Reducing the number of templates for aligned-spin compact binary coalescence gravitational wave searches using metric-agnostic template nudging

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Efficient multi-dimensional template placement is crucial in computationally intensive matched-filtering searches for Gravitational Waves (GWs). Here, we implement the Neighboring Cell Algorithm (NCA) to improve the detection volume of an existing Compact Binary Coalescence (CBC) template bank. This algorithm has already been successfully applied for a binary millisecond pulsar search in data from the Fermi satellite. It repositions templates from over-dense regions to under-dense regions and reduces the number of templates that would have been required by a stochastic method to achieve the same detection volume. Our method is readily generalizable to other CBC parameter spaces. Here we apply this method to the aligned–single-spin neutron-star–black-hole binary coalescence inspiral-merger-ringdown gravitational wave parameter space. We show that the template nudging algorithm can attain the equivalent effectualness of the stochastic method with 12% fewer templates.

I. INTRODUCTION

Compact binary coalescence systems consisting of neutron-stars and/or black-holes are key targets for the present generation of gravitational wave detectors such as Advanced LIGO [1] and Advanced Virgo [2]. The LIGO and Virgo detectors have observed a number of binary black-hole coalescence events [3–7] and recently a binary neutron-star coalescence event [8]. The detection of a coalescence of a neutron-star–black-hole binary has yet to be seen, but is increasingly likely to be observed in the future. The searches for compact binaries a priori cover a wide range of masses and spin magnitudes. A key ingredient in these searches is a suitable template bank, i.e. a collection of model gravitational wave signals which cover the desired parameter space.

For searches based on matched filtering with modeled waveforms, the traditional method of constructing a template bank was to use the parameter space metric [9, 10] for determining the spacing between adjacent templates. This method was successfully demonstrated for non-spinning systems [11, 12] and aligned-spin systems [13, 14]. In situations where the parameter space structure is not sufficiently well understood, stochastic methods are used [15–17].

Stochastic methods place templates at random points in the parameter space which are drawn from an initially chosen distribution. The chosen template is then compared with previously accepted templates, and accepted only if it is sufficiently far away each of the previously accepted templates. The procedure terminates when a certain coverage (the ratio of rejected templates over the total number of template candidates) has been achieved resulting in a final saturated template bank. These stochastic methods are more generally applicable but they are typically less efficient than the geometric methods, i.e. they require more templates than a geometric bank to achieve the same effectualness over the same parameter space. Furthermore, the construction of a stochastic template bank can be computationally demanding since, in principle, each new proposed template needs to be compared with previously chosen templates. This problem becomes particularly acute the closer the bank gets to saturation. The computational problem also becomes especially demanding when precession effects are considered as these additional degrees of freedom require a large number of templates [18, 19]. It is therefore important to consider methods of optimizing a template bank, specifically finding ways of improving effectualness for a given number of templates and reducing the computational cost.

In this paper we shall meet this computational challenge and show how stochastic methods can be improved by employing the nearest neighbor cell algorithm (NCA) [20, 21] and repositioning templates to maximize the effectualness and detection volume. It was shown in [20] that for the Fermi γ-ray pulsar search, using these methods leads to a reduction in the number of distance computations in three dimensions by about five orders of magnitude compared to other standard stochastic template bank algorithms. Here we shall apply these ideas to the Gravitational Wave (GW) Compact Binary Coalescence (CBC) problem. Specifically, we shall focus on neutron-star–black-hole (NSBH) systems since they make up 60% of the templates placed in the last LIGO-Virgo observation period [22, 23]. We expect that our method would apply to other source systems as well. We consider NSBH binaries with a black-hole of mass \( M_{\odot} \) and a neutron-star mass of \( M_{\text{NS}} \) such that \( 2M_{\odot} < M_{\text{BH}} < 16M_{\odot} \), and \( 1M_{\odot} < M_{\text{NS}} < 3M_{\odot} \). We use the inspiral-merger-ringdown-phenomenological waveform model (IMRPhenomD) [24–26] to approximate the underlying coalescence NSBH gravitational wave signal.

The plan for the paper is as follows. Sec. II briefly summarizes the necessary background material and introduces notation that will be used later. In particular we defines the maximal mismatch, effectualness and detec-
tion volume associated with a template bank. It also describes the template nudging algorithm. Sec. III B applies this to NSBH aligned spin templates. It first introduces the chirp time coordinates and calculate the template isosurfaces at some selected points in parameter space in the absence of a metric (as opposed to the method in [20, 21] which requires an analytic expression for the metric). Sec. III applies the optimization scheme to the aligned spin template bank and finally Sec. IV presents our results which are followed by concluding remarks.

II. BACKGROUND

A. Matched filtering

Matched filtering is a methodology used to determine if time series data \( x(t) \) (where \( t \) denotes time), contains some signal with parameters \( p_I \) of known form, \( h(p_I|t) \), or only instrumental noise \( n(t) \). Thus, in the absence of a signal, \( x(t) = n(t) \), and in the presence of a signal \( x(t) = h(p_I|t) + n(t) \). If the noise is stationary, we can characterize it by the single-sided power-spectral-density (PSD) \( S_n(f) \) according to

\[
\langle \hat{n}^*(f)\hat{n}(f') \rangle = \frac{1}{2}S_n(f)\delta(f-f').
\]

Here the brackets \( \langle \cdot \rangle \) denote an average expectation value over many realizations of the noise, and \( \hat{n}(f) \) denotes the Fourier transform of \( n(t) \).

The PSD is used to define the inner product between two time-series \( x(t) \) and \( y(t) \):

\[
(x|y) \equiv 4\text{Re} \int_0^\infty \frac{\hat{x}^*(f)\hat{y}(f)}{S_n(f)} df .
\]

This inner product is used to define the norm of a time series \( x(t) \) and a normalized time series \( \hat{x} \) in the usual way:

\[
||x|| \equiv (x|x)^{1/2}, \quad \hat{x} = x/||x||.
\]

For Gaussian noise, the likelihood function \( \Lambda \) is [27, 28]

\[
\log \Lambda \equiv (x(t)|h(p_I|t)) - \frac{1}{2} (h(p_I)|h(p_I)).
\]

The idealized procedure to search for a signal with unknown parameters is to compute \( \log \Lambda \) for all points (suitably discretized) in a given parameter space and to find the point where \( \log \Lambda \) is maximum. The likelihood can be analytically maximized for certain parameters (such as the initial phase \( \phi_0 \) and an overall constant amplitude) or by a Fast-Fourier transform (such as the time of arrival \( t_0 \)) (see e.g. [29]), while other parameters must be explicitly maximized over. These parameters we denote as \( \lambda_i \).

A template bank is a collection of waveforms \( \{h(p_I^j)\} \) with parameters \( \{p_I^j\} \), labeled by the index \( j \). Given a template bank, we would like to know how effective it is in recovering a given signal \( h_I \). This is quantified in terms of a number, namely the fitting-factor (FF) defined as,

\[
FF(h(p_I), \{h(p_I^j)\}) \equiv \max_j \mu(h(p_I^j), h(p_I)), \quad (7)
\]

where

\[
\mu(h(p_I^j), h(p_I)) \equiv \max_{t_0, \phi_0} (h(p_I^j)|h(p_I)(t_0, \phi_0)) \quad (8)
\]

is the match between \( h(p_I^j) \) and \( h(p_I) \). \( \mu(h(p_I^j), h(p_I)) \) represents the fraction of the optimal signal-to-noise ratio (SNR) of signal \( h(p_I^j) \) captured by the template \( h(p_I) \).

As matter of notational convenience, we define the mismatch as

\[
m_m(h(p_I^j), h(p_I)) \equiv 1 - \mu(h(p_I^j), h(p_I)) \quad (9)
\]

which represents the fraction of the optimal SNR of signal \( h(p_I^j) \) not captured by the template \( h(p_I^j) \). The fitting factor depends on a particular template bank and a particular signal \( h_I \). Since we will compute this for a fixed template bank, for notational convenience we usually drop its dependence on \( \{h(p_I^j)\} \) and write \( FF(h_I) \).

The loss in SNR can be quantified by the match between a signal and the nearest template and can be formulated geometrically [9, 30]. The match between nearby points in parameter space can be approximated as

\[
\mu(\hat{h}(\lambda), \hat{h}(\lambda + d\lambda)) = 1 - g_{ij}d\lambda^i d\lambda^j + \ldots \quad (10)
\]

with the metric

\[
g_{ij} = -\frac{1}{2} \left. \frac{\partial^2 \mu(\hat{h}(\lambda), \hat{h}(\lambda'))}{\partial\lambda_i^j} \right|_{\lambda' = \lambda}. \quad (11)
\]

This metric\(^1\) is useful in quantifying the density of templates. The higher the metric determinant, the higher the required template density is for a given allowed SNR loss (which corresponds to a given maximal mismatch). In the case of the parameter space considered in Section III B 1, we shall see that the methods presented in [20, 21] must be generalized to deal with the challenge of a priori not having an explicit expression for a metric.

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\(^1\) For Gaussian stationary noise, one can show that the metric \( g_{ij} \) is equivalent to constructing the scalar product \( \frac{1}{2} \left( \frac{\partial h}{\partial \lambda_i} \right| \frac{\partial h}{\partial \lambda_j} \right) \) and projecting out the parameters \( t_0 \) and \( \phi_0 \).
For aligned waveforms, there is an analytic expression for the inspiral-only metric for the NSBH parameter space \( [13, 14] \). However there is no known analytic expression for the metric of the inspiral-merger-ringdown (IMR) NSBH precessing parameter space \( [31] \). Numerical approximations to the metric can be ill-conditioned in curved regions of the parameter space. Therefore, a robust method for repositioning, nudging, templates must depend on numerical mismatch calculations.

### B. Template bank Effectualness

An optimally placed template bank minimizes the number of templates for a given volume and maximizes the detection volume for a given prior distribution of sources. We consider the prior distribution of NSBH binaries selected from a uniform component mass and aligned BH spin distribution. This distribution allows us to quantify the effectualness of recovering a range of NSBH injections independent of the underlying astrophysical distribution.

We quantify the effectualness of the template bank by determining the minimum recovered fitting factors of 99.9\% of a population of an injection set. To quantify the relative improvement of two or more CBC template banks, we calculate the relative improvement in detection volume \( [32] \). The detection volume, \( V \), is assumed to be proportional to the sum of the cube of the product of the optimal SNR of the injections, \( \rho_i \), with the fitting factor, \( FF_i \), obtained from attempting to recover a set of injected NSBH signals,

\[
V \propto \sum_i (FF_i \rho_i)^3. \tag{12}
\]

By taking the ratio of the detection volumes of two template banks, \( V_2 \) vs \( V_1 \), we obtain the relative detection volume quantifying which template bank will perform better in a search.

### C. Template nudging algorithm

A stochastic placement algorithm is unlikely to place templates optimally and will place more templates than are necessary to achieve the same effectualness \( [17] \). We therefore employ an algorithm (see Figure 1) to move templates slightly, templates nudging, facilitating the rearrangement of templates in a template bank into a configuration that improves the template bank’s effectualness and detection volume without the addition of more templates.

In Section III we apply the template nudging algorithm to the CBC parameter space by adapting a version of the Neighboring Cell Algorithm (NCA) method \( [20] \) that does not require an analytic expression of the metric and will work in parameter spaces where numeric approximations of the generalized metric are ill-conditioned. In our application of the method, the three dimensional regions covered by individual templates, template isosurfaces, within a pre-specified maximal mismatch (i.e. 3\%) must be determined numerically.

The following list outlines the procedure for nudging a template bank in the absence of a metric.

1. Select a template \( T \).
2. Find a set of points that are uniformly distributed distributed on the boundary of \( T \)’s isosurface.
3. Check whether each of these points is inside a neighboring template’s isosurface. If there is overlap, the considered boundary point gets zero weight. If not, it gets unit weight (i.e. the template will not be nudged toward the adjacent overlapping template isosurface).
4. If a boundary point is outside of the considered parameter space this point also gets zero weight.
5. The boundary points are averaged together using these weights into a barycenter.
6. The template is nudged (i.e. coordinates are perturbed) in a direction determined by the barycenter offset relative to the unweighted barycenter of the boundary points and a maximum relative amount, \( \epsilon \) (the pre-specified template nudge factor, the fractional distance between the original template center and the closest isosurface point beyond which templates cannot be repositioned).

This algorithm can be applied to any under-covered bank regardless of the method used to create that template bank (e.g. geometric, stochastic, or hybrid template bank placement methods which place templates with geometric and stochastic methods) and will improve the effectualness and detection volume. This provides a convenient way of enhancing any existing template bank construction method.

#### 1. Neighboring cell algorithm

In order to efficiently cover the parameter space with non-overlapping cells, the cells should be approximately the same size as the area of the chirp time volume covered by an individual template, as determined by the desired maximal mismatch between neighboring templates (3\%). In the simplest example these can be hyper spheres with a regular stacking in any choice of coordinates. To produce a list of the nearest neighboring templates we implement the following procedure.

1. Each cell is uniquely indexed.
2. Given a cell index, the indices of neighboring cells can be computed easily or can be stored in a table. Two cells are neighboring if at least one point exists
FIG. 1. Flowchart of the template nudging algorithm procedure. Preparation steps are listed in the frame on the left and the template nudging procedure is listed in the frame on the right. In Section IV we set “N” = 100 for the termination condition shown in the diamond on the lower right.

that has maximal mismatch of < 3% to each of the two cell regions.

3. Template parameters are mapped to cell indices by a cell table. The parameters of each template lie in the region of the corresponding cell.

4. Given a position of a template, the index of a cell can easily be computed by any kind of hash algorithm or a binary search. In Euclidean space and with a hyper-cubic cell lattice this can be achieved by using rounding or truncating operations on the position values of the templates.

5. A second table stores template indices and template parameters in memory.

6. Given a template, one finds all templates in the vicinity by collecting the templates in the corresponding cell and all neighboring cells. Therefore relatively few mismatches have to be computed when placing a template as opposed to other stochastic placement algorithms.

III. APPLICATION TO THE CBC ALIGNED-SPIN NSBH TEMPLATE BANK

A. Aligned-spin binaries

The parameter space of aligned-single-spin NSBH gravitational wave signals considered in this paper can be represented with three physical dimensions \( \{ M_{BH}, M_{NS}, S_{BH} \cdot \hat{L} \} \), i.e. the masses of the black hole \( M_{BH} \) and neutron star \( M_{NS} \), and the component of the black hole spin \( S_{BH} \) along the orbital angular momentum \( \hat{L} \) of the binary (\( \hat{L} \) is the unit vector along \( L \)). We require further that \(-M_{BH}^2 < S_{BH} \cdot \hat{L} < M_{BH}^2\) consistent with a Kerr black hole in general relativity. The BH spin representation can also be expressed in a dimensionless
\[
\chi_{BH} = \frac{S_{BH} \cdot \hat{L}}{M_{BH}^2}.
\] (13)

We only consider binary systems with quasi-circular orbits and we ignore any parameters associated with the internal structure of the neutron-star. We ignore neutron-star spins since they are expected to be small, but the black-hole spin will be allowed to take any magnitude which is meaningful in the Kerr metric and aligned/anti-aligned to the orbital angular momentum \(33\).\(^2\) For computational simplicity, we transform these three physical parameters into the Post Newtonian [PN] "chirp time" \(34\) coordinates, \(\{\tau_0, \tau_2, \tau_3\}\). In these coordinates the underlying mismatch metric is approximately flat at low frequencies.

B. Isosurface geometry in chirp time coordinates

Template placement efficiency is largely dependent on the geometry of the regions covered by individual templates, (i.e. the template isosurfaces). Using the physical parameters of the aligned-spin NSBH parameter space \(\{M_{BH}, M_{NS}, \chi_{BH}\}\), these template isosurfaces are non-uniform. An ideal coordinate system for template placement would yield isosurfaces that are uniformly spherical at any point in the proposal distribution. Isosurfaces that have curved or sharp edges are computationally challenging to model and tend to create holes, insufficiently populated regions in the template bank.

Since we were unable \textit{a priori} to determine an ideal coordinate system for the placement of NSBH IMRPhenomD templates, we choose three chirp time coordinates \(\{\tau_0, \tau_2, \tau_3\}\) \(35, 36\) that have been demonstrated to flatten out the TaylorF2 template bank \(13, 32\). These chirp times are defined with the following conventions where \(f_0\) denotes a reference frequency, here chosen as 30Hz:

\[
\begin{align*}
\tau_0 &= \frac{5}{256} \frac{(\pi f_0)^{-\frac{5}{2}} (M_{BH} + M_{NS})^\frac{3}{2}}{M_{BH} M_{NS}} \quad (14) \\
\tau_2 &= \frac{5}{256} \frac{M_{BH} + M_{NS}}{M_{BH} M_{NS}} (\pi f_0)^{-\frac{3}{2}} \\
&\times \left[ \frac{743}{336} + \frac{11}{4} \frac{M_{BH} M_{NS}}{(M_{BH} + M_{NS})^2} \right] \quad (15) \\
\tau_3 &= \frac{128}{(M_{BH} + M_{NS})^\frac{3}{2}} (\pi f_0)^{-\frac{3}{2}} \\
&\times \left[ 16\pi - \frac{\chi_{BH}}{6} \left( \frac{19 M_{BH} M_{NS}}{(M_{BH} + M_{NS})^2} \\
&\quad + \frac{113 M_{BH}}{M_{BH} + M_{NS}} + 94 \right) \right] \quad (16)
\end{align*}
\]

1. Template isosurfaces

In these three chirp time coordinates, \(\{\tau_0, \tau_2, \tau_3\}\), the regions of the parameter space covered by individual templates are non-ellipsoidal, thin, and therefore difficult to model analytically (see Figure 2 and 3).

However, the IMRPhenomD waveform is least sensitive to perturbations in the \(\tau_2\) degree of freedom. We found that individual cross-sections of these isosurfaces in the remaining two degrees of freedom can be modeled by two dimensional ellipsoids. We compensate for the lack of a three dimensional isosurface metric by modeling the three dimensional isosurface at a select number of cross-sections where the two dimensional projection of the isosurface is easier to model.

A further computational challenge is determining the scope of individual ellipsoid isosurface cross-section in the fewest computations possible. These ellipsoids are

\(^2\) It is possible to approximate the effect of multiple spins by representing the spins as one effective spin term. Gravitational waveforms like IMRPhenomD use single-effective-one-body-spins to model CBC binaries with two spinning component masses: \(\{M_{BH}, S_{BH}\}, \{M_{NS}, S_{NS}\}\). This effective spin is expressed as a single component mass weighted term, \(S_{EFF} = (M_{BH} S_{BH} + M_{NS} S_{NS})/(M_{BH} + M_{NS})\).
often very stretched out in these coordinates, therefore we implement the following method to select points.

The three dimensional mismatch isosurface can be obtained by computing the mismatch isosurface of individual cross-sections by keeping $\tau_2$ constant for each cross-section. The isosurface is hence a ring on the two dimensional sub-manifold parametrized by $\tau_1$ and $\tau_3$. In the first step we find the point on this plane which has the smallest mismatch with our considered template by applying a simplex amoeba gradient-free-downhill method [37]. The routine starts by finding a point on the sub-manifold which is in the allowed parameter space. The downhill method starts from there. Since the template volume ranges in the $\tau_2$ direction from one end to the other we always find a set of points on our considered $\tau_2$ plane which have mismatches with our template smaller than the allowed critical mismatch of 3%. These points are enveloped by the template isosurface and are described by:

$$P \equiv \{ \tilde{p}_i | mm (\tilde{p}_i, \bar{t}) \leq 3\% \},$$

where $\bar{t}$ is the vector of the template in the parameter space. Usually we find 12 to 25 points by using the simplex amoeba gradient-free-downhill method. We use these “inside” points to obtain a local approximation of the metric on this surface describing the distance between the point of the maximal mismatch to the points found by the simplex method. The description by a metric is not quite correct since the mismatch of the minimal point is much smaller than 3% but not zero. If this mismatch becomes significantly larger, then one might think about adding a constant to Eq. 18. However, for deriving the isosurfaces the following method works sufficiently well. We expand the metric starting with quadratic terms:

$$\tilde{m} = \sum_{p=2}^{\infty} \sum_{i_1 \geq i_2 \geq \ldots \geq i_p} \gamma_p (i_1, i_2, \ldots, i_p) \prod_{l=1}^{p} d_{i_l} ,$$

$D$ is the dimension of our manifold, in our case $D = 2$ and $d_i$ are the components of the distance vectors. In the following we restrict ourselves to the second order expansion. For each point in the set, $P$ has a real mismatch $m_k$ and the mismatch is approximated by the metric $\tilde{m}_k$. We choose the metric components $\gamma_{2(i,j)}$ such that the quantity $\chi^2 = \sum_k (m_k - \tilde{m}_k)^2$ is minimized.

The component $d^{k}_{i_l}$ is the $i_l$th component of the $k$th distance vector towards our minimal point $p_0$ on the plane. We minimize the $\chi^2$ with respect to the metric components:

$$\frac{\partial \chi^2}{\partial \gamma_{2(o,p)}} = \sum_k \left[ m_k - \sum_{i_1 \geq i_2} \gamma_2 (i_1, i_2) d_{i_1} d_{i_2} \right] d_{o} d_{p} = 0.$$  

This set of equations can be described by

$$\Rightarrow \sum_k \left[ \begin{array}{c} D_{00}^k \ D_{10}^k \\ D_{01}^k \ D_{11}^k \end{array} \right] \left[ \begin{array}{c} D_{00}^k \ D_{10}^k \\ D_{01}^k \ D_{11}^k \end{array} \right] \left[ \begin{array}{c} \gamma_2 (1,1) \\ \gamma_2 (2,1) \end{array} \right] = \sum_k m_k \left[ \begin{array}{c} D_{00}^k \\ D_{10}^k \\ D_{01}^k \end{array} \right] ,$$

where $D_{i_l}^k = (d^{k}_{i_l})^{2-i}(d^{k}_{i_l})^i$. We invert $D$ and obtain the second order expansion values for the metric $\gamma_{2(i,j)}$. This procedure can be expanded to arbitrarily high orders of the metric expansion. We can now approximate the metric in a quadratic form:

$$g = \left[ \begin{array}{c} \gamma_2 (1,1) \gamma_2 (1,2) \\ \gamma_2 (1,1) \gamma_2 (2,2) \end{array} \right] .$$
We compute eigenvalues $v_i$ and eigenvectors $\vec{c}_i$ of this metric and compute a set of $N$ points approximately in the vicinity of the isosurface. $I = \vec{p}_i$ with $\vec{p}_i = \sin(2\pi i/N)\vec{c}_1\sqrt{0.03/v_1} + \cos(2\pi i/N)\vec{c}_2\sqrt{0.03/v_2}$. These points are not equidistant but sufficiently well distributed for our purposes.

We shift the points in the radial direction with respect to the center point such that the points have a mismatch of exactly 3% using the Newton-Raphson method.

We compute the barycenter of our shifted set of points and repeat the metric approximation method and the shifting to get an even better set of points.

The scheme is applied for all distinct $\tau_2$ planes.

2. Cell Structure

In order to apply the NCA method, we provide an appropriate cell structure with the following properties:

- the cells are sufficiently small
- templates within a cell can reach the neighboring cells, but not next-nearest neighboring cells
- a cell index can be easily and quickly computed knowing the template parameter space points.

Since each template spans the entire parameter space in the $\tau_2$ direction, splitting the parameter space in this direction is not possible. On the other hand, the templates have a very small size in the $\tau_0$ direction so we can safely split the parameter space in this direction. We split the parameter space in the $\tau_0$ direction into 300 slices. We mapped the $\tau_0$ coordinates into a unit length parameter space $x = (\tau_0 - 4)/50$. The templates are curved in the $\tau_3$ direction, thus further splitting is not directly feasible. To get the templates in a compact form we applied an ad hoc transformation $y = ((\tau_2/\tau_0 - 20 + 2.5/\tau_2)\tau_0 - 480)/2500$. The new coordinate $y$ is normalized and ranges from 0 to 1. We split the parameter space into 40 slices in this direction and obtained a rectangular grid in the $x$-$y$ plane with 12,000 cells.

We tested this setup with a set of randomly distributed points. In the limit of having only one cell we compute all mismatches smaller than 3% correctly. If the cells become too small some of mismatches smaller than 3% will not be detected. This happens if the cells are so small that overlapping templates are not in neighboring cells anymore, but, for instance, in next-nearest neighboring cells.

IV. RESULTS

To test the template nudging algorithm we seeded it with a geometric lattice TF2 template bank containing 174,000 templates. Both the initial seed bank and the nudged bank were then tested against a set of 20,000 random IMRPhenomD injections. For the nudged bank only 3% of injections had a fitting factor less than 97%, compared to 10% for the original seed geometric bank (see Figure 4). For comparison, building a stochastic bank targeting 97% minimal match (i.e. 1—maximal mismatch), required 220,000 templates, with only 0.1% of templates having a fitting factor less than 97% for the same injection set.

For this comparison matches were calculated between $30 - 1024$Hz using a PSD [38] built from the harmonic mean of the Hanford and Livingston PSDs taken within a few days of GW150914 and thus comparable to what shown in [3]. The stochastic bank was generated using lalapps_cbc_sbank [23, 39–42], using with a convergence criteria of rejecting 97% trial templates. The stochastic algorithm is able to use IMRPhenomD templates directly to calculate matches. The TaylorF2 lattice was generated using pycbc_geom_aligned_bank [13, 23, 39, 41], constructed with a two dimensional lattice such that each lattice point had a maximal mismatch no larger than 3%. As can be seen in Figure 4, this target was not attained for the IMRPhenomD test signals. A known metric is required to determine the geometric lattice. Since no metric is known for the IMRPhenomD templates we used TaylorF2 3.5PN waveforms [42–44] for which an analytical metric can be calculated. Although the lattice locations were calculated using TaylorF2 waveforms, these were subsequently swapped out with IMRPhenomD templates when testing the effectualness and this results in a loss of effectualness.

The template nudging algorithm was applied for 100 iterations using one cross-section with at least 16 surface
Within each iteration, templates were nudged in parallel for the purposes of reducing computation time. Studying the effect of the template nudge factor on the convergence of the method revealed that a template nudge factor of 5% produced the largest improvements in the template bank’s effectualness in the first iteration relative to template nudge factor of 1% or 0.5%. However, in later iterations the variance of the improvement of the recovered fitting factors was also higher and a poorer fitting factor was recovered in the anti-aligned high mass regions of the bank relative to those produced by template nudge factor of 1%. We used a composite approach to improve the convergence of the method, wedding the advantages of using a bigger template nudge factor in the first iterations to the advantages of the greater precision of smaller template nudge factors in later iterations. Hence, we applied the template nudge algorithm in two batches of 50 iterations per template nudge factors: 5% and 1%. After 100 nudging iterations, the coverage had improved by 3% and the relative detection volume had increased by 0.69%.

In order to extract an additional 1% of coverage and an additional 0.80% relative detection volume beyond what is achieved with the nudging, we polished the nudged template bank by adding 20,000 templates via a final stochastic placement to fill in any remaining holes. This produced a bank with overall fitting factors comparable to the stochastic bank, but with only 194,000 templates compared to the 220,000 templates required via the `lalapps_cbc_sbank` algorithm. Therefore it is possible to achieve equivalent template bank effectualness and detection volume with 26,000 fewer templates than would be required by the purely stochastic method.

Figure 7 compares the result of injecting 40,000 IMRPhenomD signals into the nudged TF2 geometric lattice and the original TF2 geometric lattice. The nudged TF2 geometric lattice has a more even distribution of recovered fitting factors. The template nudging algorithm improves regions with poor fitting factors by repositioning templates from over-covered regions. In principle this allows the attainment of a desired minimum fitting factor across the entire parameter space with fewer templates. However, the template nudging algorithm as currently conceived does not remove templates from the bank. In some cases the nudging can result in templates from over-covered regions being nudged to nearby boundaries and piling up there. An example of this is shown in the lower right panel of Figure 7. This highly anti-aligned spin region of the NSBH parameter space, indicated by the yellow boxes in the other panels, is over-covered by the original TF2 geometric lattice, obtaining recovered fitting factors close to \( \sim 100\% \). In this region the excess templates are nudged to the boundaries and build up there, without reducing the fitting factors.

FIG. 5. The relative detection volume of the nudged bank versus the stochastic bank, 100% \( \times V_{\text{nudged}}/V_{\text{stochastic}} \), was tracked every five iterations of the template nudging algorithm using the same set of 20,000 injections drawn from a uniform component mass and aligned BH spin distribution. The improvement settles to a constant value after approximately 40 iterations.

FIG. 6. The relative detection volume improvement of the nudged bank versus the original TaylorF2 (TF2) geometric lattice bank, 100% \( \times (V_{\text{nudged}}/V_{\text{original}} - 1) \), was tracked every five iterations of the template nudging algorithm using the same 20,000 injections drawn from a uniform component mass and aligned BH spin distribution.

points to approximate individual template isosurfaces for the nudging calculations. The effectualness and improvement in the relative detection volume between successive intermediate nudges was quantified by recovering a set of 20,000 IMRPhenomD injections drawn from a uniform component mass and BH spin distribution (see Figures 5 and 6).
FIG. 7. Comparison of fitting factors for the template nudging algorithm. Coordinates are the chirp times given in Section III B. The top row shows the fitting factors for the nudged TF2 geometric lattice (left) and the original TF2 geometric lattice (right). The fitting factors are seen to be more uniform across the nudged bank. The difference in fitting factor is displayed in the bottom left panel. The regions of improvement (in red) are seen to correspond largely with the regions of insufficient fitting factor in the original geometric lattice. The lower right panel shows the template locations for three different template distributions in the highly anti-aligned spin region denoted by the yellow box in the other panels. In this region the nudging pushes templates to nearby boundaries and does not reduce the number of templates in a region that is over-covered by the geometric lattice.

V. CONCLUSIONS

We have shown how the number of gravitational wave templates needed to search a region of parameter space can be reduced by repositioning templates. In particular, we successfully implemented a method to reduce the number of templates required by the algorithm lalapps_cbc_sbank to cover an NSBH single-aligned spin parameter space. This resulted in a 12% reduction in the number of templates. For comparison, the hybrid method (utilized to build the O1 bank [22] and the O2 bank [23]) required 5% fewer templates than the stochastic method when used to build a binary-black-hole bank [41]. Given that 60% of the templates in the O2 CBC template bank are in the NSBH mass range considered in this paper, applying the template nudging algorithm in this subspace alone would already reduce the size of the bank placed on the entire mass range by 7% without sacrificing effectualness. Assuming that the percentage reduction in the number of templates is uniform across the parameter space, then we could expect an overall reduction of approximately 50,000 templates from the 400,000 template O2 bank.

The template nudging is seen to be most effective at repositioning templates from over-covered regions to under-covered regions. This is less effective in regions very close to boundaries of the desired parameter space. In some cases the shifted templates can accumulate near the boundaries as seen in the lower right panel of Figure 7. Once a template is nudged such that one or more of its isosurface points exceeds the border of the bank, there is currently no mechanism to nudge the template away.
from that border or to remove it entirely.

For the nudging the template bank is split into two-dimensional planes and the nudging takes place within each plane. The algorithm does not nudge templates between the planes and thus if there is a hole in one of the planes the method is not able to fill it with excess templates from another plane. Such banks may require additional templates after nudging in order to match the effectualness and detection volume that can be produced by a stochastic method. In our example, this would occur specifically when these holes are in regions of the bank which require nudging the $\tau_2$ coordinate since this version of the template nudging algorithm only nudges templates on the $\tau_0 - \tau_3$ plane. It is non-trivial to remove this restriction since the template isosurface cross-sections in the $\tau_0 - \tau_2$ and the $\tau_2 - \tau_3$ planes are non-ellipsoidal and would require modifying how the algorithm samples the isosurface boundary points.

Further work is needed to adapt this method to the entire aligned-spin CBC parameter space. Currently the algorithm only places templates in the NSBH and Binary Neutron Star (BNS) mass range. Extending this region would require readdressing the following two points: 1) defining the borders of the targeted chirp-time parameter space and 2) obtaining a uniform transformation across the targeted chirp-time parameter space to the two dimensional grid needed to apply our modified NCA method. In this paper, these two issues were solved ad hoc, but we believe the method is readily adaptable to other parameter regions.

An additional computational challenge is scaling the algorithm to nudge template banks with millions of templates. The preliminary application of the template nudging algorithm to the aligned-spin CBC parameter space required the use of a high-throughput computing cluster. Generating the template isosurface cross-sections in the absence of a general mismatch metric (analytic or otherwise) is particularly computationally expensive. The template nudging algorithm is considerably faster if there is a reliable (and computationally efficient) way to calculate the metric. By nudging purely BNS banks (where there is a known analytic expression of the metric), we were able to nudge an under-saturated stochastic bank on a single Lenovo Thinkpad T450s Ultrabook laptop. Similarly, the template banks used for Fermi $\gamma$-ray binary pulsar searches in [21] (which contained tens of millions of templates) could be constructed by the NCA method in a few hours on one HUAWEI RH1288 v3 Server with two 14 core CPUs since there is an analytic expression for the metric. Therefore, it is highly desirable to obtain a computationally efficient way to calculate the metric in order to apply the template nudging algorithm to larger CBC parameter spaces, or detectors with improved sensitivities.

The so-called $\theta$ coordinates produce a computationally efficient method for calculating the numeric approximation of the CBC mismatch metric [40, 45]. These coordinates are only dependent on the masses and effective spin and may produce more uniform template isosurfaces, better suited to the template nudging algorithm. Running the template nudging algorithm with these flatter coordinates may lower the computational cost and further reduce the final number of required templates for the CBC template bank.

While our current implementations of the metric agnostic template nudging algorithm are computationally inefficient, they are still a lot more flexible than the original metric dependent NCA method [20]. As long as there is a coordinate system in which cross-sections of individual template isosurfaces can be reduced to a collection of two dimensional ellipsoids, this bootstrap metric construction method is completely generalizable to any current or future CBC template bank parameter space. This makes it a potentially versatile option for building higher dimensional template banks that include the effects of precession, tidal deformation, eccentric orbits, and higher order modes.

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