detection of gravitational wave usually requires to match the measurement data with a large number of templates, which is computationally very expensive. Compressed sensing methods allow one to match the data with a small number of templates and interpolate the rest. However, the interpolation process is still computationally expensive. In this article, we designed a novel method that only requires to match the data with a few templates, yet without needing any interpolation process. The algorithm worked well for signals with relatively high SNRs. It also showed promise for low SNRs signals.

Introduction—While gravitational wave (GW) signals contain invaluable physical information, extracting this information from the noisy data is quite challenging. Most of the time, GW signals are weaker than the instrumental noise at any instant, but they are predictable and long lived [1]. This gives a way to build up signal-to-noise ratio (SNR) over time by tracking the signals coherently with matched filtering [2]. However, this requires the templates to be exactly the same as the true signal to recover the optimal SNR, or at least resemble the true signal sufficiently in order not to lose much SNR [3]. Since the template waveforms depend on several parameters, one needs to match the data with a huge number of templates in the high dimensional parameter space. Therefore, a normal grid-based search is usually computationally extremely expensive, or even prohibitive. The reduction of the computational cost lies in the center of the modern GW data analysis.

There are several categories of algorithms, successfully reducing the computational cost, such as reduced bases (RB) [4], singular value decomposition (SVD) [5] and principal component analysis (PCA) [6]. These methods make use of the fact that each template is strongly correlated with the templates in its neighbourhood in the parameter space. Therefore, its SNR can be effectively interpolated from the SNRs of the templates in its neighbourhood. In other words, the likelihood surface on the grid of the template bank has special properties (sparsity), which allows the compressed sensing [7] algorithms to apply. Instead of using all the templates in the bank, one only needs to calculate the SNRs of a few so-called basis templates (which are different from the original templates), and then interpolate the SNRs of all the other templates in the bank. It is extremely fast to perform matched filtering on that few basis templates comparing to the original bank of templates. However, the interpolation (or sometimes referred to as the reconstruction) process is still computationally expensive.

We wish to design a novel method, which requires to perform matched filtering on a few templates, and in the meantime does not require any interpolation stage (or can automatically reconstruct the parameters of the GW

signal). However, this method currently requires a relatively high SNR of the signal. The detailed description of the method and the preliminary simulation results are shown in the following.

GW data analysis routine—First of all, we briefly review the convention and notations of the GW data analysis. Usually, the measurement data can be expressed as $s = Ah_* + n$, where n is the noise, A is the amplitude of the signal, h_* is the normalized signal in the measurement, which satisfies $\langle h_* | h_* \rangle = 1$. The inner product of two time series $a(t)$ and $b(t)$ is defined as follows

$$\langle a | b \rangle = \int_{-\infty}^{\infty} \frac{\tilde{a}^*(f)\tilde{b}(f)}{S_n(f)} df, \quad (1)$$

where $\tilde{a}(f), \tilde{b}(f)$ are the Fourier transforms of $a(t)$ and $b(t)$. $S_n(f)$ is the so-called two-sided noise power spectral density (PSD), usually defined as $E[\tilde{n}^*(f')\tilde{n}(f)] = S_n(f)\delta(f - f')$.

The GW data analysis problem that we want to solve is formulated as follows. For a set of normalized candidate templates $h_i = h(\Theta_i)$ (we choose the template index $i = 1, \dots, 2^N$ for convenience) characterized by parameters Θ_i , we want to determine which one is present in the measurement, hence obtaining the parameters Θ_* of the signal. Notice that Θ denotes a set of waveform parameters. For clarity, we require the templates to be nearly independent $\langle h_i | h_j \rangle \ll 1, (i \neq j)$. This is not generally true for a whole template bank. However, one can easily divide the entire template bank into a group of smaller template banks, within which the templates are nearly independent.

We assume that the true signal h_* belongs to the template family, $* \in \{1, 2, \dots, 2^N\}$. The inner product between the measurement data and a template is denoted as

$$x_i \equiv \langle s | h_i \rangle = A \langle h_* | h_i \rangle + \langle n | h_i \rangle, \quad (2)$$

thus the expectation and the variance are

$$E(x_i) = A\delta_{*,i} \quad (3)$$

$$\begin{aligned} \text{Var}(x_i) &= E[\langle h_i|n\rangle\langle n|h_i\rangle] \\ &= \langle h_i|h_i\rangle = 1. \end{aligned} \quad (4)$$

By identifying the largest inner product x_* , we can detect the signal h_* and estimate its parameters Θ_* . When the inner product x_* is much larger than its standard deviation $\sqrt{\text{Var}(x_*)} = 1$, the significance is high. The above shows a normal search strategy, which requires to perform 2^N inner products.

The novel method—In the following, we will describe a novel search algorithm. First, we express the waveform indices i in binary, hence each index is an N -digit binary number (e.g. 001011011...). Then, we define N sets \mathcal{P}_k ($k = 1, 2, \dots, N$) such that \mathcal{P}_k consists of all the indices i whose k -th digit is 1. A new template family is defined based on these sets

$$H_k = \sum_{i \in \mathcal{P}_k} h_i. \quad (5)$$

The inner products of these new templates with the measurement data are

$$\begin{aligned} X_k &\equiv \langle s|H_k\rangle \\ &= \sum_{i \in \mathcal{P}_k} \langle s|h_i\rangle. \end{aligned} \quad (6)$$

The expectation of X_k is

$$E(X_k) = \begin{cases} A, & * \in \mathcal{P}_k \\ 0, & * \notin \mathcal{P}_k \end{cases} \quad (7)$$

The variance can be calculated as follows

$$\begin{aligned} \text{Var}(X_k) &= E[\langle n|\sum_{i \in \mathcal{P}_k} h_i\rangle^2] \\ &= \sum_{i,j \in \mathcal{P}_k} \langle h_i|h_j\rangle. \end{aligned} \quad (8)$$

Since the templates h_i are nearly independent, we have

$$\begin{aligned} \text{Var}(X_k) &= \sum_{i \in \mathcal{P}_k} \langle h_i|h_i\rangle \\ &= 2^{N-1}. \end{aligned} \quad (9)$$

Suppose $* \in \mathcal{P}_a$ and $* \notin \mathcal{P}_b$, then

$$E(X_a - X_b) = A \quad (10)$$

$$\begin{aligned} \text{Var}(X_a - X_b) &= E[\langle n|\sum_{i \in \mathcal{P}_a} h_i - \sum_{j \in \mathcal{P}_b} h_j\rangle^2] \\ &= \sum_{i \in \{\mathcal{P}_a \cup \mathcal{P}_b - \mathcal{P}_a \cap \mathcal{P}_b\}} \langle h_i|h_i\rangle \\ &= 2^{N-1}. \end{aligned} \quad (11)$$

When the expectation A is much larger than the standard deviation $2^{(N-1)/2}$, we can set some threshold \mathcal{T} between

A and $2^{(N-1)/2}$. Based on this threshold, a binary number can be obtained as follows: if $X_k > \mathcal{T}$, the k -th bit of this binary number is 1, otherwise its k -th digit is set as 0. This binary number can be converted to a decimal number i_0 . The method identifies the waveform h_{i_0} with parameters Θ_{i_0} to be most probably present in the data. In this new approach, we have used N templates instead of 2^N templates to detect the signal and estimate its parameters. The computational cost is thus reduced from $\mathcal{C} \cdot 2^N$ to $\mathcal{C} \cdot N$. Notice that, if each inner product of the data and a template provides one bit of information (above or below a certain threshold), N is the minimum required number of templates to distinguish 2^N sets of candidate parameters.

Simulation—To exemplify the performance of the novel method, we consider the following chirp waveform family

$$h(t; f, \dot{f}) = \mathcal{A} \cos(2\pi ft + \pi \dot{f} t^2), \quad (12)$$

where \mathcal{A} is the normalization constant, f and \dot{f} are the two intrinsic parameters to be estimated. We have simulated 100 seconds measurement data at 1 kHz with different SNRs. The parameters of the true signal are $f_* = 100$ Hz and $\dot{f}_* = 0.2$ Hz/s. We have considered 2^6 candidate waveforms with the parameter mesh grid

$$f = \{70, 80, 90, 100, 110, 120, 130, 140\} \text{ Hz},$$

$$\dot{f} = \{-0.3, -0.2, -0.1, 0, 0.1, 0.2, 0.3, 0.4\} \text{ Hz/s}.$$

The threshold is simply chosen as $\mathcal{T} = c \cdot \max(X_k)$, where we have tried several values of the coefficient c . The SNR varies from 8 to 50 with a uniform spacing 3. For each combination of SNR and the threshold, we carried out a Monte Carlo simulation with 1000 different noise realizations. If the algorithm identifies the true signal and its true parameters, the detection is successful. The success rate is called the detection rate. Fig. 1 shows the detection rate at different SNRs and thresholds, where the color bar indicates the value of the coefficient c . The best performance is realized by setting the coefficient c around 0.5. For signals with SNR higher than 30, the detection rate of the algorithm is above 99%. Thus, the algorithm with the least number of new templates works efficiently at relatively high SNRs. However, at low SNRs, the detection rate is low. We will see whether we could improve the detection rate by slightly increase the computational cost.

Features of the algorithm—For the set of 2^N independent templates h_i , if 2^N is smaller than the number of samples in the observation data, $x_i = \langle s|h_i\rangle$ are also independent. To characterize the performance of the algorithms, we want to examine to what extent can the noise mimic a signal. Since the signal part of x_i only contributes a DC bias to its probability distribution, we can ignore the DC part and only consider the random part of x_i , which is $\langle n|h_i\rangle$. It can be shown without much effort

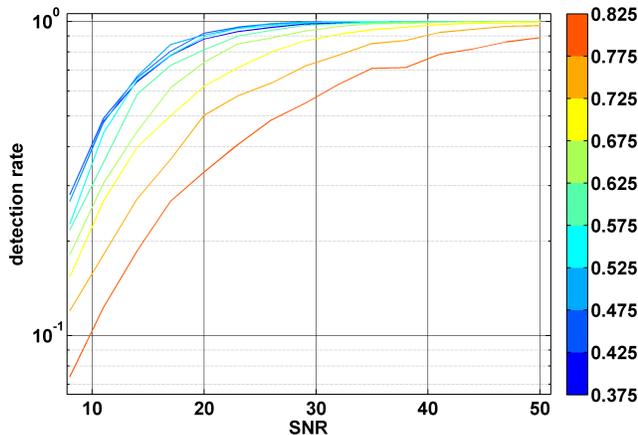


FIG. 1: Detection rate at different SNRs and thresholds. The color bar indicates the value of the coefficient c . The algorithm achieves the optimal performance, when c is around 0.5. The detection rate is above 99%, when SNR is above 30.

that the probability density function of the maximum of these 2^N random variables x_i is the following

$$p_{\max}(x) = \frac{2^N}{\sqrt{2^{2N+1}-1}\pi} \left[1 + \operatorname{erf}\left(\frac{x}{\sqrt{2}}\right) \right]^{2^N-1} e^{-\frac{x^2}{2}}, \quad (13)$$

where the error function $\operatorname{erf}(x)$ is defined as $\operatorname{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-x^2} dx$. Similarly, the probability density function of the maximum of the N random variables X_k turns out to be the following

$$p_{\max}(X) = \frac{N}{\sqrt{2^{3N}-2}\pi} \left[1 + \operatorname{erf}\left(\frac{X}{\sqrt{2^N}}\right) \right]^{N-1} e^{-\frac{X^2}{2^N}}. \quad (14)$$

For the case we considered, we have $N = 6$. The probability density functions and cumulative distribution functions of the random part of x_i and X_k are shown in Fig. 2, which tells us how large SNRs could be mimicked by pure noise. As expected, in case of X_k , the noise could mimic larger SNRs. This can also be seen from the larger standard deviation of X_k . In fact, this is the reason for the drop in the detection rate at low SNRs in Fig. 1.

Next, let us examine the role of the threshold $\mathcal{T} = \frac{1}{2} \max(X_k)$. In the previous simulations, we have six inner products X_k , ($k = 1, \dots, 6$), each corresponding to an SNR achieved by H_i . Since the detection criteria only depends on the ratio between the inner products, it is convenient to look at their pie charts. In Fig. 3, we show the pie charts for different SNRs, where the color bar represents the indices of the inner products. Take Fig. 3 (a) for instance. The inner products X_1, X_3, X_4 contribute

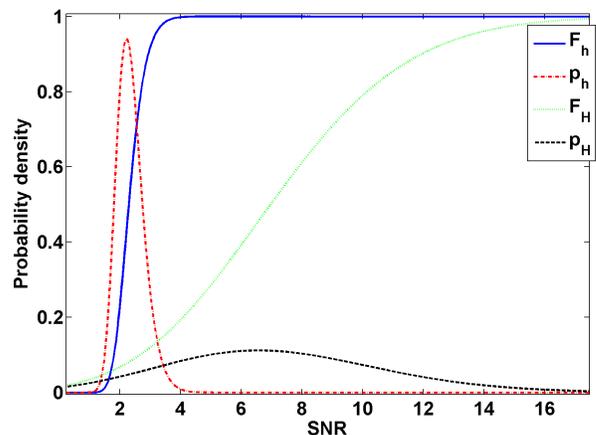


FIG. 2: The probability density functions and cumulative distribution functions of the random part of x_i and X_k , which are $\langle n|h_i \rangle$ and $\langle n|H_k \rangle$.

most part of the summation $\sum_{k=1}^6 X_k$, while X_2, X_5, X_6 are much smaller. According to the criteria we designed before, only X_1, X_3, X_4 are above the threshold. Therefore, we obtain the index $101100_2 = 44$ of the template, which most resembles the signal in the data. Similarly, Fig. 3 (b)-(e) all successfully identify the correct template in case of different SNRs. Fig. 3 (f) shows a failure case. According to the previous criteria, this pie chart gives a wrong index $101001_2 = 41$. In fact, even if one bit of the binary is wrongly determined, we end up with a completely different template (and its corresponding parameters). This is also a main reason why the detection rate at low SNRs drops so quickly.

Improve the performance of the algorithm—Now we discuss a simple and straightforward way to improve the performance of the algorithm by slightly increasing the computational cost. Let us look at the failure case in Fig. 3 (f) again. The largest inner product is X_1 , which contributes 30 percent of the entire SNR pie. The threshold, which was set to half of the largest inner product, turns out to be 15 percent. Therefore, among the six inner products, X_1, X_3 are significantly above the threshold, X_2, X_5 are significantly below, while X_4, X_6 are close to the threshold. In the end, the binary bits corresponding to X_4 and X_6 (i.e. the 4th and 6th bits) were determined wrongly, which leads to a detection failure. However, the binary bits corresponding to X_1, X_2, X_3 and X_5 are correctly determined, and we are confident about that in the blind search. In fact, we are not so confident about the bits corresponding to X_4 and X_6 , since they are just slightly above or below the threshold. If we leave these two binary bits undetermined, we end up with a binary number $101y0y_2$, where we have used y to denote undetermined bits. It implies that the true signal might match

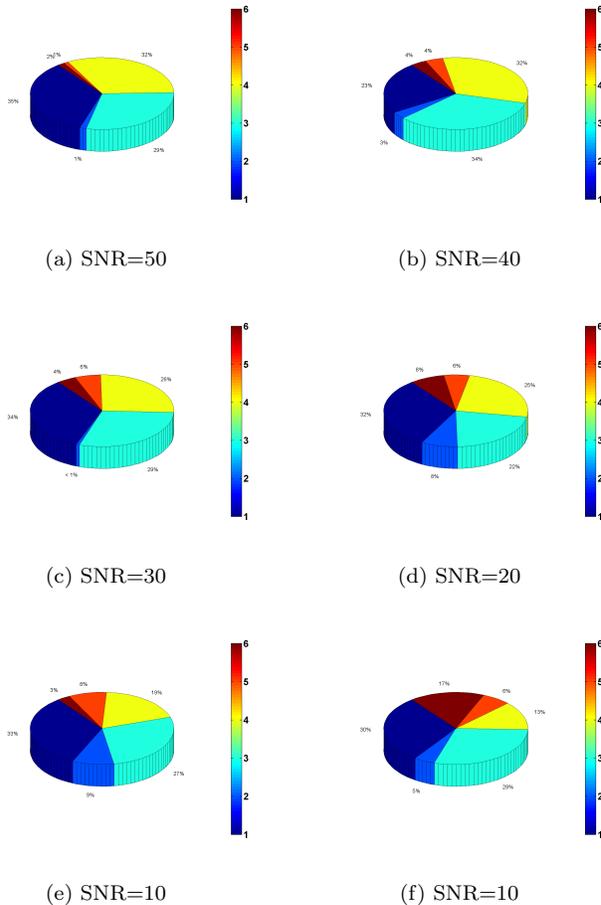


FIG. 3: Pie charts of X_k for different SNRs. (e) and (f) are for the same SNR with different noise realizations. The color bar denotes the index of X_k .

one of the four templates $101000_2 = 40$, $101100_2 = 44$, $101001_2 = 41$ and $101101_2 = 45$. By simply calculating the inner products of the data and these four templates, we will know which one matches the true signal.

Hence, we can modify the algorithm according to the above procedure. In the beginning, we calculate X_k , ($k = 1, \dots, 6$) and the threshold $\mathcal{T} = c \cdot \max(X_k)$. Then, we identify two X_k , which are closest to the threshold \mathcal{T} , and leave two binary bits corresponding to these two X_k undetermined. We determine other binary bits in the same way as before. A binary number with two unknown bits is thus constructed. It corresponds to four original templates h_i . In the end, we calculate the inner product between the data and these four templates, and detect the signal. Following this procedure, we carried out a similar simulation as before. The detection rate is plotted in Fig. 4 with different combinations of c values and SNRs. Comparing with Fig. 1, the modified algorithm has significantly improved the performance. The detection rate is increased at all SNRs. We also observe that

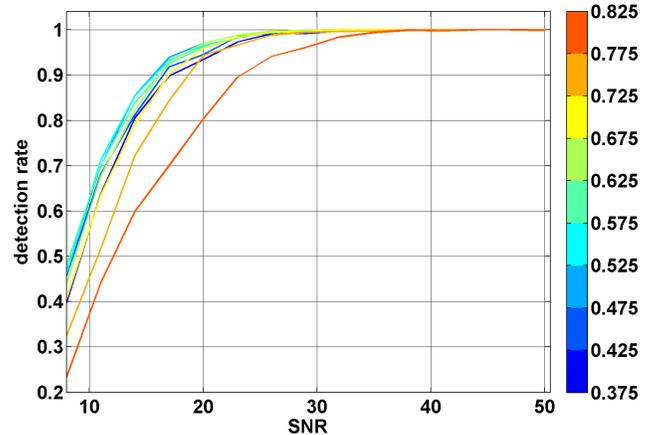


FIG. 4: Detection rate at different SNRs and thresholds. The color bar indicates the value of the coefficient c .

$c = 0.5$ is still the optimal choice. For the curve $c = 0.5$, the detection rate is 100% above SNR 30 and 96% at SNR=20. This strategy can be easily generalized by assigning a probability to each binary bit according to X_k , hence obtaining the probability of each h_i present in the data. However, this is out of the scope of the current article. We will discuss it in the future work.

Conclusion and future work—We have designed a novel algorithm for GW data analysis. Instead of using 2^N normal waveform templates, this new algorithm uses only N combinations of the original waveforms as the new templates. By calculating the inner products between these N new templates with the data and comparing these inner products with some threshold, we can construct a binary number with N bits. From this binary number, we can determine which normal template in the original template bank best matches the signal in the data, without any reconstruction process. Therefore, this new algorithm can greatly reduce the computational cost in certain circumstances. However, it requires relatively high SNRs. We have discussed a simple and straightforward way to improve the performance of the algorithm. By leaving two most unconfident binary bits undetermined and calculating four additional inner products, we can significantly improve the performance of the algorithm at low SNRs. The detection rate of the modified algorithm is 100% for 1000 different noise realizations for each SNR larger than 25. For SNR lower than 25, further improvements are demanded. We reserve that for future work.

One possible way to improve the algorithm is to construct additional H_k , ($k = N + 1, \dots$) for auxiliary use, such as to determine unconfident binary bits, to suppress the noise in X_k , etc. One can also set more sophisticated

thresholds. We have used a threshold only depending on the relative values between the inner products X_k for simplicity. A threshold also depending on the absolute values of the inner products would help, since the probability distribution of the random part of X_k depends only on the absolute SNRs.

We have only carried out simulations for a bank of nearly independent templates. In the future, we will do a simulation for an entire template bank. The correlation between templates need also to be studied, since it could be used to reduce the noise in the detection statistic.

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